

# **Stage Senior 2010 – Livello Advanced**

**Stampato integrale delle lezioni**

Autori vari



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# SENIOR 2010 - ADVANCED

Titolo nota

06/09/2010

**SERIE**

$$\sum_{n=0}^{\infty} a_n$$

$$S_m = a_0 + a_1 + \dots + a_m$$

$$S_m \rightarrow ? \quad L \in \mathbb{R} \text{ o } \pm\infty$$

NON ESISTE LIM

Se  $a_n \geq 0$  ci sono solo 2 poss:  $L \in \mathbb{R}$ ,  $+\infty$ 

Esempio 1 Serie telescopica  $\sum_{n=1}^{\infty} \frac{1}{n(n+1)} = \sum_{n=1}^{\infty} \left( \frac{1}{n} - \frac{1}{n+1} \right) = 1$

Si dim. che  $S_m = 1 - \frac{1}{m+1}$

Esempio 2  $\sum_{n=1}^{\infty} \log\left(1 + \frac{1}{n}\right) = \sum_{n=1}^{\infty} [\log(n+1) - \log n]$

$$S_m = \log(m+1) - \log 1 \rightarrow +\infty$$

Esempio 3  $\sum_{n=0}^{\infty} a^n$  (a parametro) GEOMETRICA

$$S_m = \frac{a^{m+1} - 1}{a - 1}$$

- se  $a \geq 1$  DIV. a  $+\infty$
- se  $a \in (-1, 1)$  Conv. a  $\frac{1}{1-a}$
- se  $a \leq -1$  Non ha limite

Esempio 4  $\sum_{n=1}^{\infty} \frac{1}{n^a}$  (armonica generalizzata)

- converge se  $a > 1$
- diverge a  $+\infty$  se  $a \leq 1$

**Dim 1** Con  $a = 1$  diverge

$$1 + \frac{1}{2} + \frac{1}{3} + \frac{1}{4} + \frac{1}{5} + \dots + \frac{1}{8} + \dots \geq$$

$$1 + \underbrace{\frac{1}{2}}_{\frac{1}{2}} + \underbrace{\frac{1}{4} + \frac{1}{4}}_{\frac{1}{2}} + \underbrace{\frac{1}{8} + \dots + \frac{1}{8}}_{\frac{1}{2}} + \dots$$

Formalmente:  $S_{2^m} \geq \frac{m}{2}$

A maggior ragione: con  $a < 1$  diverge

Con  $a=2$  converge visto ieri  $S_m \leq 2 - \frac{1}{m}$

Oppure  $\frac{1}{n^2} \leq \frac{1}{n(n-1)} \Rightarrow \sum_{n=2}^{\infty} \frac{1}{n^2} \leq \sum_{n=2}^{\infty} \frac{1}{n(n-1)} = 1$   
↑  
VISTA SOPRA

A maggior ragione: con  $a \geq 2$  converge

Btw:  $\sum_{n=1}^{\infty} \frac{1}{n^2} = \frac{\pi^2}{6}$  (Non banale)  $\sum_{n=1}^{\infty} \frac{1}{n^{2k}}$  c'è formula ( $k \in \mathbb{Z}$ )

Restano gli  $a \in (0, 1)$

Dim 2 Idea di ieri:  $S_m^a = \frac{1}{1^a} + \dots + \frac{1}{m^a}$   $S_m^a \leq M_a - \frac{N_a}{m^{a-1}}$

Cerco  $M_a$  ed  $N_a$  in modo da fare induzione

P.I.  $S_{m+1}^a = S_m^a + \frac{1}{(m+1)^a} \leq M_a - \frac{N_a}{m^{a-1}} + \frac{1}{(m+1)^a} \leq$   
↑  
Hp  
 $\leq M_a - \frac{N_a}{(m+1)^{a-1}}$   
↑  
Ho pe

serve  $\frac{N_a}{(m+1)^{a-1}} + \frac{1}{(m+1)^a} \leq \frac{N_a}{m^{a-1}}$

$$N_a \left( \frac{1}{m^{a-1}} - \frac{1}{(m+1)^{a-1}} \right) \geq \frac{1}{(m+1)^a}$$

$$N_a \geq \frac{m^{a-1} (m+1)^{a-1}}{(m+1)^a [(m+1)^{a-1} - m^{a-1}]} \sim \frac{m^{a-2}}{m^{a-2}}$$

serve che sia LIMITATO

$$\frac{m^{a-1}}{m \left(1 + \frac{1}{m}\right)^{a-1} \left[ \left(1 + \frac{1}{m}\right)^{a-1} - 1 \right]} = \frac{1}{m \left[ \left(1 + \frac{1}{m}\right)^{a-1} - 1 - \frac{1}{m} \right]}$$

$(1+x)^\alpha \geq 1 + \alpha x$

VERA per  $x \geq 0$  e  $\alpha \geq 1$  (se è così è ok solo per  $a \geq 2$ )

$$\geq \frac{1}{m \left( 1 + \frac{a}{m} - 1 - \frac{1}{m} \right)} = \frac{1}{a-1} \text{ ok !!}$$

$N_a \rightarrow \infty$  per  $a \rightarrow 1$  (ottimo!)

Dim. 3 Criterio di condensazione di Cauchy

$$\sum_{n=1}^{\infty} a_n \quad \text{Ipotesi: } \begin{cases} \bullet a_n \geq 0 \\ \bullet a_n \text{ debolm. decresc.} \end{cases}$$

Allora  $\sum a_n$  converge  $\Leftrightarrow \sum 2^n \cdot a_{2^n}$  converge

$\uparrow$   
indice = potenza di 2

$$a_1 + a_2 + a_3 + a_4 + \dots + a_8 + \dots \geq a_1 + a_2 + a_4 + a_8 + \dots + a_8$$

$$= a_1 + a_2 + 2a_4 + 4a_8 + 8a_{16} \quad \text{se serie } \geq$$

$$a_1 + a_2 + a_3 + \dots + a_8 \leq \overbrace{a_1 + a_2 + a_2}^{2 \cdot 3} + \overbrace{a_4 + \dots + a_4}^{4 \cdot 7} + \overbrace{a_8 + \dots}^{8 \cdot 15}$$

$$a_1 + 2a_2 + 4a_4 + 8a_8 + \dots$$

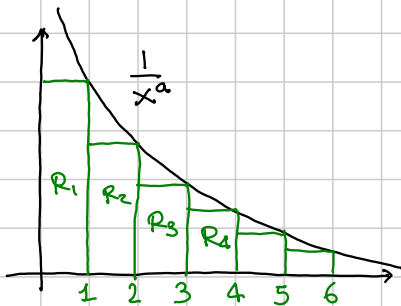
Applicazione:

$$\sum_{n=1}^{\infty} \frac{1}{n^a} \text{ converge } \Leftrightarrow \sum_{n=1}^{\infty} 2^n \cdot \frac{1}{2^{an}} = \sum_{n=1}^{\infty} \left[ \frac{2}{2^a} \right]^n$$

$$= \text{geometrica che converge } \Leftrightarrow \frac{2}{2^a} < 1$$

$$\Leftrightarrow a > 1.$$

— o — o —

Dim. 4 Confronto serie integrali (convergenza per  $a > 1$ )

$$\text{Area}(R_i) = \frac{1}{i^a}$$

$$R_2 + R_3 + \dots + R_m \leq \int_1^m \frac{1}{x^a} dx$$

area sotto il  
grafico da 1 ad m

$$\int x^a dx = \frac{1}{a+1} x^{a+1}$$

$$a = -a$$

$$= \frac{1}{1-a} \left[ \frac{1}{x^{a-1}} \right]_1^m$$

$$= \frac{1}{1-a} \left[ \frac{1}{m^{a-1}} - 1 \right]$$

$\downarrow$   
0 se  $a > 1$

IMO 1991? - 6 Determinare se esiste una successione  $x_n$  LIMITATA t.c.

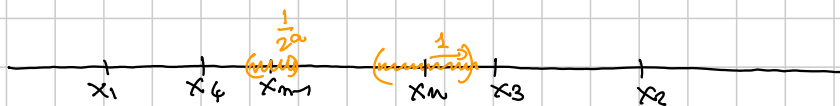
$$|x_m - x_n| \geq \frac{1}{|m-n|^a} \quad \forall m \neq n$$

(a parametro  $> 1$ )

Oss.1 Se non ci fosse limitata sarebbe banale!

Oss.2 Posso prendere  $x_n$  crescente? No! Ognuno dista almeno 1 dal precedente!

Tenta per induzione. Dati  $x_1, \dots, x_n$ , dove devo prendere  $x_{n+1}$ ?



Devo escludere  $I(x_m, \frac{1}{2a})$  ← intervallo di centro  $x_m$  e raggio  $\frac{1}{2a}$   
 $I(x_{m-1}, \frac{1}{2a})$   
 $I(x_{m-2}, \frac{1}{3a})$   
 $\vdots$

Il caso peggiore è se sono DISGIUNTI

L'area della zona esclusa è  $\leq 2 \left( 1 + \frac{1}{2^a} + \dots + \frac{1}{n^a} \right)$

$\leq 2Ma$  ← non dipende da  $n$ !!!

Ricostruzione della diu:

- Fisso un intervallo di ampiezza  $> 2Ma$ .
- ad ogni passo resta sempre almeno un p.to buono nell'interv.

(Nota bene: gli intervalli esclusi passo-passo possono anche scappare da quello iniziale!)



### Criterio del confronto asintotico

$$\sum_{n=0}^{\infty} a_n$$

$$\sum_{n=0}^{\infty} b_n$$

Ipotesi:

- $a_n \geq 0$
- $b_n > 0$

- $\lim_{n \rightarrow +\infty} \frac{a_n}{b_n} = l \neq 0 \neq +\infty$

Allora le 2 serie hanno lo stesso comportamento  
(le somme possono essere diverse)

Dim. basta osservare che  $\frac{1}{2}l \leq \frac{a_n}{b_n} \leq 2l$  per  $n$  grande

$$\frac{1}{2}l b_n \leq a_n \leq 2l b_n$$

Esempio (serviva in RMM 2009)  $\sum_{n=1}^{\infty} \underbrace{\arctan \frac{1}{n}}_{a_n} = +\infty$

Prendo  $b_n = \frac{1}{n}$ :

$$\lim_{n \rightarrow +\infty} \frac{a_n}{b_n} = \lim_{n \rightarrow +\infty} \frac{\arctan \frac{1}{n}}{\frac{1}{n}} = \lim_{x \rightarrow 0} \frac{\arctan x}{x} = 1$$

$\uparrow$   
 $x = \frac{1}{n}$

Poiché  $\sum \frac{1}{n}$  diverge, anche l'altra diverge (per  $x \rightarrow 0$ )

Serie a segno alterno  $\sum_{n=0}^{\infty} (-1)^n d_n$

### CRITERIO DI LEIBNITZ

- Ipotesi:
- ①  $d_n \geq 0$
  - ②  $d_n \rightarrow 0$
  - ③  $d_{n+1} \leq d_n \quad \forall n \in \mathbb{N}$

Tesi: la serie converge

"Dim."  $S_{2m} \searrow \quad S_{2m+1} \nearrow \quad S_{2m} \geq S_{2m+1} \quad \forall m \quad \forall n$





Esempio  $1 - \frac{1}{2} + \frac{1}{3} - \frac{1}{4} + \dots = \sum (-1)^{n+1} \frac{1}{n}$  converge!

(e si sa anche la somma:  $\log 2$ )

$$\sum (-1)^n \frac{1}{n^a} \text{ converge } (\Leftrightarrow) \quad a > 0$$

— 0 — 0 — 0 —

$$1 - \frac{1}{2} - \frac{1}{4} + \frac{1}{3} - \frac{1}{6} - \frac{1}{8} + \frac{1}{5} \quad \text{Converge?}$$

Achtung!!!!  $\sum a_n$  converge +  $\sum |a_n|$  converge

⇓

posso riordinare i termini!!!!

Se  $\sum a_n$  converge, ma  $\sum |a_n| = +\infty$ , allora riordinando posso ottenere qualunque somma!

— 0 — 0 —

$$\sum_{n=1}^{\infty} \frac{1}{n^2} = \frac{\pi^2}{6} \quad \text{"Dimostrazione"} \quad \sin x = x - \frac{1}{6}x^3 + \frac{1}{5!}x^5 - \frac{1}{7!}x^7 + \dots$$

$$f(x) = \frac{\sin x}{x} = 1 - \frac{1}{6}x^2 + \frac{1}{5!}x^4 + \dots$$

$$g(x) = \frac{\sin(\pi x)}{x} = \frac{1}{x} \left( \pi x - \frac{1}{6}\pi^3 x^3 + \dots \right) = \pi - \frac{1}{6}\pi^3 x^2 + \dots$$

$$g(x) = 0 \Leftrightarrow x \in \mathbb{Z} \setminus \{0\}$$

In un polinomio, la somma dei reciproci delle radici è  $-\frac{a_1}{a_0}$

$$P(x) = a_0 + a_1 x + a_2 x^2 + a_3 x^3 + \dots \quad r_1, \dots, r_m \text{ radici}$$

$$\frac{1}{r_1^2} + \frac{1}{r_2^2} + \dots + \frac{1}{r_m^2} = \text{Somma}^2 - \text{doppi prodotti}$$

$$= \left(\frac{a_1}{a_0}\right)^2 - 2 \frac{a_2}{a_0}$$

Se fosse vero per le serie di potenze ...  $\sum_{n \in \mathbb{Z} \setminus \{0\}} \frac{1}{n^2} = -2 \frac{a_2}{a_0} = \frac{\pi^2}{3}$

↑  
contano 2 volte...

Max e min  $[a, b] \subseteq \mathbb{R}$  intervallo

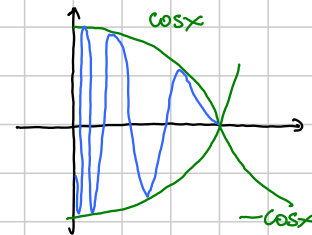
Teorema di WEIERSTRASS  $f: [a, b] \rightarrow \mathbb{R}$  continua, allora  
 esistono  $\max \{ f(x) : x \in [a, b] \}$   
 $\min \{ \quad \quad \quad \}$

*estremi compresi*

Esercizio  $\exists f: (0, 1]$  continua che non ammette né max, né min?

Esempio 1:  $f(x) = \frac{1}{x} \sin \frac{1}{x}$  e se la volessi pure limitata...

Esempio 2:  $f(x) = \cos x \cdot \sin \frac{1}{x}$



Dove trovare max/min

3 categorie di candidati:

- ① p.ti  $x \in (a, b)$ :  $f'(x) = 0$  (stazionari interni)
- ② " "  $f'$  non esiste (singolari interni)
- ③ Bordo  $x = a$  e  $x = b$ .

Esempio  $[-2, 2]$   $f(x) = |x|$   $\min = 0$  p.to di min  $x = 0$  ②  
 $\max = 2$  p.ti di max  $x = \pm 2$  ③

In 2 o più variabili  $A \subseteq \mathbb{R}^n$   $f: A \rightarrow \mathbb{R}$

- A si dice limitato se è contenuto in una opportuna palla

$$B(x_0, r) = \{ x \in \mathbb{R}^n : \text{dist}(x, x_0) \leq r \}$$

- A si dice chiuso se (brutalmente: contiene il bordo)  
 rigorosamente  $\forall x \notin A \exists r > 0$  t.c.  $B(x, r) \cap A = \emptyset$



- A si dice compatto se è chiuso + limitato

Teo. Weierstrass  $f: A \rightarrow \mathbb{R}$  con  $f$  continua e  $A$  compatto  
 $\Rightarrow$  esistono max e min.

Dove trovare p.ti di max/min

- ① P.ti  $x \in$  interno di  $A$  (p.ti  $x \in A$  per cui  $\exists B(x, \nu) \subseteq A$ )  
 con derivate parziali nulle

$$\left. \begin{cases} f_x = 0 \\ f_y = 0 \end{cases} \right\} \text{ in 2 variabili}$$

$$\nabla f = (f_x, f_y) = (0, 0)$$

↑  
gradiente  
di  $f$

- ② P.ti  $x \in$  interno di  $A$  in cui "qualche deriv. pars. non esiste".

- ③ Bordo (in generale sono  $\infty$  p.ti)

### MOLTIPLICATORI DI LAGRANGE

Caso di 1 moltiplicatore

$$V = \{x \in \mathbb{R}^n : \underbrace{\Phi(x) = 0}_{\text{equazione di } V}\}$$

Dato  $f(x)$ , trovare max/min di  $f$  in  $V$

Teorema Supponiamo  $x_0$  p.to di max/min per  $f$  in  $V$  (se  $V$  fosse compatto, esisterebbero tali punti).  
 Allora ci sono 3 possibilità

- ① In  $x_0$  non esiste  $\nabla f$  o non esiste  $\nabla \Phi$
- ② In  $x_0$   $\nabla \Phi(x_0) = 0$  (cioè si annullano tutte le deriv. parziali)
- ③ In  $x_0$   $\nabla \Phi(x_0)$  e  $\nabla f(x_0)$  sono paralleli, cioè  $\exists \lambda \in \mathbb{R}$  b.c.

$$\nabla f(x_0) = \lambda \nabla \Phi(x_0)$$

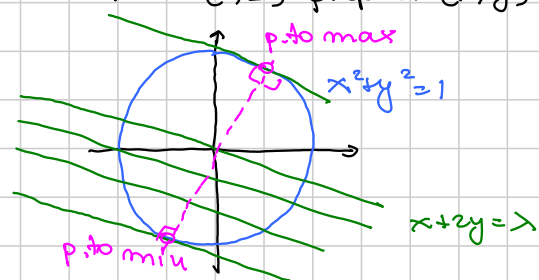
↑  
moltiplicatore

Esempio 1 max/min  $\{x+2y : x^2+y^2=1\}$

Per via elementare: AM-QM pesata o C-S:

$$x+2y \leq \sqrt{5} \sqrt{x^2+y^2} \quad \text{con} = (\Leftrightarrow) \quad (1,2) \text{ prop a } (x,y)$$

Via linee di livello:



Via moltiplicatori:

- Max e min esistono per  $W$ .

- Vedo 3 possibili situazioni  $\Phi(x,y) = x^2 + y^2 - 1$   $f(x,y) = x + 2y$

① Nulla  $\emptyset$

$$\textcircled{2} \begin{cases} 2x = 0 \\ 2y = 0 \\ x^2 + y^2 = 1 \end{cases} \quad \nabla \Phi(x,y) = (2x, 2y)$$

Nulla  $\emptyset$  (di solito è così!)

- ③ Devo cercare i p.ti  $(x,y) \in V$  t.c.  $\nabla f(x,y) = \lambda \nabla \Phi(x,y)$   
 $(1, 2) = \lambda(2x, 2y)$

$$\begin{cases} 1 = 2\lambda x \\ 2 = 2\lambda y \\ x^2 + y^2 = 1 \end{cases} \quad \text{3 equ in 3 incognite}$$

Dalla 1<sup>a</sup> e 2<sup>a</sup> ottengo  $y = 2x$   
 sostituisco nella 3<sup>a</sup> e fine!  
 — o — o —

Esempio 2 AM-GM  $\sqrt[n]{x_1 \cdot \dots \cdot x_n} \leq \frac{x_1 + \dots + x_n}{n}$

Osservo che è omogenea, quindi posso assumere LHS = 1 oppure RHS = 1

Scelgo LHS = 1 Vuol dire che

$$V = \{(x_1, \dots, x_n) \in \mathbb{R}^n : x_1 \cdot \dots \cdot x_n = 1, x_1 \geq 0, \dots, x_n \geq 0\}$$

È compatto? NO!  $(\frac{1}{m}, m, 1, \dots, 1) \in V \quad \forall m$  NO LIMITATO

Scelgo RHS = 1  $V = \{( ) \in \mathbb{R}^n : x_1 + \dots + x_n = 1, x_1 \geq 0, \dots, x_n \geq 0\}$

Ora  $V \subseteq [0, 1]^n \Rightarrow$  limitato

Ora so che esiste Max del LHS (per  $W$ .)

cerco i p.ti di max in ①, ②, ③

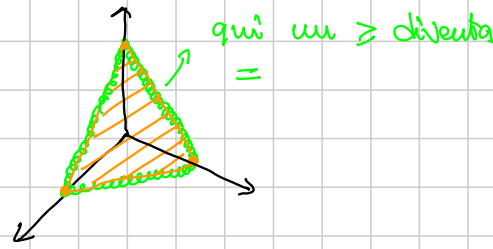
$$\Phi(x_1, \dots, x_n) = x_1 + \dots + x_n \quad f(x_1, \dots, x_n) = x_1 \cdot \dots \cdot x_n$$

①  $\emptyset$  ②  $\emptyset$   $\nabla \Phi = (1, \dots, 1)$  ③  $\nabla f = \lambda \nabla \Phi$

$$\begin{cases} x_2 \cdot \dots \cdot x_n = \lambda \\ x_1 \cdot x_3 \cdot \dots \cdot x_n = \lambda \\ \vdots \\ x_1 + \dots + x_n = 1 \end{cases} \quad \left. \begin{array}{l} \text{Tutte le variabili sono uguali} \\ \text{(se posso dividere)} \end{array} \right\}$$

$\rightarrow x_1 = \dots = x_n = \frac{1}{n}$  e si conclude

Achtung! Se  $V$  è descritto da 1 equazione + un po' di disuguaglianze  
 compattiamo i "bordi dei bordi"  
 $m=3$



Il metodo va bene per trovare i candidati nell' "interno" di  $V$ , cioè dove le disug. sono strette.

- Cosa dire
- Max e min esistono per  $W$ .
  - Se p.to di max sta nel "bordo del bordo" allora LHS=0
  - Se non sta lì, posso dividere e finisco

Esempio 3  $\min \{x^2+y^4+z^6 : xyz = m\} \quad x, y, z \geq 0$   
 non compatto

Mettendo  $x=m, y=1, z=1$  ottengo  $\min \leq m^2+2$

Quindi min originario =  $\min \{x^2+y^4+z^6 : xyz = m, x \leq m, y \leq m, z \leq m\}$   
 compatto !!!

Quindi il minimo esiste, e non è su un bordo del bordo (perché altrimenti cambio  $m$  con  $2m$ ...)

I p.ti di minimo stanno nel gruppo ③  $\nabla f = \lambda \nabla \phi$   
 $(2x, 4y^3, 6z^5) = \lambda (yz, xz, xy)$

$$\begin{cases} 2x = \lambda yz & \cdot x \\ 4y^3 = \lambda xz & \cdot y \\ 6z^5 = \lambda xy & \cdot z \\ xyz = m \end{cases} \quad 2x^2 = 4y^4 = 6z^6 \quad \text{e da qui si chiude}$$

basta sostituire nella 4ª:

$$x^2 + y^4 + z^6 \geq C [xyz]^{\frac{12}{11}}$$

Achtung! Alla fine i punti vanno sostituiti per trovare max/min

$$f(x,y) = x^2 - y^2$$

## SENIOR 2010 - ADVANCED - A

Titolo nota

08/09/2010

### Più moltiplicatori di Lagrange

$$V = \{ x \in \mathbb{R}^n : \phi_1(x) = \dots = \phi_k(x) = 0 \} \quad k \text{ equazioni } (k < n)$$

$$f: \mathbb{R}^n \rightarrow \mathbb{R} \quad \max/\min \{ f(x) : x \in V \}$$

Teorema Se  $x_0 \in V$  è p.to di max/min, allora succede una di queste 3 cose

- $\nabla f$  o qualche  $\nabla \phi_i$  non esiste in  $x_0$
- Costruiamo la matrice che ha come righe i  $k$  gradienti

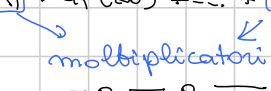
$$\begin{array}{c} \nabla \phi_1 \\ \nabla \phi_2 \\ \vdots \\ \nabla \phi_k \end{array} \quad \text{Matrice } k \times n \quad \begin{array}{cc} \odot & \odot \\ \odot & \odot \end{array}$$

[ Rangho di una matrice = dimensione della massima sottomatrice quadrata  $r \times r$  con  $\det \neq 0$   
 = max numero di righe lin. indep.  
 = " " " colonne " "  
 = dimensione dell'immagine ]

In  $x_0$  il rango è  $< k$  (cioè tutte le sottomatrici  $k \times k$  hanno determinante = 0 ( $k$ ) condizioni)

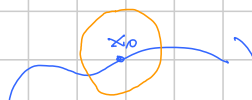
- In  $x_0$  si ha che  $\nabla f$  è combinazione lineare dei  $\nabla \phi_i$ , cioè  
 $\exists \lambda_1, \dots, \lambda_k$  b.c.

$$\nabla f(x_0) = \lambda_1 \nabla \phi_1(x_0) + \dots + \lambda_k \nabla \phi_k(x_0)$$



Idea per una dim. con 1 moltiplicatore:

$$V = \{ x \in \mathbb{R}^n : \phi(x) = 0 \} \quad x_0 \in V \text{ p.to di min per } f(x) \text{ ristretta a } V$$

 Pseudo  $B(x_0, \pm)$  chiusa

Per ogni  $n$  calcolando  $\min \{ f(x) + n \phi^2(x) : x \in B(x_0, \pm) \}$

Esiste ed è raggiunto in un p.to  $x_n$ . Si può dimostrare che  $x_n \rightarrow x_0$ ,

In  $x_n$  si ha che  $\nabla (f(x) + n \phi^2(x)) = 0$ , cioè  $\nabla f(x) + 2n \phi(x) \nabla \phi(x) = 0$

cioè in  $x_n$   $\nabla f(x_n)$  è multiplo di  $\nabla \phi(x_n)$ . Passo al limite

Lemma sui vettori  $V_m \rightarrow V_\infty$   $\forall m$  si ha che  $V_m$  è multiplo  
 $W_m \rightarrow W_\infty$  di  $W_m$

Allora  $\rightarrow 0$   $W_\infty = 0$   
 $0$   $V_\infty$  è multiplo di  $W_\infty$   
 $\rightarrow 0 \rightarrow 0 \rightarrow$

Con  $k$  equazioni la dim. è la stessa

$$\min \{ f(x) + m \phi_1^2(x) + \dots + m \phi_k^2(x) \}$$

Nel p.to di min  $x_m$  si ha che

$$\nabla f(x_m) + 2m \phi_1(x_m) \nabla \phi_1(x_m) + \dots + 2m \phi_k(x_m) \nabla \phi_k(x_m) = 0$$

cioè  $\nabla f(x_m)$  è comb. lineare dei  $\nabla \phi_i(x_m)$

Lemma Dato  $k+1$  succ. di vettori  $V_m \rightarrow V_\infty$   
 $W_m^1 \rightarrow W_\infty^1$   
 $W_m^k \rightarrow W_\infty^k$

Hip: per ogni  $m$  si ha che  $V_m$  è comb. lin. dei  $W_m^i$

Test:  $\rightarrow 0$  il limite è ancora comb. lin.

$\rightarrow 0$  i  $W_\infty^i$  non sono comb. lin. indep., cioè la matrice da essi formata non ha rango max.

Dim. Ipotesi dice che la matrice  $\begin{pmatrix} V_m \\ W_m^1 \\ \vdots \\ W_m^k \end{pmatrix} (k+1) \times n$

ha rango  $\leq k$ , cioè tutte le s. matrici  $(k+1) \times (k+1)$  hanno  $\det = 0$

Quindi anche la matrice limite ha rango  $\leq k$ , cioè

$\rightarrow 0$  i  $k$  sotto sono dip. tra di loro

$\rightarrow 0$  il primo dipende dagli altri.  
 $\rightarrow 0 \rightarrow 0 \rightarrow$

A meno di cambiare  $f(x)$  con  $f(x) + \text{dist}(x, x_0)$  possiamo supporre  $x_0$  p.to di minimo stretto (basta il  $\square$ ).  
 $\rightarrow 0 \rightarrow 0 \rightarrow$

Problemi ad applicare i moltiplicatori

$\rightarrow$  sapere che max/min esistono (guadagnare compattezza)

$\rightarrow$  bordi dei bordi

$\rightarrow$  risolvere il sistema!  
 $\rightarrow 0 \rightarrow 0 \rightarrow$

## CONVESSITÀ

Def. Sia  $A \subseteq \mathbb{R}^m$  un insieme convesso (per ogni coppia di p.ti contiene tutto il segmento)

$f: A \rightarrow \mathbb{R}$  si dice convessa se

$$f(\lambda x + (1-\lambda)y) \leq \lambda f(x) + (1-\lambda)f(y) \quad \forall x \in A \quad \forall y \in A \\ \forall \lambda \in [0,1]$$

Oss. In una variabile si può verificare con  $f''(x) \geq 0$   
 In più variabili ci sarebbe la matrice Hessiana, oppure sfruttando che è somma di funzioni convesse semplici (distanza da un p.to, distanza<sup>2</sup> da un p.to, distanza da una retta, ...)

Oss.  $f$  è convessa  $\Leftrightarrow$  il sopragrafiteo  $\{(x,y) \in A \times \mathbb{R} : y \geq f(x)\}$   
 è un insieme convesso di  $\mathbb{R}^{m+1}$   
 — o — o —

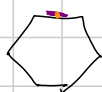
Teorema Se  $A \subseteq \mathbb{R}^m$  è un convesso compatto, allora  
 $\max \{f(x) : x \in A\}$ ,  
 se  $f$  è convessa e continua, esiste e tra i p.ti di max c'è almeno un p.to estremo di  $A$  (solo estremi se  $f$  è strettamente convessa)

Def.  $x_0 \in A$  è p.to estremo  $\Leftrightarrow$  quando  $x_0 \in$  segmento  $\subseteq A$  si ha che  $x_0$  è un estremo del segmento

Esempi



p.ti estremi = bordo



p.ti estremi = vertici

— o — o —

Esempio Trovare in un triangolo il p.to per cui è massima la somma delle distanze dai vertici

Soluzione La somma delle distanze è convessa (anzi strettamente convessa perché i 3 vertici non sono allineati ...), quindi il max si realizza in un p.to estremo, cioè in un vertice (quello opposto al lato minore)

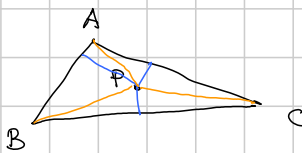


Come si dim. per bene che è strett. convessa.

$$f(x) = \text{dist}(x, A) + \text{dist}(x, B) + \text{dist}(x, C).$$

La disug. " $\leq$ " nella definizione è facile. Se ci fosse " $=$ " i p.ti  $x$  e  $y$  dovrebbero essere allineati con  $A, B, C$  (il che è impossibile).  
Infatti dovrebbe esserci " $=$ " sui 3 addendi.

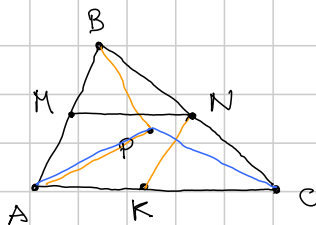
Esempio 2 La somma delle distanze dai vertici di un  $\Delta$  è convessa, anzi è affine nel triangolo (il grafico è un piano, perché somma di 3 piani).  
Quindi il max è almeno in un vertice, quindi è l'altezza + lunga



Esempio 3 IMOSL 99 Dim. che in un triangolo ABC se prendo un p.to P si ha che

$$\underbrace{AP + BP + CP + \min\{AP, BP, CP\}}_{LHS} \leq AB + BC + CA$$

Dim.



Facile:  $AP + CP \leq AM + MN + NC$   
 $AP + BP \leq BN + NK + AK$

Somma:

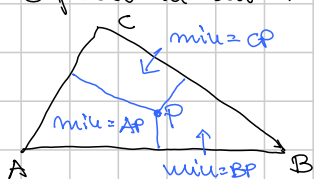
$$LHS \leq 2AP + BP + CP \leq \frac{1}{2} AB + \frac{1}{2} AC + \frac{1}{2} BC + \frac{1}{2} BC + \frac{1}{2} AB + \frac{1}{2} AC = \text{Perim.}$$

Osservazione: è sempre vero che P è dentro almeno 2 "trapezi"

Dim. "bovina"

Pongo  $f(P) = AP + BP + CP + \min\{AP, BP, CP\}$

È convessa? FORGET IT!! Su tutto il triangolo NO, ma lo è nei 3 pezzi in cui il minimo so chi è.



Nelle 3 zone  $f(P)$  è convessa (strett.) perché somma di convesse

$$2AP + BP + CP$$

Basta quindi controllare i punti estremali delle 3 zone, cioè

→ vertici

→ punti medi dei lati

→ circocentro

Da qui si finisce.

— o — o —

Equivalente in  $n$  variabili di  $f''(x)$ . Derivate parziali seconde! (sono  $n^2$ )

$$\begin{pmatrix} f_{xx} & f_{xy} & f_{xz} \\ f_{yx} & f_{yy} & f_{yz} \\ f_{zx} & f_{zy} & f_{zz} \end{pmatrix}$$

La matrice è simmetrica

Def. Una Matrice simmetrica  $n \times n$  si dice semidefinita positiva se

$$x \cdot Ax \geq 0 \quad \forall x \in \mathbb{R}^n$$

Def. Forma quadratica in  $\mathbb{R}^n =$  somma di monomi di 2° grado in  $n$  variabili

Esempio ( $n=2$ )

$$q(x,y) = ax^2 + by^2 + 2cxy$$

( $n=3$ )

$$q(x,y,z) = ax^2 + by^2 + cz^2 + 2dxy + 2eyz + 2fzx$$

Una forma è semidefinita pos. se viene  $\geq 0$  per ogni  $x \in \mathbb{R}^n$

Forme  $\leftrightarrow$  matrici simmetriche

$$ax^2 + by^2 + 2cxy \quad \begin{pmatrix} x & y \end{pmatrix} \begin{pmatrix} a & c \\ c & b \end{pmatrix} \begin{pmatrix} x \\ y \end{pmatrix}$$

Teorema  $f$  è convessa in  $A \iff$  la matrice delle derivate seconde è semidefinita positiva


se ci sono le derivate

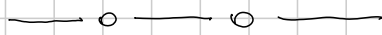
Come stabilire se una forma è positiva, negativa o "indefinita"

→ autovalori della matrice

→ ogni forma è somma / differenza di quadrati

$$x^2 + 3xy - y^2 = x^2 + 2x \cdot \frac{3y}{2} + \frac{9}{4}y^2 - \frac{9}{4}y^2 - y^2 = \left(x + \frac{3y}{2}\right)^2 - \frac{13}{4}y^2$$

→ condizione sufficiente (ma non necessaria) per essere positiva  tutte del > 0



DISUG. DI JENSEN  $I \subseteq \mathbb{R}$  convessa  $f: I \rightarrow \mathbb{R}$  convessa

$$f(\lambda_1 x_1 + \dots + \lambda_n x_n) \leq \lambda_1 f(x_1) + \dots + \lambda_n f(x_n) \quad \begin{matrix} x_1, \dots, x_n \in I \\ \lambda_1 + \dots + \lambda_n = 1 \\ \lambda_i \geq 0 \end{matrix}$$

Dim. Induzione UP and DOWN

1) Vera per  $n \Rightarrow$  vera per  $2n$  ( $n=2$ : def. di funzione convessa)

2) Vera per  $n \Rightarrow$  vera per  $(n-1)$ : basta usare un  $\lambda$  nullo

Dim. 1)  $f(\lambda_1 x_1 + \dots + \lambda_n x_n + \lambda_{n+1} x_{n+1} + \dots + \lambda_{2n} x_{2n}) =$

$$f\left( (\lambda_1 + \dots + \lambda_n) \frac{\lambda_1 x_1 + \dots + \lambda_n x_n}{\lambda_1 + \dots + \lambda_n} + (\lambda_{n+1} + \dots + \lambda_{2n}) \frac{\lambda_{n+1} x_{n+1} + \dots + \lambda_{2n} x_{2n}}{\lambda_{n+1} + \dots + \lambda_{2n}} \right)$$

$$\leq (\lambda_1 + \dots + \lambda_n) f\left( \frac{\lambda_1 x_1 + \dots + \lambda_n x_n}{\lambda_1 + \dots + \lambda_n} \right) + (\lambda_{n+1} + \dots + \lambda_{2n}) f\left( \frac{\lambda_{n+1} x_{n+1} + \dots + \lambda_{2n} x_{2n}}{\lambda_{n+1} + \dots + \lambda_{2n}} \right)$$

Ora uso Jensen su  $n$  vero per ipotesi induttiva (i coefficienti sono

$$\frac{\lambda_1}{\lambda_1 + \dots + \lambda_n}, \dots, \frac{\lambda_n}{\lambda_1 + \dots + \lambda_n} \text{ che fanno somma 1}$$

Occiso a trattare il caso della divisibilità per zero ...

Oss. Caso particolare  $f\left(\frac{x_1 + \dots + x_n}{n}\right) \leq \frac{f(x_1) + \dots + f(x_n)}{n}$

Volendola dim. direttamente faccio 1) come prima, ma per 2)

è più complicato:

solo dati  $x_1, \dots, x_{n-1}$  e devo scegliere  $x_n = \frac{x_1 + \dots + x_{n-1}}{n-1}$ ,

Mi ritorna

$$\underbrace{f\left(\frac{x_1 + \dots + x_{n-1}}{n-1}\right)}_{\text{LHS}} = f\left(\frac{x_1}{n} + \dots + \frac{x_{n-1}}{n} + \frac{x_n}{n}\right)$$

$$x_1 \left(\frac{1}{n} + \frac{1}{n(n-1)}\right) = x_1 \frac{n}{n(n-1)}$$

$$\leq \frac{1}{n} f(x_1) + \dots + \frac{1}{n} f(x_{n-1}) + \frac{1}{n} f(x_n) \quad \underbrace{\hspace{10em}}_{\frac{1}{n} \text{ LHS}}$$

Porto a sx e viene

$$\text{LHS } \frac{m-1}{m} \leq \text{RHS} \quad \text{--- } 0 \text{ ---}$$

Applicazione 1 AM-QM (Jensen su  $f(x) = x^2$ )  
(N.B. vale senza Hp di segno)

Applicazione 2 GM-AM (Jensen sulla concava  $f(x) = \log x$ ;  
qui servono  $x_i > 0$ )

$$\log\left(\frac{x_1 + \dots + x_m}{m}\right) \stackrel{\text{concava}}{\geq} \frac{1}{m} (\log x_1 + \dots + \log x_m) \\ = \log [x_1 \dots x_m]^{\frac{1}{m}}$$

Usando i  $\lambda$  si ottengono versioni pesate.

Applicazione 3  $M_p(x_1, \dots, x_m) \leq M_q(x_1, \dots, x_m)$   $p \leq q$   
ci si riduce facilmente a  $0 < p < q$  e  $p=1$   $q=\alpha > 1$   
Ora è Jensen su  $x^\alpha$  che è convessa per  $\alpha \geq 1$  ( $x \geq 0$ )

Applicazione 4  $ab \leq \frac{1}{p} a^p + \frac{1}{q} b^q$  se  $a > 0, b > 0, p > 0, q > 0$   
 $\frac{1}{p} + \frac{1}{q} = 1$

Jensen su  $\log x = f(x)$   $x = a^p$   $y = b^q$   $\lambda = \frac{1}{p}$   $(1-\lambda) = \frac{1}{q}$

$$\log(\lambda x + (1-\lambda)y) \geq \lambda f(x) + (1-\lambda)f(y)$$

$$\log\left(\frac{1}{p} a^p + \frac{1}{q} b^q\right) \geq \frac{1}{p} \log(a^p) + \frac{1}{q} \log(b^q)$$

$$= \log a + \log b = \log(ab)$$

Analogo  $abc \leq \frac{1}{p} a^p + \frac{1}{q} b^q + \frac{1}{r} c^r$   $\frac{1}{p} + \frac{1}{q} + \frac{1}{r} = 1$

Senza per dimostrare C-S a 3 specie (Hölder)

$$\sum a_i b_i c_i \leq \left(\sum a_i^p\right)^{\frac{1}{p}} \left(\sum b_i^q\right)^{\frac{1}{q}} \left(\sum c_i^r\right)^{\frac{1}{r}} \quad \frac{1}{p} + \frac{1}{q} + \frac{1}{r} = 1$$

Per omogeneità posso assumere

$$\sum a_i^p = \sum b_i^q = \sum c_i^r = 1 \quad \text{e poi sommo sui } r \text{ le YOUNG.}$$

Caso speciale :  $i \in \{1, 2\}$   $p_1 = p_2 = \dots = p_n = m$

$$(a_1, \dots, a_m) + (b_1, \dots, b_m) \leq (a_1^m + b_1^m)^{\frac{1}{m}} \cdot \dots \cdot (a_m^m + b_m^m)^{\frac{1}{m}}$$

$$= \sqrt[m]{(a_1^m + b_1^m) \cdot \dots \cdot (a_m^m + b_m^m)}$$

$(a_1, b_1)$   
 $(a_2, b_2)$   
 $\vdots$   
 $(a_m, b_m)$

Dimostrazione diretta : UP and DOWN

Applicazione 3  $\min \{ f(x_1) + \dots + f(x_n) ; x_1 + \dots + x_n = A \}$

Se  $f$  è convessa, allora il minimo è quando sono tutti uguali

$$f(x_1) + \dots + f(x_n) \geq n f\left(\frac{x_1 + \dots + x_n}{n}\right) = n f\left(\frac{A}{n}\right)$$

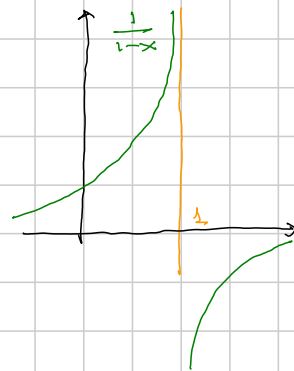
Nesbit  $\frac{a}{b+c} + \frac{b}{c+a} + \frac{c}{a+b} \geq \frac{3}{2}$   $a, b, c > 0$

Per omogeneità assumo  $a+b+c = 1$   $\frac{a}{1-a} + \frac{b}{1-b} + \frac{c}{1-c} \geq \frac{\frac{1}{3}}{1-\frac{1}{3}} \cdot 3$

Se  $f(x) = \frac{1}{1-x}$  è convessa ho finito

Rigoroso :  $f''(x) = \dots$  oppure

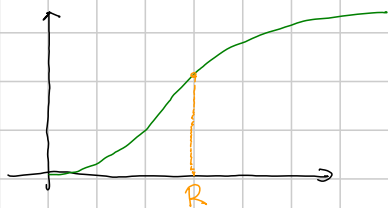
cambio variabile e diventa  $\frac{1}{x}$  per la quale uso la definizione.



Achtung! Precipare sempre DOVE si lavora.  
 In questo caso :  $[0, 1)$

**LEMMA CONVEX-CONCAVE**  $\min \{ f(x_1) + \dots + f(x_n) ; x_1 + \dots + x_n = A, x_i \geq 0 \}$

Supponiamo  $f$  convex-concave. Per il minimo se la giocano 2 configurazioni



- $x_1 = \dots = x_m$  nella zona convessa
- $x_1 = \dots = x_{n-1}$  nella zona convessa e  $x_n$  nella zona concava per aggiustare la somma

Ponendo  $x = x_n$  mi riduco a studiare

$$f(x) + (n-1) f\left(\frac{A-x}{n-1}\right)$$

oppure, ponendo  $x_1 = x$  :  $(n-1) f(x) + f(A - (n-1)x)$   
e si può fare studiando la funzione.

Idee alla base del convex-concave:

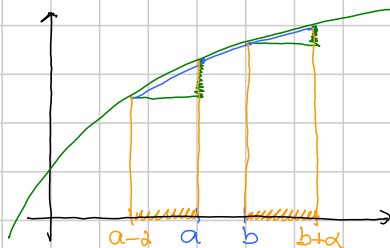
- 1) min esiste se  $f$  è continua in  $[0, A]$  : Weierstrass.
- 2) tutti gli  $x_i$  che stanno nella zona convessa si possono rendere uguali conservando la somma ed abbassando  $\sum f(x_i)$
- 3) A max ci c'è uno nella zona concava : se ce ne sono 2 di "allargare" conservando la somma fino a quando il  $x_i$  cade nella zona convessa (il dx non può arrivare ad A senza che tutti gli altri siano zero)

Lemma di allungamento (versione concava)

$$f(a-d) + f(b+d) \leq f(a) + f(b)$$

il che equivale a

$$\underbrace{f(b+d) - f(b)}_{\alpha f'(c_1)} \leq \underbrace{f(a) - f(a-d)}_{\alpha f'(c_2)}$$



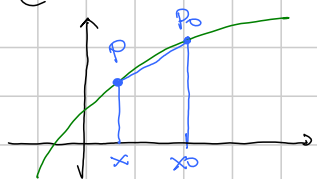
Più in generale : se  $y > x$  e  $\alpha > 0$ , allora

$$f(y+\alpha) - f(y) \leq f(x+\alpha) - f(x)$$



Non ho garanzia di ordinamento tra  $c_1$  e  $c_2$ .

Idea: fissato un punto  $x$ , il rapporto incrementale è monotono (crescente o decrescente a seconda di convessità/concavità)



il coeff. angolare della retta  $PP_0$   
diminuisce quando  $P$  sale (fissato  $P_0$ )  
(si dimostra a partire dalla def. di convessità)

**DISUGUAGLIANZA DI KARAMATA**  $f: I \rightarrow \mathbb{R}$  convessa

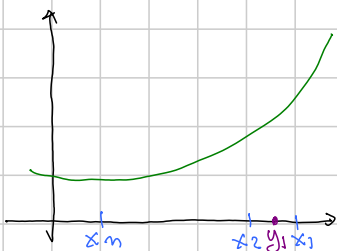
$(x_1, \dots, x_n) \succ (y_1, \dots, y_n)$  in stile bunching, cioè

- $x_1 \geq y_1$
- $x_1 + \dots + x_k \geq y_1 + \dots + y_k \quad k=1, \dots, n-1$
- $x_1 + \dots + x_n = y_1 + \dots + y_n$  (controllare!!!!)
- $x_1 \geq x_2 \geq \dots \geq x_n, \quad y_1 \geq y_2 \geq \dots \geq y_n$

Allora:  $f(x_1) + \dots + f(x_n) \geq f(y_1) + \dots + f(y_n)$

KARAMATA  $\Rightarrow$  JENSEN      Ordiniamo gli  $x_i$  e poi uso  $y_1 = \dots = y_n = \frac{x_1 + \dots + x_n}{n}$

Idea della dim: induzione + Lemma di allargamento



Prendo  $x_1$  e lo sposto verso  $y_1$  e contemporaneamente alzo  $x_n$  finché non ho "contatti" tra  $x$  e  $y$ .  
 Dopo il contatto semplifico e scendo di uno. Così però mangerei la somma  $x_1 + x_2 + x_3$  è scesa ...  
 Forse è meglio lavorare con  $x_1$  e  $x_2$  ...  
 Alternativa: spostare  $x_1$  e  $y_1$  a sx  
 Fino a quando ...

NB. Le funzioni convesse stanno sopra le proprie rette tangenti !!

Idea: derivata = rapporto incrementale con  $P=\mathbb{R}$

Applicazioni (Bernoulli con esponente  $\alpha \geq 1$ )

$$(1+x)^\alpha \geq 1+\alpha x \quad \forall x \geq -1 \quad \forall \alpha \geq 1$$

$f(x) = (1+x)^\alpha$  è convessa per  $\alpha \geq 1$        $f(x) \geq 1+\alpha x$   
 $f(0) = 1$      $f'(x) = \alpha(1+x)^{\alpha-1}$      $f'(0) = \alpha$       *retta tg. in  $x=0$*

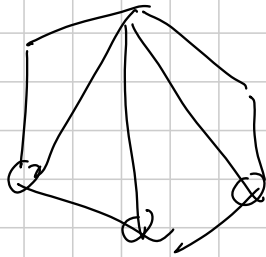
# COMBINATORICS

Titolo nota

06/09/2010

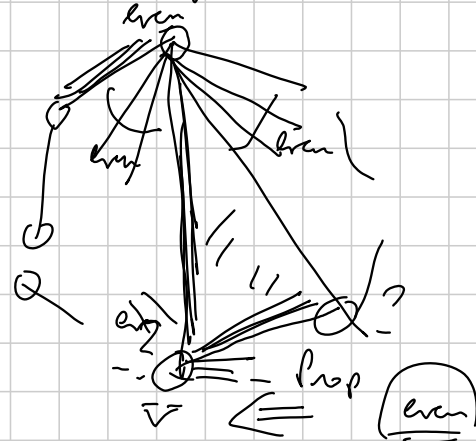
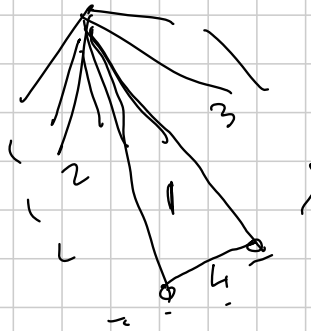
## 1. Induction.

1.1. Given a convex polygon<sup>n</sup>, triangulated by diagonals, such that the number of segments meeting at each vertex is even.



Question? When?

- $n = 3$  ✓
- $n = 4$  —
- $n = 5$  —
- $n = 6$  ✓



Claim  $v$  also has even # segments

Polygon + sides diagonals

$\Rightarrow$  Graph  $G = (V, E)$



$\deg v = \#$  edges incident at  $v$ .

$$\sum_{v \in V} \deg v = 2|E| \text{ (even)}$$

# (vertices of odd degree)

$$\sum_{\deg v = \text{odd}} \deg v = \underline{\underline{\text{even}}}$$



⇒ even : Proposition The number of vertices of odd degree in a finite graph is even.

Start again. Define the induction predicate.

Induction on number of vertices have to be ℕ.

Sufficient condition

— show how (exhibit.) Necessary condition

- 1) to just describe example
- 2) to inductively build it.

Find some "special value" — "threshold" value

Example -  $\{1, 2, \dots, 2n\}$

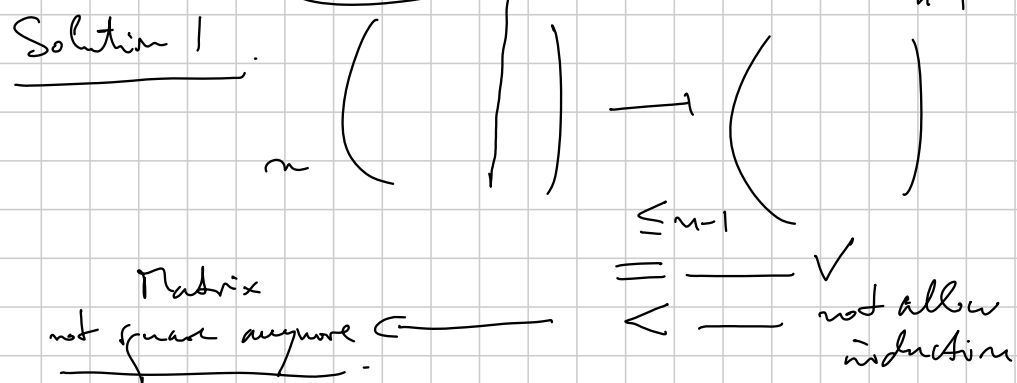
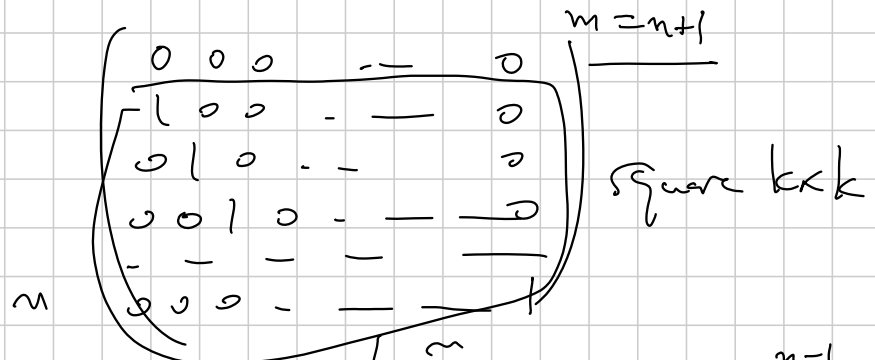
What is minimum cardinality of a subset so that there must exist 2 coprime numbers in that subset.

(D) more than  $|S| > t$  value

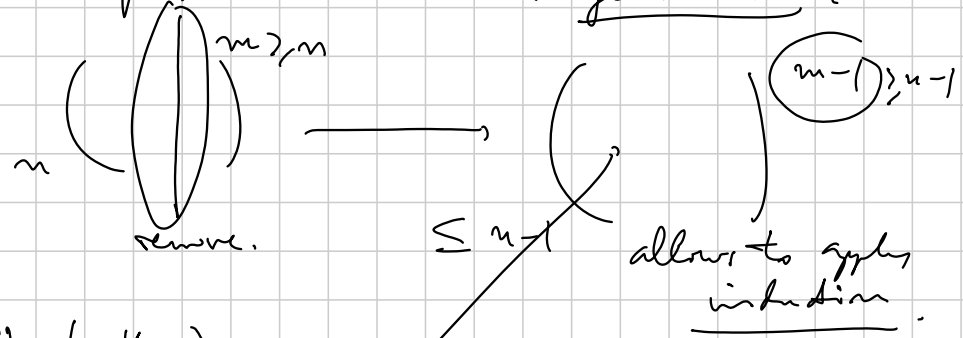
(E) ~~more than~~  $|S| = t - 1$  not true. example.

Problem. Given is  $n \times n$  matrix, entries are symbols  $a_{ij}$ . Each row  $i$  seen  $\binom{a_{i1}, a_{i2}, \dots, a_{in}}{\in \text{space of } n \text{ symbols}}$

Given if the fact  $R_i \neq R_j$  as vectors.  
 Show that we may choose to remove  
 some column, so that in the new  
 $n \times (n-1)$  matrix, the rows are still  
distinct (as vectors).



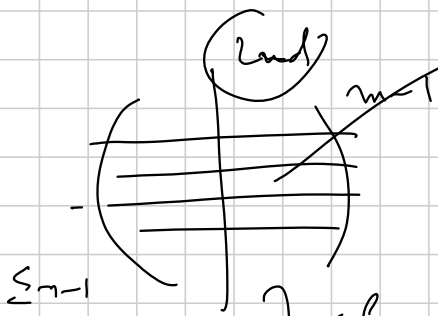
"Patch" to this slide: replace our problem with a more general one:



(Check the Ind. Hyp for small values  $n=1, 2$ )  
Ind. Hyp  $\rightarrow$  YES, I can remove a column s.t. remaining rows distinct.

$n=m=1$ , remove the single column  $\rightarrow$  original statement true?

by a vacuous argument  $\rightarrow$  YES



3 columns may be removed, rows to remain distinct.

Now  $\rightarrow$  we go back, remove column from  $n \times m$  matrix.  $\square$ .

Solution 2. (Story) Ross Hornberger.

Assume that any  $C_k$  column we remove, there appear duplicate rows  $R_i$  and  $R_j$ . For  $k \neq k \rightarrow$  could they be the same? NO.

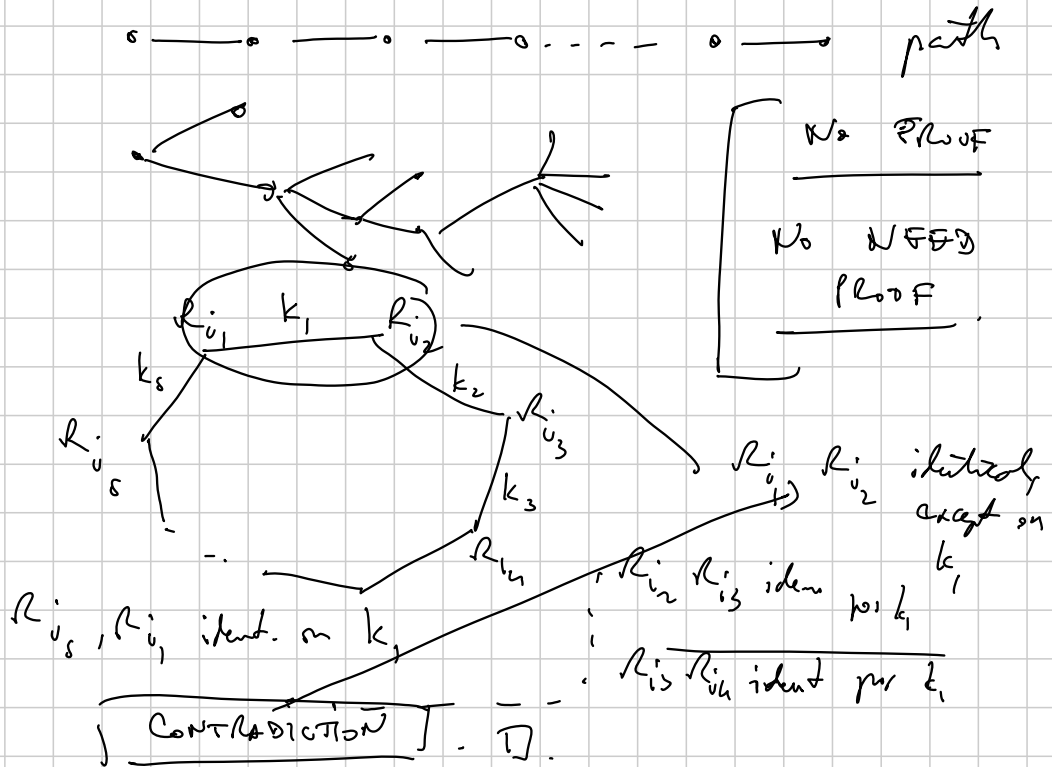
Build graph  $G$ : vertices  $R_i$   
edges between  $R_i, R_j$   
 $|G| = n, ||G|| = n$ .

What do we know about such graphs?

There must be a cycle! (YES)

Because the largest number of edges graph  $n$  vertices may have, is that no cycles are present  $\rightarrow$  if  $n-1$

$n-1$  the maximal # edges of an acyclic graph. Are called trees:



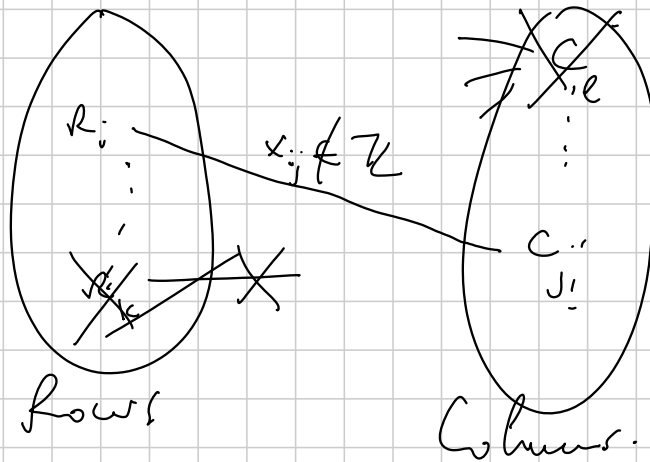
Generalisation. Given  $n \times m$  matrix.

$m > n \rightarrow$  we may remove  $\lfloor m-n+1 \rfloor$  columns  $\rightarrow$  rows to remain distinct.

Problem. We have some rectangular array  $n \times m$ , whose entries real numbers  $x_{ij}$ . We know that sums by rows and by columns are integer numbers.

Prove that we can replace (if need is)  $x_{ij}$  with  $y_{ij}$  integer:  $y_{ij} \in \{\lfloor x_{ij} \rfloor, \lceil x_{ij} \rceil\}$ , such that the sums by rows and columns stay the same.

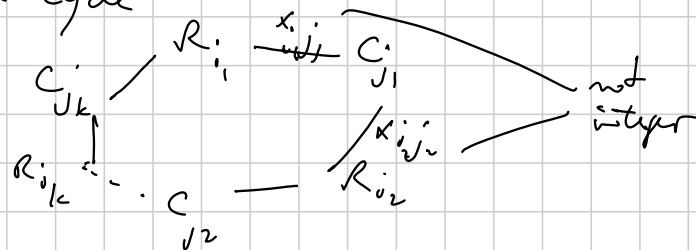
Solution. Those  $x_{ij} \in \mathbb{Z}$  do not matter.



Graph, where vertices partitioned 2 classes, and edges are only between classes.

Results about bipartite graphs is called Bipartite  
 $\Leftrightarrow$  Cycles are of even length.

Let's take a cycle



Let us denote  $S_j \in \mathbb{Z}$  the least distance from any of the  $x_{ij}$  to their neighborly integers.

$$x_{i+1, j+1} \xleftarrow{\varepsilon} x_{i, j} \xrightarrow{\varepsilon} x_{i+2, j+2}, \quad x_{i+2, j+2} \xleftarrow{\varepsilon} x_{i+1, j+1}, \quad \text{and so on.}$$

$$x_{i+1, j+1} \xrightarrow{\varepsilon} x_{i, j+1}$$

End  $\rightarrow$  some ( $\geq 1$ )  $x_{ij}$ 's  $\rightarrow$  become integers and all the conditions are obeyed.

Recurse induction  $\rightarrow$  scale algorithm (loop)  
 Build a new bipartite graph (strictly less edges).

Problem . Noga Alon (Combinatorial Nullstellensatz).

Consider  $\{1, 2, \dots, n\}$ . Consider two permutations  $\sigma, \tau \in S_n$ .

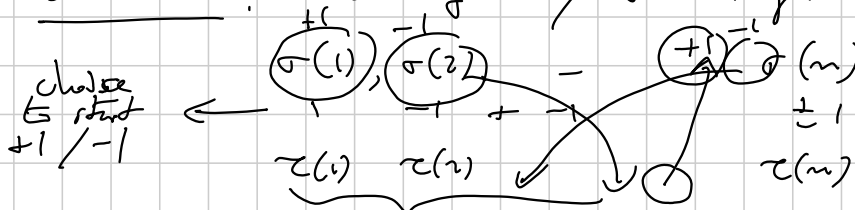
$$\begin{matrix} \sigma(1), \sigma(2), \dots, \sigma(n) \\ \tau(1), \tau(2), \dots, \tau(n) \end{matrix} \parallel \text{rows, } f \pm 1's.$$

We look at functions  $f: \{1, 2, \dots, n\} \rightarrow \{-1, +1\}$ .

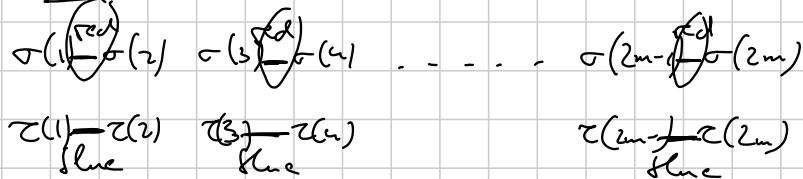
We look at  $\sum_{k=i}^{k=j} f(\sigma(k))$  (sums of consecutive terms in the  $\sigma$ -sequence)

Question is? Can we find such an  $f$ .  
 So that all these "partial sums", for  $\sigma$ , and  $\tau$ , they all are at most  $\underline{2}$  in absolute value?

Solution. What if only  $\sigma$  was given?



Subcase  $[n = 2m]$  (even)  $-1 \quad +1$



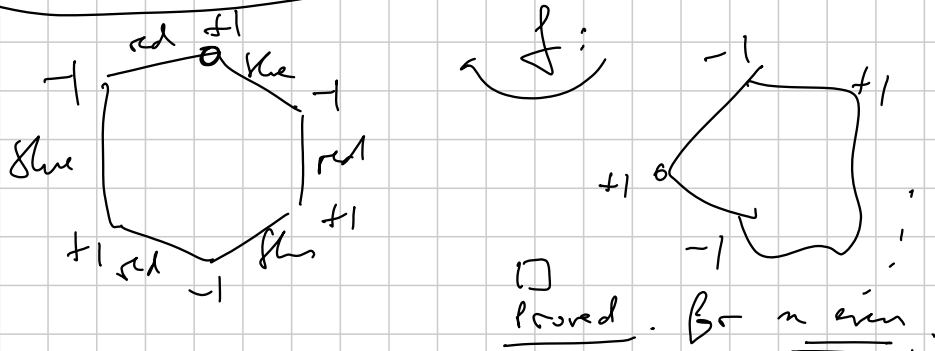
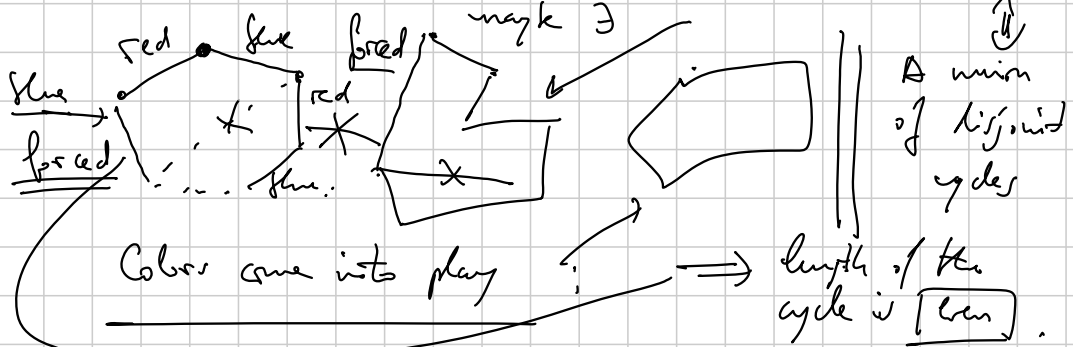
the vertices of my graph  $G: \{1, 2, \dots, n \in 2m\}$   
 edges?

Some red, some blue.

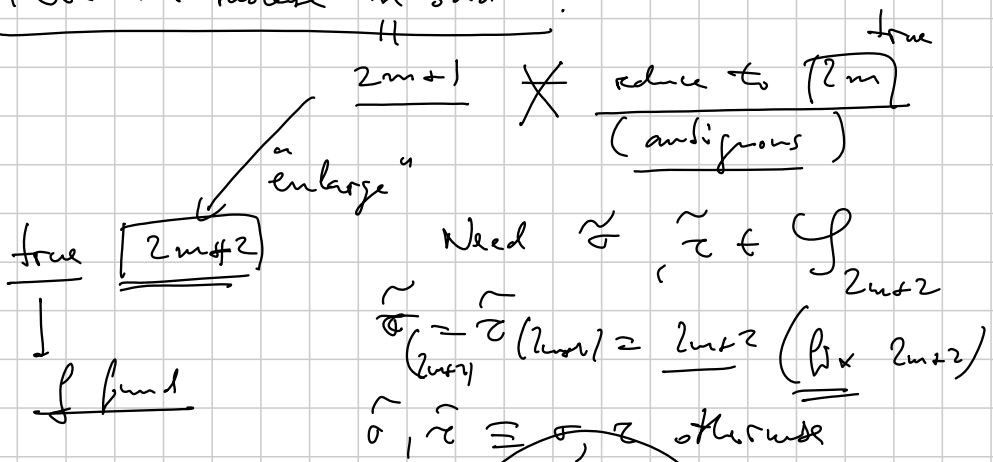
Structure of  $G$ ?  $\deg(v) = 2$  (1 red + 1 blue).

$G$  is a 2-regular graph.

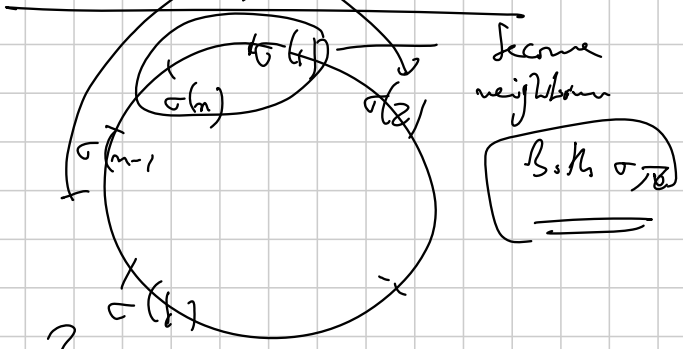
What do we know about 2-regular graphs?



Kill the subcase  $n$  odd!



One last word

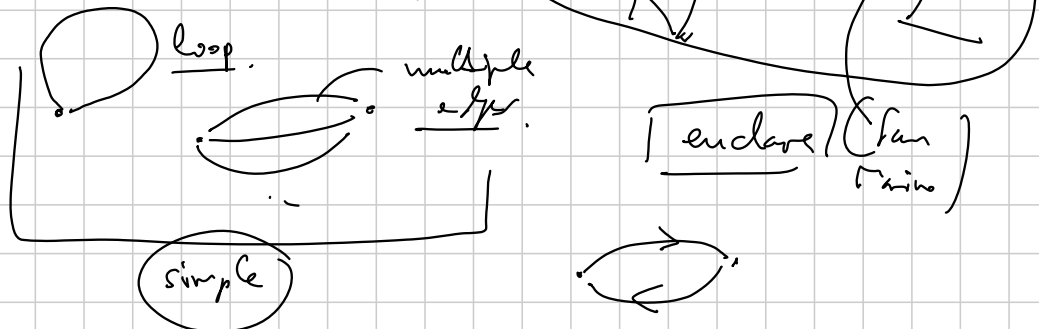


Q: Does the result stay true?

For  $n$  even — YES.  
 $n$  odd — NO — Need a model.  
 $\sigma = id$ .  
 $\tau = \begin{pmatrix} 1 & 2 & \dots & n \\ & 1 & & \end{pmatrix}$   
all  $\sigma$  (evens) | all  $\sigma$  (odds)      $\pi = \begin{pmatrix} 1 & 2 & 3 & \dots & n \\ u_1 & u_2 & & & u_n \end{pmatrix}$

Problem .  $G = (V, E)$

Oriented, directed, (digraph)



A complete digraph = tournament.  
 (Round-robin).

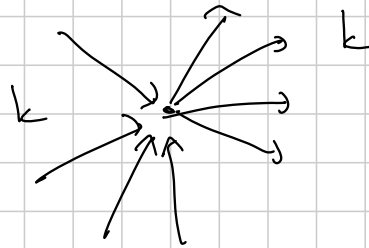
A graph  $G$  is called connected if  
 any two vertices are linked by some path.



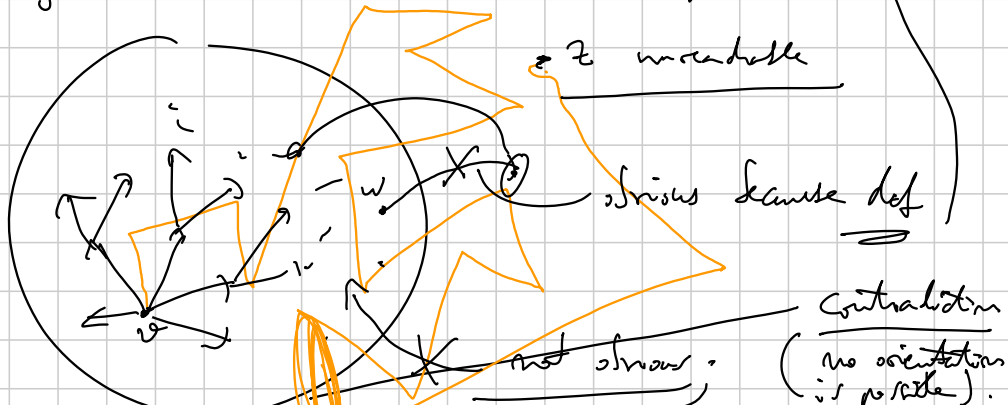
Definition For a digraph  $G$  → strongly connected if  $\forall A, B \rightarrow$  path.  
weakly connected but if then ignoring the orientation, it is



Problem. Any  $k$ -regular digraph which is weakly connected, is also strongly connected.

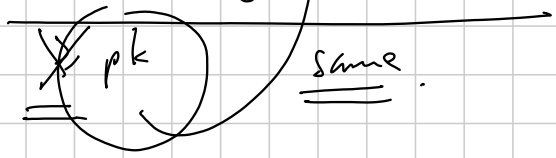


Solution. Start  $v$ ; build the set  $C_v$ , made of all  $w$  which are reachable from  $v$ .



$|C_v| = p$ ; ? arrows come out of  $x \in C_v$ ?

? arrows come in  $x \in C_v$ ?



SENIOR 2010 - DAN (COMBI + LINEAR ALGEBRA)

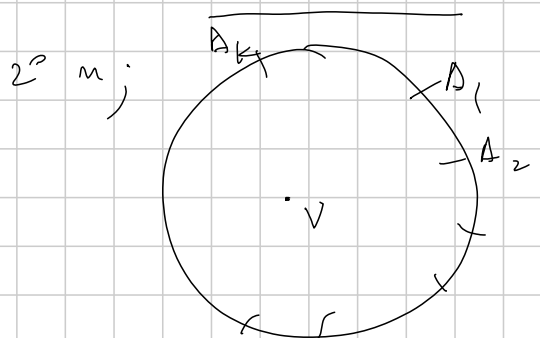
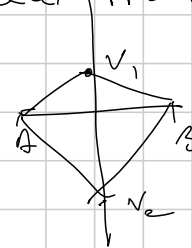
Titolo nota

08/09/2010

For starters IMO 1989? (death counting)  
 Given  $n, k$ . Given  $n$  points in the plane,  
 no 3 collinear (in general position); for each  
 point — there are  $\geq k$  other points situated  
 at a same distance from that point.  
 How large (find a bound) can  $k$  be?

Solution Count all isosceles triangles.

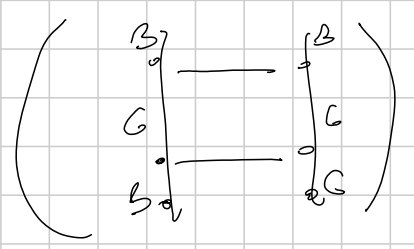
1°  $\binom{n}{2}$  ;  $N \leq n(n-1)$



Any  $A_i, A_j$   
 $N \geq n \frac{k(k-1)}{2}$

Problem 2 (China 2005?)

We have a  $2n \times n$  array. Boys & girls.



Any two columns,  
 There are at least as  
 many differences in gender  
 (by row) as there are  
 coincidences of genders.

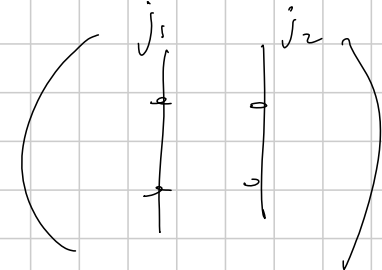
To establish a bound for  $B$  boys.

$i \begin{pmatrix} b_i & g_i = n - b_i \end{pmatrix}$

$N = \text{total \# of differences of gender}$

$$N = \sum_{i=1}^{2m} b_i j_i \quad \sum_{i=1}^{2m} j_i (n - j_i) \quad \dots$$

$$N \geq n \binom{2m}{2} = n \frac{2m(2m-1)}{2}$$

$$\sum_{i=1}^{2m} j_i = B$$


$$f(x) = x(n-x)$$

$$\sum_{i=1}^{2m} j_i (n - j_i) \leq 2 \frac{\sum_{i=1}^{2m} j_i}{2m} \left( n - \frac{\sum_{i=1}^{2m} j_i}{2m} \right) =$$

$$= 2 \frac{B}{2m} \left( \frac{2mn - B}{2m} \right)$$

$$n \frac{2m(2m-1)}{2} \leq N \leq B \cdot \frac{2mn - B}{2m}$$

$$B^2 - 2mnB + 2m^2(n-1) \leq 0 \quad \Delta = \frac{4m^2n^2 - 4m^2n(n-1)}{4m^2}$$

$$mn - n\sqrt{m} \leq B \leq mn + n\sqrt{m}$$

$$2m = 22, \quad m = 75$$

$$mn - n\sqrt{m} \approx$$

ELEMENTS OF LINEAR ALGEBRA  
(VECTOR)

$(K)$  "body" of numbers,  $+ \begin{pmatrix} 0 \\ 1 \end{pmatrix}$   
field  $(\mathbb{Q}, \mathbb{R}, \mathbb{C}, \mathbb{Z}_p)$ .  
of scalars.

$\nabla$  - set of vectors.  $\left\langle \begin{array}{l} \text{add, } \mathbf{0}, -v \\ \text{multiplication: scalar} \cdot \text{vector} \end{array} \right.$

$$\lambda(u+v) = \lambda u + \lambda v \quad \parallel \quad 1 \cdot v = v$$

$$\lambda(\mu u) = (\lambda\mu)u \quad \parallel \quad 0 \cdot v = \mathbf{0}$$

$$\lambda \cdot \mathbf{0} = \mathbf{0}$$

$$\sum_{i=1}^n \lambda_i v_i = \text{linear combination}$$

Can it be equal to  $\mathbf{0}$ ? without all  $\lambda_i = 0$

all  $v_i$ 's  
being null.

$l = 0$  but not trivial = (relation)

Consider those set of vectors where there are no relations. Such a set  $S$  linearly independent vectors.  $\emptyset \notin S$ .

Take any l.c.  $\sum_{i=1}^n \lambda_i v_i = u, v_i \in S$ .

unique.

~~$\sum_{j=1}^m \mu_j w_j, w_j \in S$~~

Basis.  $V =$  lin. ind set  $B$  such that any  $u \in V$  can be represented as  $\sum_{i=1}^n \lambda_i v_i, v_i \in B$ .

Finite  $B = \{j_1, j_2, \dots, j_d\}$

$K^d = \underbrace{K \times K \times \dots \times K}_d = \left\{ (x_1, x_2, \dots, x_d) \mid x_i \in K \right\}$   
 $K$ -plets

$0 = (0, 0, \dots, 0)$   
 $-v = (-x_1, -x_2, \dots, -x_d)$

$\lambda v = (\lambda x_1, \lambda x_2, \dots, \lambda x_d)$

Canonical basis :  $e_1 = (1, 0, 0, \dots, 0)$   
 $e_2 = (0, 1, 0, \dots, 0)$   
 $\vdots$   
 $e_d = (0, 0, \dots, 0, 1)$

$v = \sum_{i=1}^d x_i e_i$

$b_i \leftrightarrow e_i$

$V \cong K^d$   
isomorphic.

Theorem If there exists a finite basis  $B$  with  $d$  elements, then any other basis has also  $d$  elements.

Lemma (The Exchange Lemma)

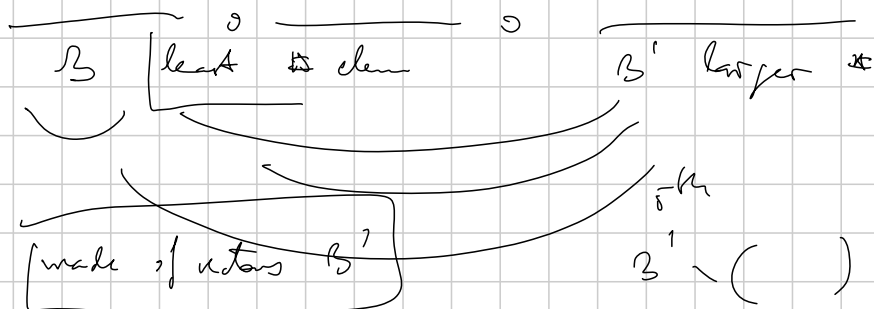
If  $\mathcal{V}$  have two bases  $\mathcal{B}, \mathcal{B}'$ ,  
 and chose  $b' \in \mathcal{B}'$ , there exists some  $b \in \mathcal{B}$ ,  
 so that  $\boxed{(\mathcal{B} \setminus \{b\}) \cup \{b'\}}$  is a basis.

Proof. Because  $\mathcal{B}$  basis,  $b' = \sum_{i=1}^m \lambda_i s_i$  ;  
 there must be  $\lambda_i \neq 0$

$$\lambda_i s_i = \sum_{j \neq i} \lambda_j s_j - b' \quad ; \quad \boxed{j \text{ divide by } \lambda_i}$$

$$s_i = \dots$$

□



$$\dim \mathcal{V} = \dim(K^d) = d = |\mathcal{B}|$$

$$u = (u_1, u_2, \dots, u_d)$$

$$v = (v_1, v_2, \dots, v_d)$$

$$u \cdot v = \langle u, v \rangle \stackrel{\text{def}}{=} \sum_{i=1}^d u_i v_i \in K$$

Example: All polynomials of degree  $\leq d-1$   
 coefficients in  $K$

$$\text{Basis} = \{1, x, x^2, \dots, x^{d-1}\}$$

$$\mathbb{R}^2 = \text{vectors}$$

$$\mathbb{C}$$

$$\mathbb{R}[x]$$

$$\langle u_1 + u_2, v \rangle = \langle u_1, v \rangle + \langle u_2, v \rangle$$

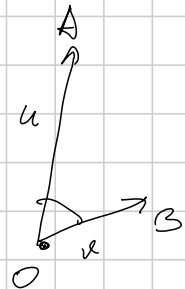
$$\langle \lambda u, v \rangle = \lambda \langle u, v \rangle$$

Dot product = called for a bilinear form

$$\langle u, u \rangle = \sum_{i=1}^d u_i^2 = \text{square of the Euclidean distance from } 0 \text{ to } u$$

$(K = \mathbb{R})$

What is  $\frac{\langle u, v \rangle}{\|u\| \cdot \|v\|} = \cos(u, v)$



Parseval - Schwarz

$$0 \leq \|u + \lambda v\|^2 = \langle u + \lambda v, u + \lambda v \rangle \quad (\text{variational methods})$$

$$= \langle u, u \rangle + 2\lambda \langle u, v \rangle + \lambda^2 \langle v, v \rangle =$$

$$= \|u\|^2 + 2\langle u, v \rangle \lambda + \|v\|^2 \lambda^2$$

$$\Delta \leq 0 \implies \langle u, v \rangle^2 - \|u\|^2 \|v\|^2 \leq 0 \quad (C-S)$$

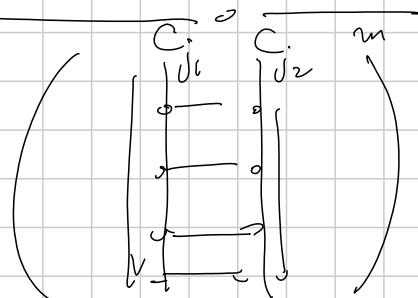
Back Chinese problem

Boys - +1

Girls - -1

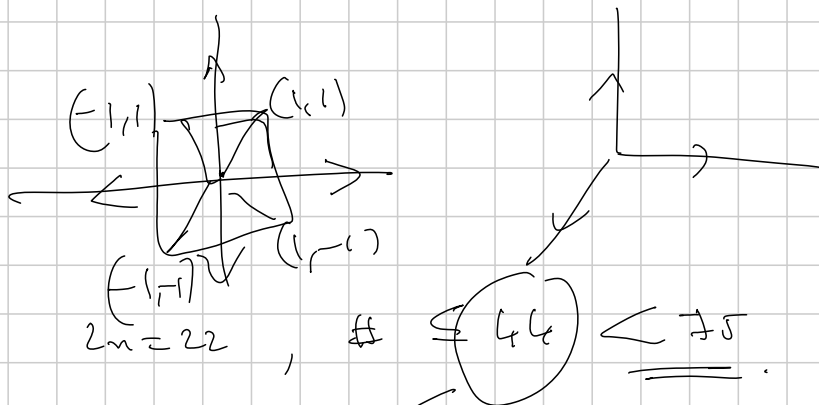
$$\langle \underline{c}_1, \underline{c}_2 \rangle = \sum_{i=1}^m l(c_1^i) \cdot l(c_2^i) \leq 0$$

$\mathbb{R}^{(2m)}$



Translates the given problem

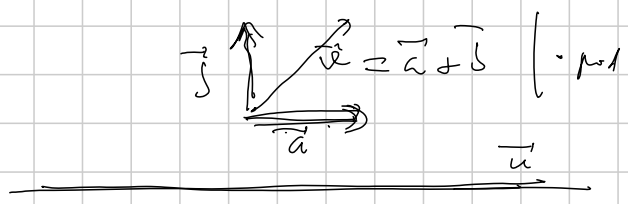
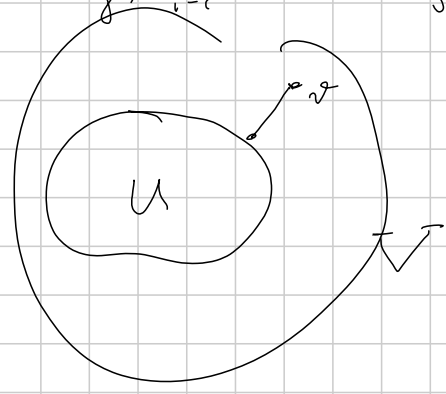
Result The max. number of vectors, which 2 by 2 make  $\geq \frac{\pi}{2}$  angle is twice the dimension of the space  $(4m)$



Definition:  $\langle u, v \rangle = 0$ , orthogonal  $u \perp v$   
 Assuming  $\|u\| = 0 \iff u = \mathbf{0}$

Proposition: A set of pairwise orthogonal vectors is linearly independent.

Proof: Assume  $\sum_{i=1}^n \lambda_i v_i = \mathbf{0}$  |  $\cdot v_j$   
 $\langle v_j, \sum_{i=1}^n \lambda_i v_i \rangle = \lambda_j \|v_j\|^2 + \sum_{i \neq j} \lambda_i \langle v_j, v_i \rangle = \lambda_j \|v_j\|^2$   
 (Note:  $\langle v_j, v_i \rangle = 0$  for  $i \neq j$ )



Problem (PUTNAM 2007?)

Given a  $m \times n$  array whose entries are  $\pm 1$ .  
 Look for the  $+1$ . Find some rows, columns  
 = subarray made of just  $+1$ 's I J

Prove  $|I| \cdot |J| \leq \lfloor \frac{n}{2} \rfloor$

We know that for any two rows  $R_i, R_j$

$$\sum_{k=1}^m x_{ik} \cdot x_{jk} = 0$$

$\langle R_i, R_j \rangle = 0 \quad ; \quad \boxed{R_i \perp R_j}$

$n=1$   $(1, 1, \dots, 1)$   $n=2m'$

$n=2$

$$\begin{pmatrix} +1 & - & +1 & +1 & - & +1 \\ +1 & - & +1 & - & - & - \end{pmatrix}$$

Solution. My rows = vectors,  $n$ -dimensional.

Assume  $I, J, R = \sum_{i \in I} R_i$  ;

$$\begin{aligned} \|R\|^2 &= \langle R, R \rangle = \langle \sum_{i \in I} R_i, \sum_{j \in J} R_j \rangle = \\ &= \sum_{i \in I} \|R_i\|^2 = n|I|. \\ \|R\|^2 &= \sum_1^m (\text{coordinates})^2 = \sum_{j \in J} (\ )^2 + \sum_{j \notin J} (\ )^2 \\ &\geq \sum_{j \in J} (\ )^2 = |I|^2 \cdot |J| \\ n|I| &\geq |I|^2 \cdot |J| \rightarrow \boxed{|I| \cdot |J| \leq n} \quad \square \end{aligned}$$

Problem. Set  $[A] = m, A_1, A_2, \dots, A_m \subseteq A$   
 $|A_i| = \text{odd}, |A_i \cap A_j| = \text{even}, i \neq j$ .  
 Prove  $\underline{m \leq n}$ . Trivial example if we take all singletors.



$$B \subseteq A, \chi_B : A \rightarrow \{0, 1\} = \mathbb{Z}_2$$

$$|A|=m \quad \chi_B(a) = \begin{cases} 0 & \text{if } a \notin B \\ 1 & \text{if } a \in B \end{cases}$$

$$\chi_B : \left( \begin{matrix} 0/1, & 0/1 & \dots & 0/1 \\ \text{1st} & \text{2nd} & & \text{nth} \end{matrix} \right) \rightarrow$$

$$B \rightarrow \chi_B = \left( \begin{matrix} \phantom{0} \\ \phantom{0} \\ \phantom{0} \\ \phantom{0} \end{matrix} \right) \in \mathbb{Z}_2^m$$

$$B \Delta C = (B \cup C) \setminus (B \cap C)$$

$$(B \setminus C) \cup (C \setminus B)$$

$$\chi_B + \chi_C = \chi_{B \Delta C}$$

$$A, \left( \mathcal{P}(A), \Delta, \cap \right) \text{ Boolean ring.}$$

$$A_1, A_2, \dots, A_m \quad \chi_{A_1}, \chi_{A_2}, \dots, \chi_{A_m}$$

$$|A_i| = \text{odd} \iff \|\chi_{A_i}\| = 1$$

$$|A_i \cap A_j| = \text{even} \iff \langle \chi_{A_i}, \chi_{A_j} \rangle = 0$$

$$\{ \chi_{A_i} \} \text{ linearly independent } \subset \mathbb{Z}_2^m \text{ - } n\text{-dimensional}$$

$m \leq n$

Problem. ("baby" Lindstrom). (China)

$$\text{Again } |A|=m, \quad A_1, A_2, \dots, A_{n+1} \subseteq A$$

$$\text{Prove there exist } I, J, \quad I \cap J = \emptyset, \quad \text{non-empty}$$

$\{1, 2, \dots, n+1\}$

$$\text{such that } \bigcup_{i \in I} A_i = \bigcup_{j \in J} A_j.$$

Counterexample for  $n$  subsets: all midpoints.

Solution.  $A_i \rightarrow X_{A_i}$  vectors. not such.  
 $\prod_{i=1}^m \mathbb{R}^m$  (same  $\mathbb{Z}_2^m$ )  
 $n$ -dimensional vector space

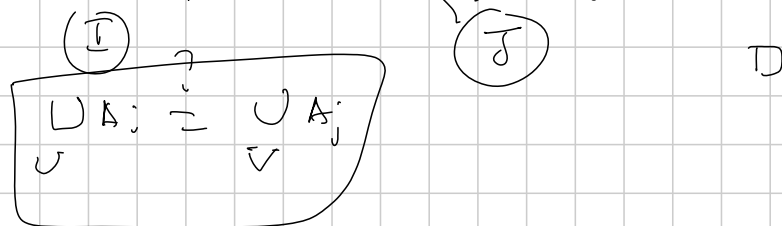
Applying Lin. Alg.

$\Rightarrow$  There must exist a non-trivial relation between  $m+1$  vectors:  $\sum_{i=1}^{m+1} \lambda_i X_i = 0$ ,  
 without all  $\lambda_i$  being null. all not null

$\Rightarrow$  (there must be  $\geq 2$  such not-null scalar  $\lambda_i$ 's)

$\Rightarrow \sum_{i \in U} \lambda_i X_{A_i} = \sum_{j \in V} (-\lambda_j) X_{A_j}$   $U, V \subseteq \{1, 2, \dots, m+1\}$   
 $\lambda_i \neq 0$

/ positive  $\lambda_i$                       / negative  $\lambda_j$



Problem (Lindström).  $A_1, A_2, \dots, A_{m+2}$

one can find  $I, J$  sets of indexes, disjoint, s.t.

$\cup_{i \in I} A_i = \cup_{j \in J} A_j$ , sub-alge  $\cap_{i \in I} A_i = \cap_{j \in J} A_j$

Solution. Counterexample; all singletons, and  $A$

2 spaces.  $1^{st}$   $A_i \rightarrow (x_{A_i}, x_{\bar{A}_i})$   $2n$ -vector  
 $(x_1, x_2, \dots, x_n, 1-x_1, 1-x_2, \dots, 1-x_n) \in \mathbb{R}^{2n}$

(set of vectors  $S \subseteq V$ ; consider all  $\sum_{i=1}^n \lambda_i x_i$ ,  $x_i \in S$ )  
 $\langle S \rangle \subseteq V$  a linear space generated by  $S$ .

$(A| = n, \text{ all } B \subseteq A$   $(x_B | x_A) \in \mathbb{R}^{2n}$   
 $\dim \langle S \rangle \geq n$  ? generate a copy of  $\mathbb{R}^n$   
 $=$  ?  $\boxed{n+1}$

There must be a non-trivial relation.

$$\sum_{i=1}^{n+2} \lambda_i \begin{pmatrix} x_{A_i} \\ x_{B_i} \end{pmatrix} = 0$$

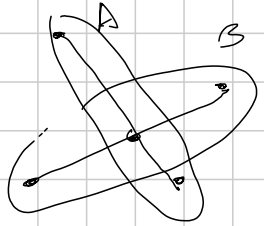
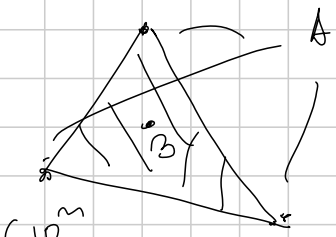
We separate positive  $\lambda$ 's = negative  $\lambda$ 's

2<sup>nd</sup> Helly's Theorem. If a finite collection of convex sets intersect 3 by 3; then they all have a common point.

$\mathbb{R}^m$ , need to intersect in pairs of  $n+1$

Radon's Theorem If  $n+2$  points in  $\mathbb{R}^m$  then there is at least one way to partition them into two non-empty sets  $A, B$  such that  $\boxed{\text{Conv}(A) \cap \text{Conv}(B) \neq \emptyset}$

$\mathbb{R}^2$ :  $\boxed{n+2=4}$ :



$x_{A_i}$   $1, 2, \dots, n+2$  can be partitioned.  $A, B$   
 $\implies x \in \text{Conv}(A) \cap \text{Conv}(B) \neq \emptyset$ .

$$\| \sum_{i \in A} \lambda_i x_i = \sum_{j \in B} \mu_j x_j \|$$

$x$  writes in  $\text{Conv}(A)$        $x$  writes  $\text{conv.}(B)$

affine combination (non scalar is)

not just linear combinations they are convex combinations

scalars  $\geq 0$  their sum = 1

R Memento → How many points,  $(N)$   
 (no 3 collinear), can we have in the plane,  
 note that for any three → at least one distance  
 equal to 1

$N \geq 6$



$N \geq 7$

third would  
 more  
 & cannot  
 /B

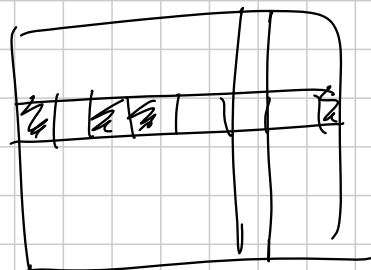
ADVANCED - 2010 - COMBINATORICS MISCELLANEA

Titolo nota

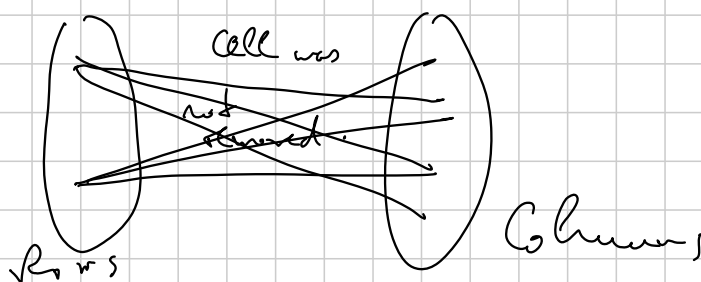
10/09/2010

Problem 1. A  $2n \times 2n$  chessboard; remove  $k$  cells, but not more than  $n$  on any row or column.

Prove that one can still place  $2n$  rooks so that they do not attack each other



Solution - If no cells removed  $(2n)!$  ways.



Answer = perfect matching

Marriage Theorem. Bipartite graph  $|A|=|B|$ ,

if for any  $S \subseteq A$ ,  $|S| \leq |\Gamma(S)|$ .

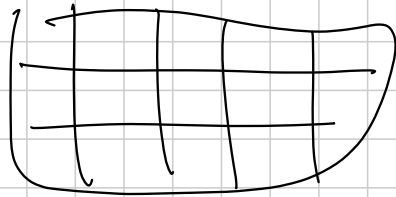
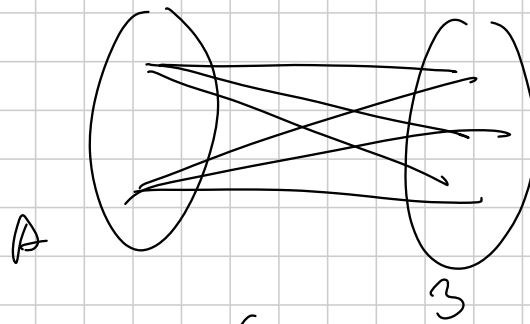
$\implies$  there is a perfect matching.

If  $|S| \leq n$ , ok: even one row  $R \in S$

If  $|S| > n$ ,  $|\Gamma(S)| = 2n$

By M.T.  $\implies$  Perfect matching can be done.

Particular case  
k-regular



Set (finite) =  $mk$  elements  
 $\pi$  partition into  $m$  classes  
of  $k$  elements each.

Find  $m$  elements which represent both partitions.

$G$  group,  $H$  subgroup

$$G = \bigcup_{i=1}^r Hx_i = \bigcup_{i=1}^r y_i H$$

index                      index  
 $x_i$                        $y_i$   
 $z_i$                        $z_i$

Problem 2. (A.M.R) Island A,  $n$  families.

Ministry Agriculture : partitioned A into  $n$  equal

Ministry Hunting :  $1, 1, \dots, 1$

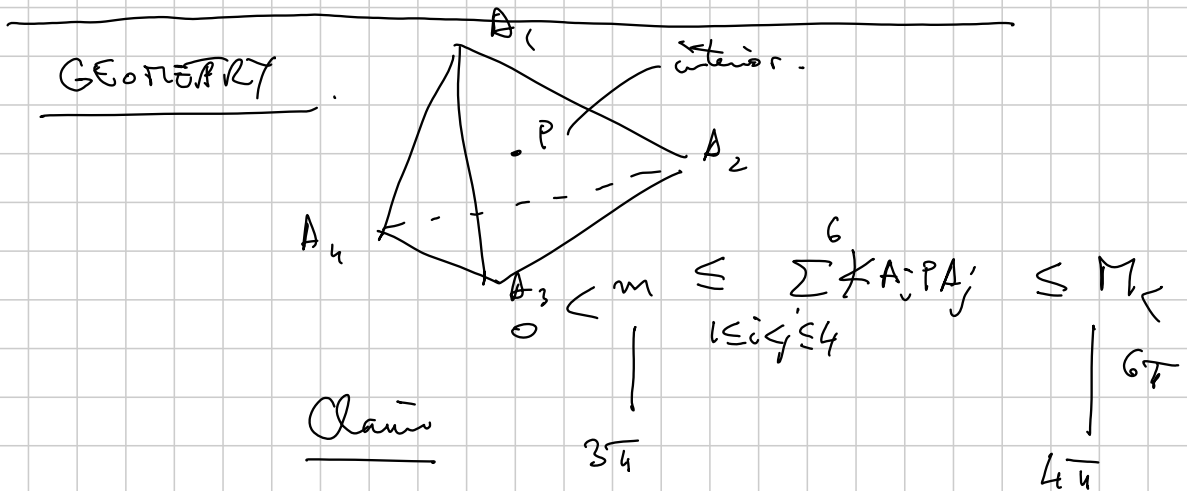
Ministry Family : we should allocate to each family one each so that they interested

they succeeded!

Ministry Religion : declared a miracle!

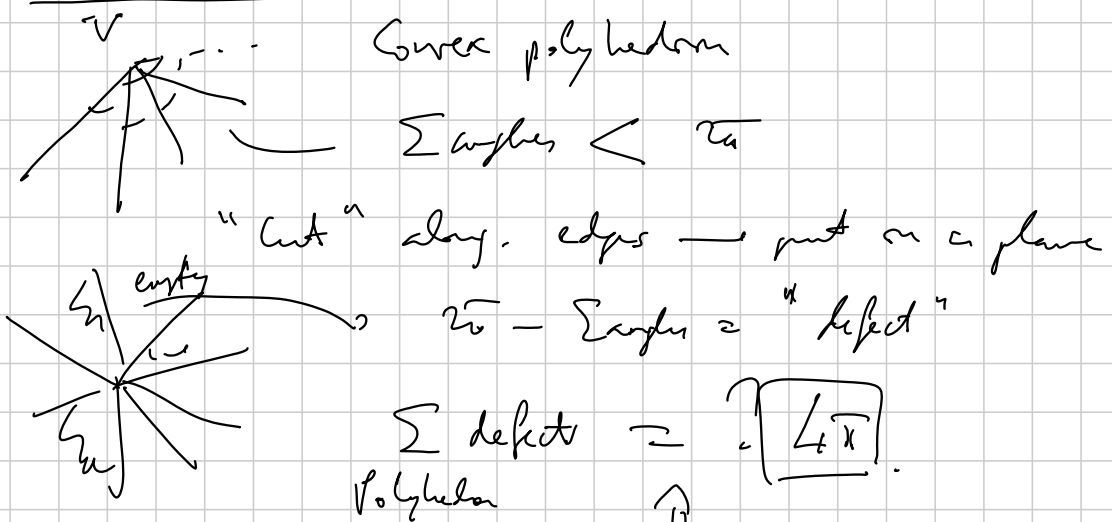
a) Math. Institute :  $\rightarrow$  reported: it could always be done!

b) Find the fraction  $\lambda$  of sea B island, so that by a dove allocation, each family receives a  $|C_i \cap H_i| \geq \lambda$ .



$\therefore \sum \approx 3 \cdot 1 + 3 \cdot 0 = 3T$

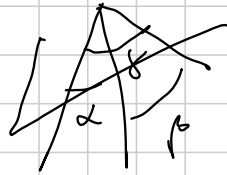
$\therefore \sum \approx 2 \cdot 0 + 4 \cdot 1 = 4T$



Euler:  $V + F = E + 2$

Proof.  $\sum$  angles  $< 2\pi$   
around a vertex

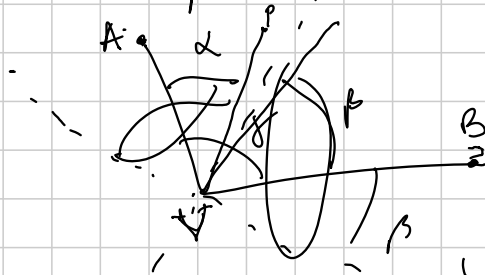
show Lemma. Vertex of a ~~trichedron~~ ~~trichedron~~.



$\alpha, \beta, \gamma$  have the "triangular" inv. property

$$\begin{cases} \alpha + \beta \geq \gamma \\ \beta + \gamma \geq \alpha \\ \gamma + \alpha \geq \beta \end{cases}$$

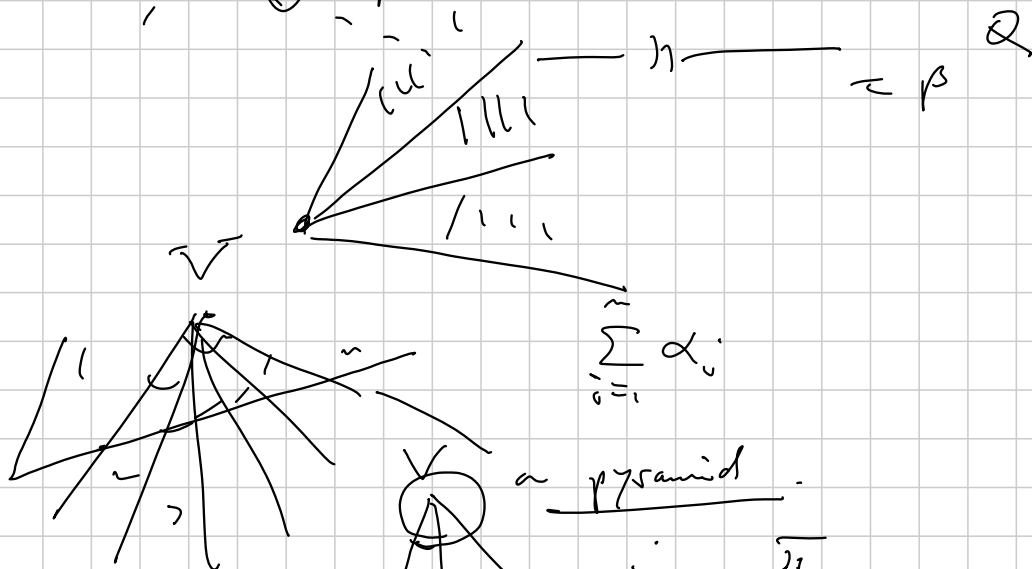
In the plane draw the locus of  $\gamma$ :



Assume  $\alpha + \beta < \gamma$ .

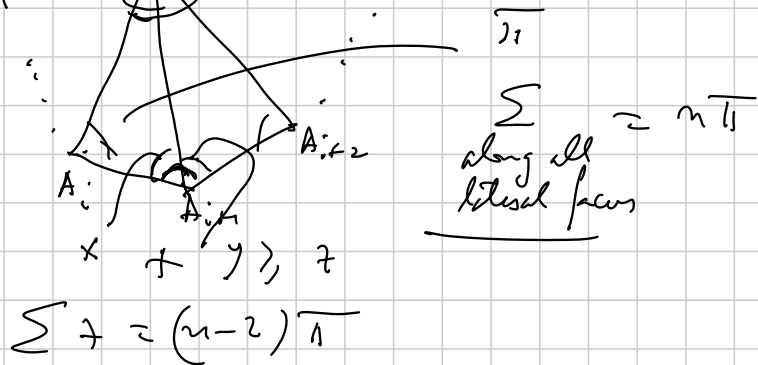
Locus in  $\mathbb{R}^3$  of points P

such that  $\angle PVA = \alpha$  ?



$$\sum_{i=1}^n \alpha_i$$

a pyramid.



$\sum$  along all lateral faces  $\approx n\pi$

$$x + y \geq \gamma$$

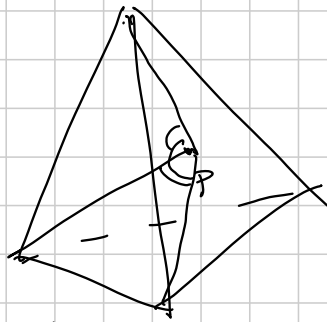
$$\sum \gamma = (n-2)\pi$$



$$\sum_{i=1}^6 \angle A_i P A_i \leq 4\pi$$

$$\sum \leq 2\pi$$

for all pairs of 3 out of the 4 vertices



Four times  $\frac{2\sum}{4} \leq \frac{4\pi}{4}$

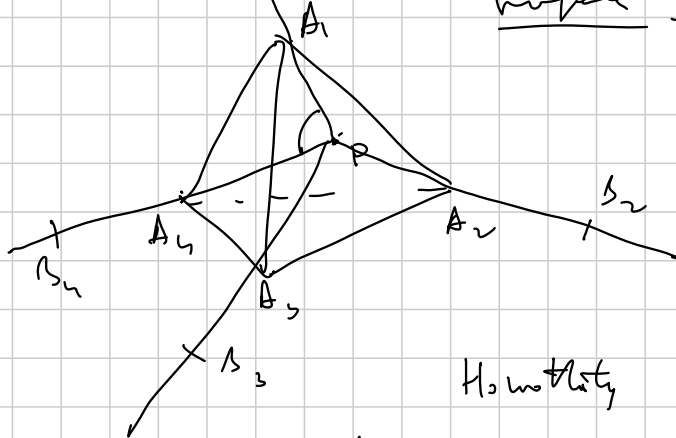
$$3 \leq \sum \leq \widehat{A_1 P A_3} + \widehat{A_3 P A_4} + \widehat{A_4 P A_2} + \widehat{A_2 P A_1}$$

(1, 3, 4, 2)

Lateral. Why is this happening (is it?)

One way  $\rightarrow$  you do a nice drawing on a sphere.

"geodesic" = "arc" of minimal length that connects 2 points on a surface.

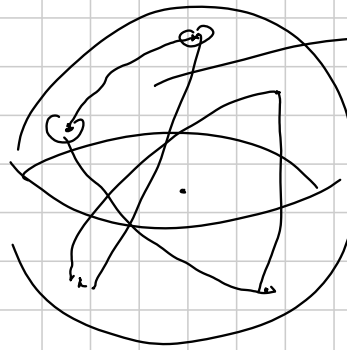
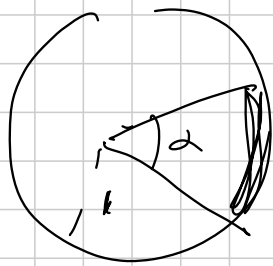


Can take  $B_i$  such that  $PB_i = \text{rad}$ .

$P$  center of the sphere  $\triangle B_1 P B_3 B_4$  inside.

Homothety  $= 1$ .

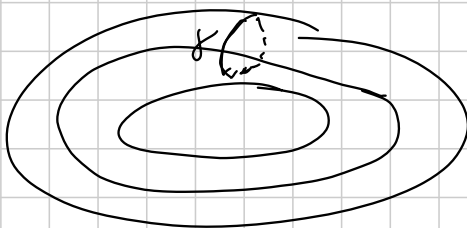
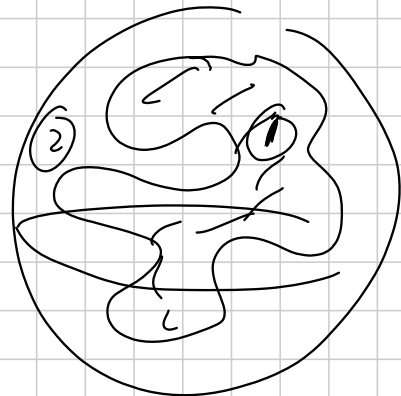
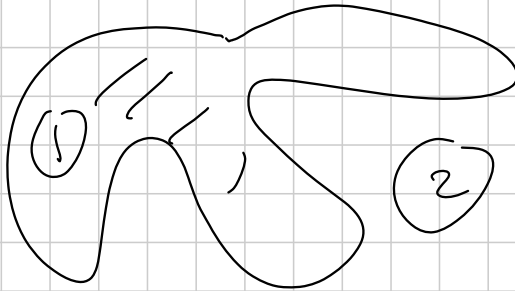
The advantage: measure of  $\angle$  = length of the arc great circle is subtended.



length of  $h$   $\geq 2l$

WRONG PICTURES

Combi fact 1. If  $\gamma$  on surface of sphere of radius  $l$ , curve  $\gamma$  simple (does not auto-intersect), closed, rectifiable (its length can be calculated). And if  $\gamma$  not a circle, divides surface of sphere into 2 regions, equal area, Then length  $|\gamma| \geq 2l$ .



Torus (donut)

Proof.  $\varphi: S^2 \rightarrow S^2$   
 $p \rightarrow p' = \text{antipode}$   
 $\varphi(\gamma) = \text{curve}$   
 $\gamma$  blue  
 $\varphi(\gamma)$  red

$\gamma \subset \text{Hemisphere}$

$|\gamma \cap \varphi(\gamma)| \geq \frac{2(\text{pairs})}{p}$

$P, P' \in \gamma \quad (\varphi(\gamma))$

$\lambda(\gamma) = \widehat{PP'} + \widehat{P'P} \geq \frac{1}{2} \text{circ} + \frac{1}{2}$

$= \text{full part circle} \leq \pi$

□

$\lambda(\gamma) < \pi$

$\exists! P'$

$\lambda(PP') = \lambda(P'P)$

arc of great circle  $PP'$

plane  $\perp$  OM

" middle "

$\gamma \subset \text{Hemipl.}$

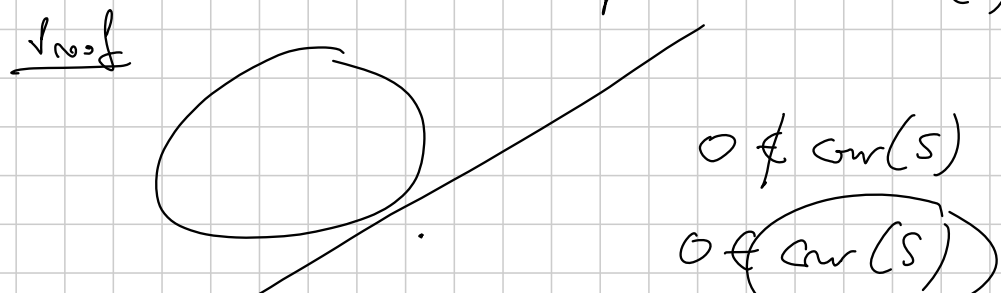
Either  $(PP')$   $PP' \cap$  "equator"

Assume not

$\text{equator} \cap P' \rightarrow$  symmetric with respect to  $O\Gamma$   
 $\text{arc} \subset S$   
 $\text{same length} = \frac{1}{2} \lambda(\gamma)$   
 $X' \downarrow$  also on equator (opposite)

Instead looking at  $\gamma$  :  $(PP' \cup \text{symmetric arc})$   
 contains  $X, X'$   $= \lambda(\gamma)$

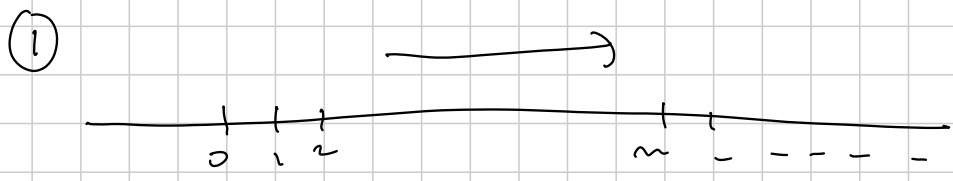
My solution : (1) The fact that a set  $SCS^2$  is not contained in ANY hemisphere  $\Leftrightarrow \text{Conv}(S) \ni O$ .



(2) Carathéodory Theorem  $\forall \gamma$   
 There must be  $k$  line  $A_1, A_2, A_3, A_4 \ni O$

(3)

Monty Python : And now, on a different...  
 problems with hats, names...



Establish an equivalence relation among infinite binary words

wwwwww...  $\xrightarrow{\quad}$

$X \sim Y$  ~~def~~ they differ only in finitely many positions.

Agree on a representative  $X \in \hat{X}$  (Axiom of choice)

Each one  $\rightarrow$  write the color on the position that he skipped.

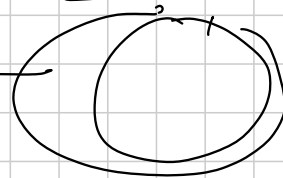
Finite variations  $1 \rightarrow n$

The  $n$  persons again not allowed to speak

$n$   $\rightarrow$  a tag write down a color, they also pronounce it.

But instead  $\rightarrow$

$(n)$  people,  $(n)$  colors of hats



Strategy 1<sup>st</sup> give themselves some number  $1 \rightarrow n$   
2<sup>nd</sup> give the color some color  $1 \rightarrow n$

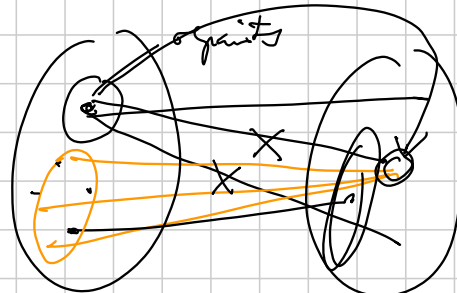
Each will write:  $\left[ \left( \sum \text{color he sees} \right) - \text{his position} \right] \pmod n$

$= \sum_{\text{subset}}^{\text{all color}}$  - the color of that person hat

In a forest there  $n$  deer  $\left\{ \begin{array}{l} \text{blue nose} \\ \text{pink} \end{array} \right.$

Each one, in turn, visits all his friends.  
 If strictly  $> 1/2$  his friends have a lower  
 different color than his  $\rightarrow$  he repaints.  
 Show, after some time  $\rightarrow$  there is no need  
 to repaint.

Sikinia.



BLUE PINK

The number of edges between B/P is  $\downarrow$  strictly

The (2n) drawers — game:

In cabin there are 2n boxes, each  
 containing the name of one dwarf.

One by one, each enters cabin, allowed to  
 open  $\leq n$  boxes  $\rightarrow$  finding his name.

They win if ALL find their names.  
 What strategy to maximize the chance to win?

From the "idiot"  $\rightarrow \left(\frac{1}{2}\right)^{2n}$  chance  
 a  $\frac{1}{n}$  chance by improvement

Sketch:  $k \rightarrow$  opens box  $k$  — if  $D$   
 if not opens box  $\leftarrow$   
 $\sigma(1) \sigma(2) \dots \sigma(2n)$

$$\sigma = (1, \sigma(1), \sigma(\sigma(1)) \dots) (3, \sigma(3) \dots) \dots \binom{2n}{\sigma(k)} =$$

$$\Rightarrow \sigma = \prod_{i=1}^s \gamma_i, \quad \sum \lambda(\gamma_i) = 2n.$$

They are same of  $\lambda(\gamma_i) \leq n, \forall i$ .

Estimate  $\frac{\# \sigma \text{ with cycles } \leq n}{\# \text{ perm}} = \frac{\# \sigma \text{ with cycles } \leq n}{(2n)!}$ .

Count  $N$  of  $\sigma$  which do have cycles of  $\lambda > n$

Let's count  $\# N_k$  of  $\sigma$  with a cycle  $\lambda = n+k$   
 $k=1, 2, \dots, n$

$$\binom{2n}{n+k} \cdot \text{rest} \cdot \text{arbitrary} \cdot (n-k)!$$

$$\binom{2n}{n+k} \cdot (n+k-1)! \cdot (n-k)! = \frac{2n!}{n+k}$$

$$p = \frac{N}{(2n)!} = \frac{1}{n+1} + \frac{1}{n+2} + \dots + \frac{1}{2n} \approx \ln 2$$

$$p = 1 - \ln 2 \approx 0.7 \approx 30\%$$

LARS → (Martin GARDNER)

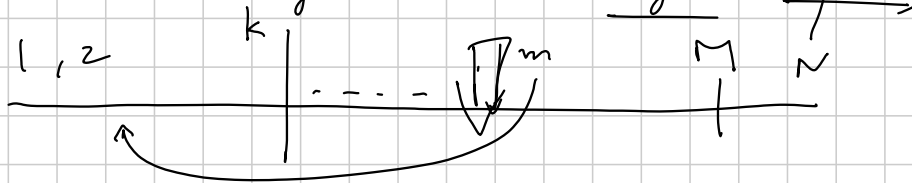
You are given  $N$  folded papers, each a real number

You pick the largest number distinct.

The "idiot" — pick one  $\frac{1}{N}$

the only possible strategy is:

- pick  $1 \leq k \leq N-1$
- open  $k$  pages — put them aside.
- start opening, until hit one larger than those rejected



$k \rightarrow p_k = \text{probability of missing best which } k \text{ max}$

$$k \approx \frac{N}{e}, \quad p_k \approx \frac{1}{e}$$

1...2...3, SCAPPATE!!! XD

MA MEGLIO MALE CHE CI SONO I CAPRONI

"PUO' BACIARE LA SPOSA!"

"MA E' UN LUPO!!!"

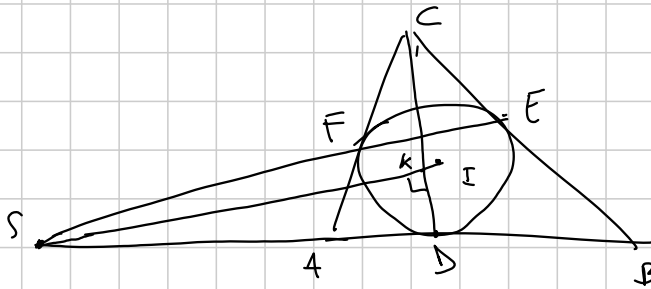
"NOI PRENDIAMO X ENO' NON LA PAI LA PROLETTIVITA' CHE' MANA P ALL' INFINITO!"



Titolo nota

10/09/2010

1. ABC triangolo, I incentro, D, E, F punti di tangenza, S = EF ∩ BA. Th: SI ⊥ CD

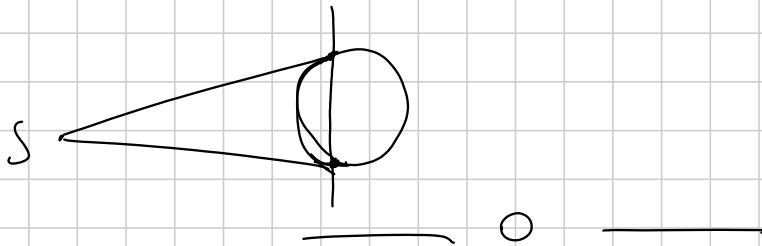


$D \in \text{pol}(S)$

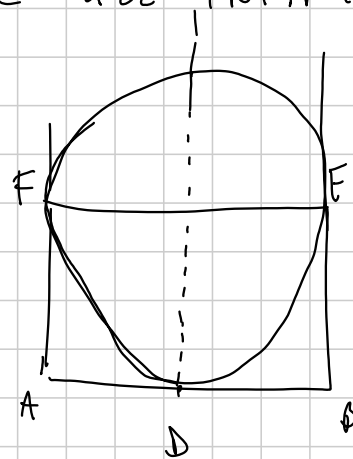
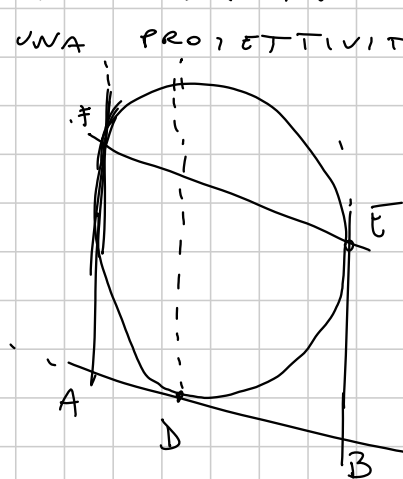
$S \in \text{pol}(C) \Rightarrow$

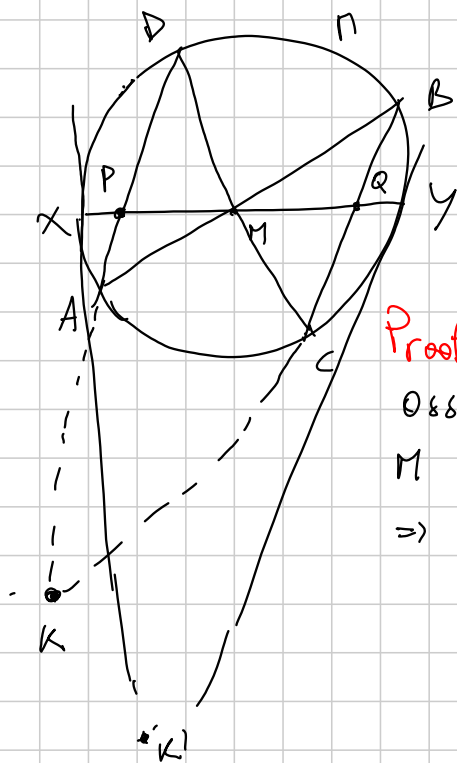
$C \in \text{pol}(S)$

$\Rightarrow \text{pol}(S) = CD$



2<sup>a</sup> SOL. MANDO LA RETTA SC ALL'INFINITO CON UNA PROIETTIVITA'



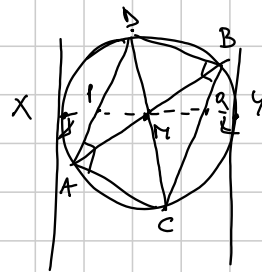


Butterfly theorem

H<sub>0</sub>:  $XM = MY \rightsquigarrow X, M, Y, \infty$   
 quater. armonica

Th:  $PM = MQ \rightsquigarrow P, M, Q, \infty$   
 armonica

Proof: Mendo la retta  $KK'$  all'infinito  
 Osservo  $CAF$  e' LA POLARE di  
 $M \Rightarrow KK' \cap \Gamma = \emptyset$   
 $\Rightarrow \Gamma$  si puo' mantenere cfr.

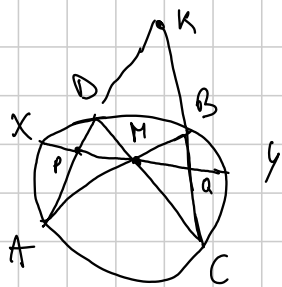


$pol(M) \perp OM$   
 $OM \perp XY$   
 $\Rightarrow pol(M) \parallel XY$   
 $\Rightarrow \infty \in pol(\Gamma)$

Nel disegno nuovo:

- le tangenti in  $x$  e  $y$  a  $\Gamma$  sono parallele  
 $\Rightarrow xy$  e' diametro di  $\Gamma$
- $X, M, Y, \infty_{xy}$  rimane una quaterne armonica  
 $\Rightarrow MX = MY \Rightarrow M$  e' il centro di  $\Gamma$
- $M \in BC \Rightarrow BC$  e' diametro  $\Rightarrow \hat{CAD}$  e' retto  
 e similmente gli altri angoli di  $ACBD$
- Ora la figura risulta simmetrica  $\Rightarrow PM = MQ$   
 $\Rightarrow P, M, Q, \infty_{xy}$  e' quaterne armonica  
 $\Rightarrow$  Thesis !!

□.

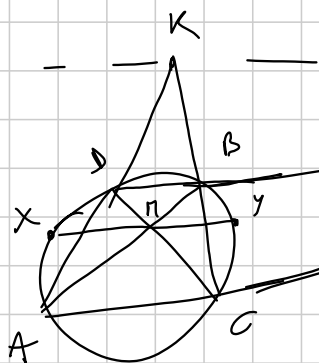


Birapporto  $X, B, D, Y$   
 proiettando su  $c$  ottago

$$(X, B, D, Y) = (X, Q, M, Y) = (X, M, P, Y)$$

$$\frac{XM \cdot QY}{XY \cdot QM} = \frac{XP \cdot DY}{XY \cdot MP}$$

$\Rightarrow$   $Q, P$  dividono due segmenti uguali ( $XM$  e  $MY$ ) nella stessa proporzione  $\Rightarrow PM = PQ$

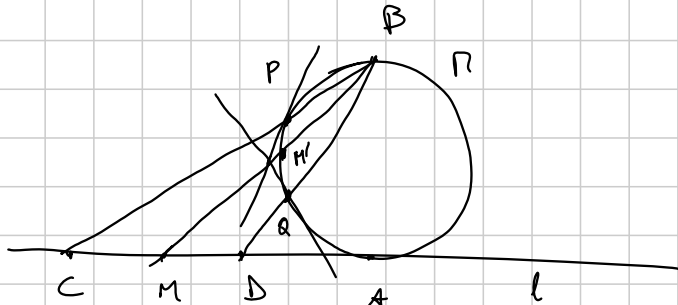


$KJ = \text{pol}(M)$   
 $KJ \cap XY = J'$

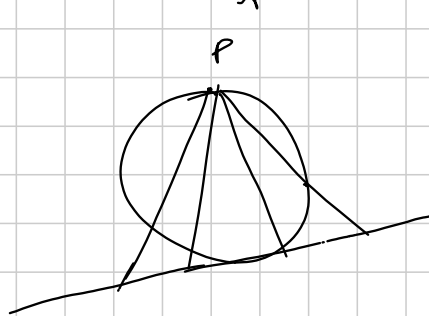
$$\boxed{|(X, Y, M, J') = -1| \quad ??}$$

Così dimostro  $KJ \parallel XY$

$C, M, D, A \in l$        $CM = MD$        $w$  tangente  $l$  in  $A$   
 $B$  diam. opp. ad  $A$  in  $w$



tangente in  $P$   
 tangente in  $Q$   
 $BM$  } Concorrono



1. invertito in  $B$  dimodeste  
 $P$  va in  $C$  (di raggio  $BA$ )

$$P \leftrightarrow l$$

$$\boxed{BCD \sim BQP}$$

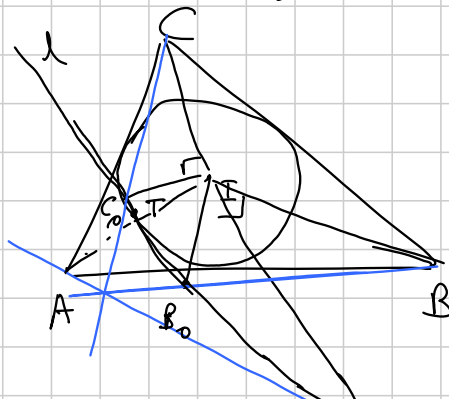
$BM$  e' simmediata  
 $\Rightarrow$  lemma simmediata

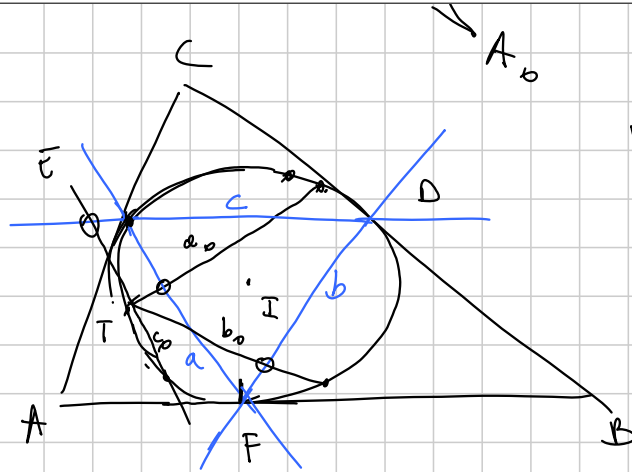
$BPM'Q$  e' quadrilatero armonico  
 $\Rightarrow$  tangenti e diagonali concorrono



18.  $ABC$  triangolo  $I$  incentro.  $SM$   $l$   
 tangente all'incirchio. siano  $A_0, B_0, C_0 \in l$   
 $t_{A_0} C$        $\widehat{AA_0} = \widehat{BB_0} = \widehat{CC_0} = 90^\circ$

Th,  $AA_0, BB_0, CC_0$  concorrono





retta di Simson

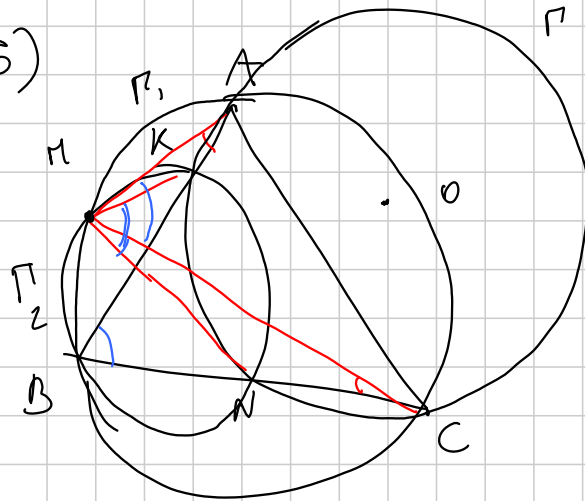
$$c \perp C_0$$

$$a \perp a_0$$

$$b \perp b_0$$

$T \in \Gamma \Rightarrow$  triangolo pedale di degenere  $\Rightarrow$  i punti sono allineati

(IMO 1985)

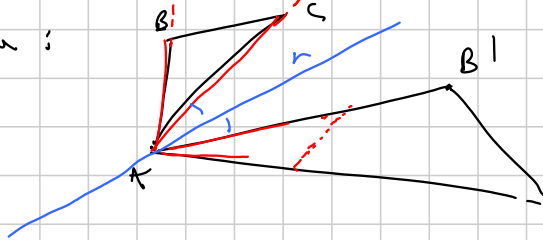


Th:  $OM \perp MB$

$KN, MB, AC$  sono concorrenti (assi radicali)

$\widehat{MAK} = \widehat{MCN}$   
 $\widehat{KAN} = \widehat{CBN}$   
 $\widehat{ABC} = \widehat{AMC}$

Lemma:



$$\widehat{ABC} \sim \widehat{AB'C'}$$

$\Rightarrow$   $\exists$  inv. in A  
 composta con una sim. assiale

che manda  $B \rightarrow C'$  e  $C \rightarrow B'$

$$\text{Preto } r_{\text{aggio}}^2 = AB \cdot AC \stackrel{\text{sim.}}{=} AC \cdot AB$$

Simetria wrt  $r$  fa esattamente quello che vogliamo.

$\Gamma$  rimane dunque in se stessa.

1. poiché il raggio rimane lo stesso  $\Gamma \rightsquigarrow \Gamma$  con l'inversione.

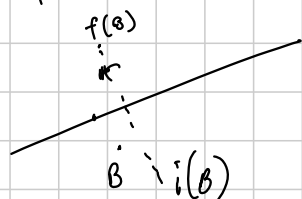
2.  $\Gamma \rightsquigarrow \Gamma$  con la riflessione  $\Rightarrow O \in r$

B dove va?  $B = \Gamma_1 \cap \Gamma_2$

$$i(B) = i(\Gamma_1) \cap i(\Gamma_2) = i(A)i(C) \cap i(K)i(N)$$

$$f(B) = f(A)f(C) \cap f(K)f(N) = NK \cap CA \in MB$$

$\Rightarrow f(B)$  è su  $MB$ , ma anche  $i(B) \in AB$



$$\Rightarrow i(B)P(B) \cap r = M$$

$$i(B)P(B) \perp r$$

$$r \perp i(B)M = BM$$

$$OM \perp BM$$

$i(P)$  = immagine di  $P$  sen  
l'inversione

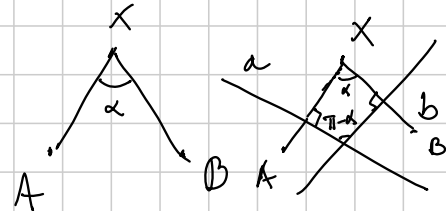
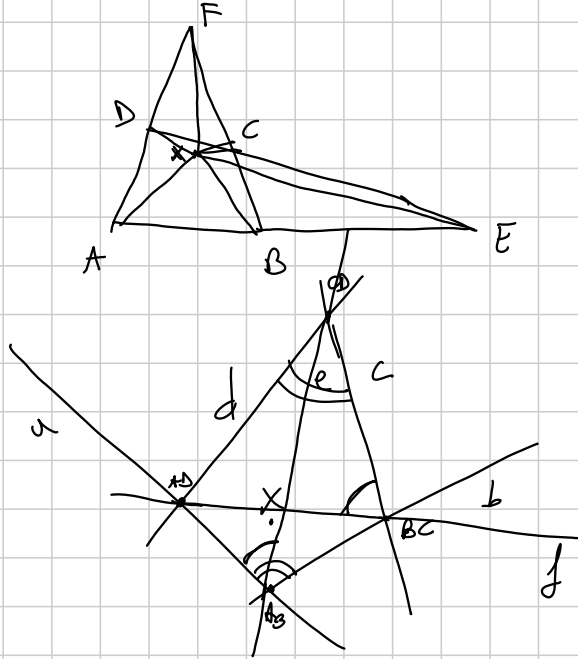
$f(P)$  = immagine di  $P$  sen  
inversione + riflessione

□

ROBA DI VERTECHI: <http://veluca.altervista.org/olimpiadi/>

Yufei Zhao (billzho su ML). (fonte degli esercizi)

14.  $ABCD$  quadrilatero convesso.  $AB \cap CD = E$ ,  $AD \cap BC = F$   
 Esiste  $X$  dentro il quad. s.t.  $\widehat{AXE} = \widehat{CXF}$ . Th:  $\widehat{AXB} + \widehat{CXD} = 180^\circ$

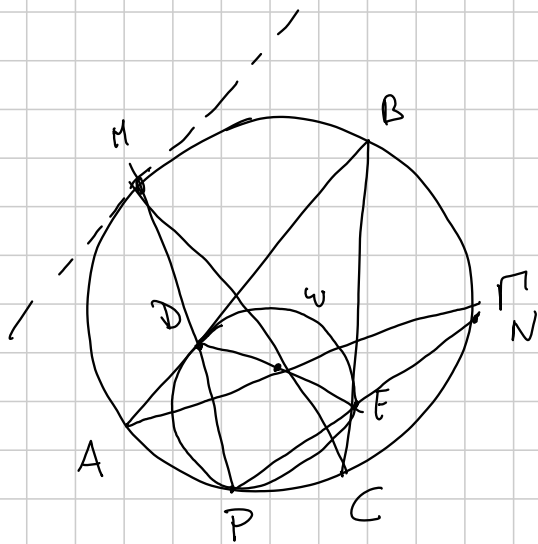


$$\widehat{AXE} = 180^\circ - \widehat{CXF}$$

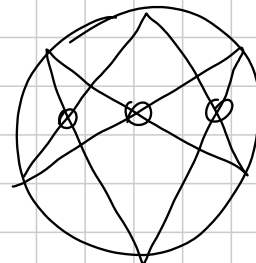
↓  $\times$  interno + ciclicita.

$$\widehat{AXB} = 180^\circ - \widehat{CXD}$$

□



Th: p.to medio di  $DE$  e' l' incentro di  $ABC$



Pappo-Pascal

Pascal ;  $D, E, CM \cap AN$  sono allineati

$M, N$  punti medi degli archi  $\Rightarrow CM$  e  $AN$  sono bisettrici in  $ABC \Rightarrow CM \cap AN \in l$

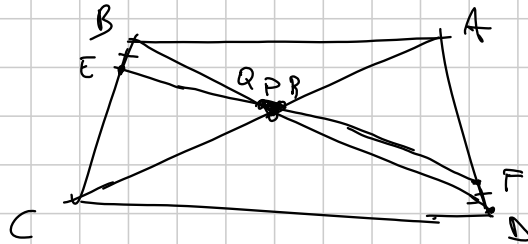
$D, I, F$  allineati. Considero  $\triangle DBF$ , isoscele  
 dunque  $I$  è pto medio di  $DF$ ,

$ABCD$  q. convesso  $BC \cong AD$  ma  $BC \nparallel AD$

$E \in BC$   $F \in AD$  t.c.  $BE = DF$ .

$AC \cap BD = P$ ,  $BD \cap EF = Q$ ,  $EF \cap AC = R$

Considero  $PQR$  al vrtine di  $E$  ed  $F$ . Dimostrare  
 che tutti i circoli hanno un punto in comune  
 diverso da  $P$





# TEORIA DEI NUMERI 3

Titolo nota

09/09/2010

## GRUPPO

$$G \times G \rightarrow G \quad f(a, b) = a \cdot b$$

$\exists$  ELEMENTO NEUTRO

$$e \in G: \quad a \cdot e = e \cdot a = a \quad \forall a \in G$$

- ASSOCIATIVA:  $(a \cdot b) \cdot c = a \cdot (b \cdot c)$

-  $\exists$  INVERSO:  $\forall a \exists b: \quad a \cdot b = b \cdot a = e$

$$b = a^{-1}$$

---

COMMUTATIVITÀ  $ab = b \cdot a$

ANELLO COMMUTATIVI

$A$  è OPERAZIONE

$A$  è un gruppo rispetto al +  
prodotto:

- ASSOCIATIVO

- DISTRIBUTIVO  $a(b+c) = ab+ac$

- COMMUTATIVO  $ab = b \cdot a$

$A$  ANELLO CON UNITÀ

se  $\exists e: \quad a \cdot e = e \cdot a = a \quad \forall a \in A$

- EL MFTPO ΣΟΠΛΑ 20
- EL MFTPO ΠΡΟΔΟΤΟ 21

ΔΟΠΙΝΙ ΔΙ ΙΝΤΕΓΡΙΤΑ'

A anello  $\mathcal{A}$  in dominio

se  $\forall a, b \neq 0, a \cdot b \neq 0$

$a$  può NON AVERE INVERSO MOLTIPLICATIVO

$$a \cdot b = a \cdot c \not\Rightarrow b = c$$

$$\left( \begin{array}{l} 2 \cdot 4 = 2 \cdot 2 \\ \mathbb{Z}_6 \end{array} \right.$$

✓ VERBA NEI ΔΟΠΙΝΙ

$$a \cdot b = a \cdot c \Rightarrow a \cdot (b - c) = 0 \Rightarrow a \cdot (b - c) = 0$$

$$a \neq 0, \forall b = c$$

ΣΑΠΡΙ  $\mathbb{K}$

- ΣΟΠΛΑ ΓΡΥΠΠΟ ΑΒΕΛΙΑΝΟ

- ΠΡΟΔΟΤΟ  $\mathbb{K} \setminus \{0\}$   $\mathbb{K}^*$

ΓΡΥΠΠΟ ΑΒΕΛΙΑΝΟ ΣΟΠΥΤΑΤΙΥΟ

$$a \cdot (b + c) = a \cdot b + a \cdot c$$

! ΣΑ ΠΡΙ ΣΟΝΟ ΔΟΠΙΝΙ.

DOMINI EUCLIDEI  $D$

ESISTE UNA FUNZIONE (GRADO)

$$g: D \setminus \{0\} \rightarrow \mathbb{N}$$

$$\forall m, n \in D \quad \exists k, r \in D:$$

$$m = kn + r$$

DIVISIONE CON RESTO

$$\text{E } g(r) < g(n)$$

OPPURE  $r=0$

$\forall a, b \neq 0$   
 $g(ab) \geq g(a)$   
 con U.C.V.A.  
 S.S.E. & INVERT.  
 ESERCIZIO

GRUPPO  $\mathbb{Z}_n$  RISPETTO ALLA SOMMA

ANELLO  $\mathbb{Z}_n, +, \cdot$  ANELLO CON UNITA

$\mathbb{Z}_n, +, \cdot$  NON È UN DOMINIO

$n | a \text{ e } n | b \Rightarrow n | a + b$   $n$  COMPONETE

PRIMO  $\neq$  IRRIDUCIBILE

$$n | ab \Rightarrow n | a \vee n | b$$

IRRIDUCIBILE SE NON SI FATTORIZZA

$$nza \cdot b \quad a = \pm 1 \vee b = \pm 1$$

$$n | ab \quad n | a \vee n | b$$

$\Downarrow \mathbb{Z}_p$  È UN DOMINIO

DOMINIO PRIMO  $\Rightarrow$  CAMPO

$\mathbb{Z}_p, +, \cdot$  è un campo,  $p$  PRIMO

FATTORIZZAZIONE UNICA (DOMINIO D)

def:  $u \in D$  è UNITA' se è INVERTIBILE  
 se  $\exists v^{-1} \quad u \cdot v^{-1} = 1$

def: DOMINIO D è A FATT. UNICA (UFD)  
 se OGNI  $d \in D$  si ESPRIME COME  
 PRODOTTO DI PRIMI IN MODO UNICO  
 (A meno di UNITA').

RISOLUZIONE DIOPHANTINA CI SERVE !!

- DOMINI EUCLIDEI SONO UFD !!

1) PRIMO  $\Leftrightarrow$  IRRIDUCIBILE

PRIMO  $\Rightarrow$  IRRIDUCIBILE

$p \in D$  PRIMO NON IRRIDUCIBILE

$p = a \cdot b$   $a, b$  NON SONO UNITA'

$p | a \Rightarrow p | a \cdot b$   $p | a$   $a = p \cdot c$

$p = p \cdot c \cdot b$   $1 = c \cdot b$

2) IRRIDUCIBILE  $\Rightarrow$  PRIMO

$d \in D$  IRRIDUCIBILE

supponiamo  $d \mid ab$   $d \nmid a$  e  $d \nmid b$

$$d \nmid a \quad \exists j, k \in D: ja + kb = 1$$

$$I = \{ na + md \mid n, m \in D \}$$

$(a, d) \in I$  L'ELEMENTO DI  $I$  A GRADO MINORE

$$a \in I \quad d \in I$$

$$\forall n \in I, c \mid n$$

$$n = ja + kd \quad j, k \in D$$

$$n = j_1 a + k_1 d$$

$$n = q \cdot c + r$$

$$g(r) < g(c)$$

$$\forall r = 0$$

$$g(r) < g(c)$$

$$j_1 a + k_1 d = q(ja + kd) + r$$

$$a(j_1 - qj) + d(k_1 - qk) = r \in I \quad g(r) < g(c)$$

ES VIEN FUORI LO STESSO E A PENO DI UNITA'

$$\begin{array}{l} a, d \in D \text{ eucl.} \quad (a, d) \\ \exists j, k \quad ja + kd \mid a \quad ja + kd \mid d \end{array}$$

$(a, d)$  SI A VALORE CHE'

$$l|a \quad \vee \quad l|d \Rightarrow l|(a, d)$$

$$(a, d) | a \quad (a, d) | d$$

$$d|al \quad (d|a \wedge d|b)$$

$$l_1 = \begin{bmatrix} j_1 & d + k_1 a \end{bmatrix}$$

$$l_1 | d \quad l_1 | a$$

$$d = l_1 \cdot c$$

$$l_1 = d \cdot c^{-1} \Rightarrow$$

$$d | l_1 \Rightarrow d | a$$

$$l_1 \text{ INVERTIBILE} \quad \exists l_1^{-1}$$

$$l_1^{-1} \begin{bmatrix} j_1 & d + k_1 a \end{bmatrix} = l_1^{-1} \cdot l_1$$

$$\begin{bmatrix} l_1^{-1} j_1 & l_1^{-1} \cdot k_1 \end{bmatrix} a = 1$$

$$d \text{ IRR.} \quad d \nmid a \quad d \nmid b$$

$$\Rightarrow \exists m_1, n_1 : m_1 d + n_1 a = 1$$

$$m_2 d + n_2 b = 1$$

$$(m_1 d + n_1 a) (m_2 d + n_2 b) = 1$$

$$m_1 m_2 d^2 + m_1 d \cdot n_2 b + n_1 a \cdot m_2 d + n_1 n_2 a b = 1$$

$$d | ab \Rightarrow d | 1 \quad \text{ASSURDO}$$

$$\text{PRIMO} \Leftrightarrow \text{IRRIDUCIBILE}$$

de  $D$  INDUO SU  $f(d)$

$$f(a) = f(b) \quad a = \prod_{i=1}^n q_i \quad q_i \in \mathbb{R}.$$

di  $a \cdot b$   $a, b$  NON UNITARI!

$$f(a) > f(b), \quad f(a) > f(b)$$

$$a = \prod_{i=1}^n q_i \quad b = \prod_{i=1}^m p_i$$

$$ab = \left( \prod_{i=1}^n q_i \right) \left( \prod_{i=1}^m p_i \right)$$

OGNI  $d$  SI FATTORIZZA IN PRIMI

$$d = \prod_{i=1}^n q_i = \prod_{i=1}^m p_i$$

$$q_n \mid p_j \quad q_n \cdot d = p_j$$

SEMPRE FALCO È HO CONTRO ESEMPLO + PICCOLO!!

$$f(ab) \geq f(a) \Leftrightarrow b \text{ INV.}$$

$$a = kab + r \quad \Rightarrow \quad f(r) \geq f(a) \quad r = 0$$

$$a = kab \quad l = kb$$

ESEMPIO

$$g(n) = |n| \quad |mn| \geq |m| \quad \checkmark$$

$$\mathbb{C}[i] = \{a + bi \mid a, b \in \mathbb{R}\}$$

$$\overline{a + bi} = a - bi \quad z \in \mathbb{C}[i] \quad \bar{\bar{z}} = z$$

$$|z| = z \cdot \bar{z}$$

$$A \xrightarrow{f} B$$

$$\text{sono } \begin{cases} f(a+b) = f(a) + f(b) & f(0) = 0 \\ f(a \cdot b) = f(a) \cdot f(b) & f(-a) = -f(a) \end{cases}$$

ISOMORFISMO = OMOMORFISMO BICETTIVO

$$A, A^{-1}$$

$$\overline{(a+bi) + (c+di)} = \overline{(a+bi)} + \overline{(c+di)}$$

$$\overline{(a+bi)(c+di)} = \overline{(ac - bd + bci + adi)}$$

$$= (ac - bd - bci - adi)$$

$$= (a - bi)(c - di)$$

---


$$\text{inoltre } (a+bi)(a-bi) = a^2 + b^2$$



$$w, z \quad |wz| = wz \cdot \overline{wz} = w \cdot z \cdot \overline{w} \cdot \overline{z}$$

$$|wz| = |w| |z|$$

$$= w\overline{w} \cdot z\overline{z}$$

$$a + bi$$

da  $\forall \alpha \in \mathbb{R}$

$$a + bi$$

$$a \neq 0 \\ b \neq 0$$

$$(a + bi) = (c + di) (\alpha + \beta i)$$

$\alpha, \beta$  non interi

$$|(a + bi) - (c + di) (\alpha + \beta i)| =$$

$$\alpha \in \mathbb{R} \quad |\alpha - A| \leq \frac{1}{2}$$

$$A + Bi$$

$$= |(a + bi) - (c + di) (\alpha + \beta i)| =$$

$$= |(a + bi) - (c + di) (\alpha + \beta i)| < |c + di|$$

$$|c + di| |(\alpha - A) + (\beta - B)i| < |c + di|$$

$$\Downarrow \\ |(\alpha - A) + (\beta - B)i| < 1$$

$$(\alpha - A)^2 + (\beta - B)^2 < \left(\frac{1}{2}\right)^2 + \left(\frac{1}{2}\right)^2 = \frac{1}{2} < 1$$

$$x^2 + 1 = y^2$$

$$x, y \in \mathbb{Z}$$

$$(x + i)(x - i) = y^2$$

$$(x + i, x - i) = (x + i, 2i)$$

$$p \in \mathbb{Z}[i] \quad |p| = q \quad \text{primo in } \mathbb{Z}$$

$$p \text{ primo} \quad p \equiv a-bi \quad |p| = |a| \cdot |b|$$

$$\} \quad \} \quad q = 1 \cdot q$$

$$z = -i(1+i)^2 \quad (1+i)(1-i) \quad \boxed{1+i \text{ PRIMO}}$$

$$p \mid n^2 + 1 = (n+i)(n-i)$$

$$a \mid b \quad \exists c \in \mathbb{D} : ac = b$$

$$\exists a, b \quad p \mid (a+bi) = n+i$$

$$p = (a+bi)(c+di)$$

$$(a+bi)(c+di) \in \mathbb{R}$$

$$c+di = k(a-bi) \quad k \in \mathbb{R}$$

$$(a+bi)(c+di) = |a+bi| \cdot k$$

$$p = (a+bi)(a-bi) \cdot n$$

$$p^2 = |a+bi| |a-bi| \cdot |n| \quad \begin{matrix} p \\ \uparrow \\ p \end{matrix}$$

$$p^2 = (|a+bi| \cdot n)^2 \quad p = |a+bi| \cdot n$$

$$p = (a+bi)(a-bi)$$

$$p \equiv 1 \pmod{4} \quad \forall a, b \in \mathbb{Z} \quad p = a^2 + b^2$$

$$p \equiv 3 \pmod{4} \quad \text{sono irriducibili} \Rightarrow p \text{ primo}$$

$$(a+bi)(a-bi) = m \cdot n \quad (a, b) = 1$$

$$\boxed{\text{WLOG}} \quad a+bi \mid m \quad \text{e } a-bi \mid n$$

$$m = (a+bi)(c+di) = l \cdot (a+bi)(a-bi)$$

$$m = l \cdot m \cdot n$$

$$(x+ni)(x-ni) = y^3 \quad x^2+1 = y^3 \quad y^3 \equiv 2 \pmod{4}$$

$$(x+ni, x-ni) = (x+ni, 2) = (x+ni, (x+ni)^2)$$

$$= (ni, 2) \quad \begin{matrix} 1 \equiv 1^3 \\ -1 \equiv (-1)^3 \\ i \equiv (i)^3 \\ -i \equiv (-i)^3 \end{matrix}$$

$$(x+ni) = (a+bi)^3$$

$$x+ni = a^3 + 3a^2bi + 3ab^2i - b^3$$

$$1 = 3a^2b - b^3 \quad b \mid 1 \quad b = \pm 1$$

$$1 = 3a^2 - 1 \quad a^2+1 = 1^3$$

$$1 = -3a^2 + 1 \Rightarrow a = 0$$

$$x^{2+1} = y^p$$

$$(x+i)(x-i)$$

$$(x+i) = (a+bi)^p$$

$$i = \sum_{k=0}^{p-1} a^{2k} (bi)^{p-2k} \cdot \binom{p}{2k}$$

$$i = \binom{p}{2} a^2 + \dots$$

$$0 = \sum_{k=1}^{p-1} a^{2k} (\pm i)^{p-2k} \binom{p}{2k} a$$

$$k \geq 1$$

$$\sqrt{2} \left( a^{2k} \binom{p}{2k} \right) \geq \sqrt{2} \left( a^2 \binom{p}{2} \right)$$

$$a^{2(k-1)} \frac{(p-2)! 2!}{(2k)! (p-2k)!} = a^{2(k-1)} \cdot \binom{p-2}{2k-2} \cdot \frac{2}{2k \cdot (2k-1)}$$

$$\sqrt{2} \left( a^{2(k-1)} \binom{p-2}{2k-2} \frac{2}{2k(2k-1)} \right) \geq 0$$

$$\sqrt{2} \left( a^{2(k-1)} \binom{p-2}{2k-2} \right) \geq \sqrt{2} (k)$$

$$\geq \sqrt{2} \left( a^{2(k-1)} \right) \geq 2(k-1)$$

$$6^{-} \text{ POSSIBILI US ARE } \sqrt{2} (k) \geq \frac{2(k-1)}{2}$$

$$x^{2+1} = y^p$$

$$(x+i)(x-i) = 1$$

$$y^p = (a+bi)^p (a-bi)^p$$

$$y^p = (a^2 + b^2)^p$$

$$p-2k > 0$$

$$b \neq 1$$

$$b = \pm 1$$

$\mathbb{Q}^{\sqrt{-2}} / \mathbb{Q}$      $K \supset \mathbb{Q}$     ASSURDO!

$\mathbb{Z}[\sqrt{-2}]$ ,  $\mathbb{Z}[\omega]$ ,  $\mathbb{Z}\left[\frac{1+\sqrt{-7}}{2}\right]$

$\mathbb{Z}[\sqrt{-5}]$      $1+\sqrt{-5}, 1-\sqrt{-5}$      $(1+\sqrt{-5})(1-\sqrt{-5})=2$

UFD, EUCLIDEO  
 PRIME    ...    1) PRIMA IN  $\mathbb{Z}$

$p \in \mathbb{Z}$  PRIMA IN  $\mathbb{Z}$      $q(x)$  NON HA RADICI mod  $p$

$\mathbb{Z}[\alpha]$      $a + b\alpha + c\alpha^2 + \dots$

$a_0 + a_1 X + \dots + a_n X^n$      $q(x)$  POL.

$\mathbb{Z}\left[\frac{1+\sqrt{-7}}{2}\right] = \mathbb{Q}(\sqrt{-7}) \cap \text{INT. ALG.}$  (ESERCIZIO)

$\mathbb{Z}[\alpha]$      $(x-\alpha)(x-\bar{\alpha})$      $|\bar{\alpha}| = \alpha - \bar{\alpha}$

$a + b\alpha \rightarrow a + b\bar{\alpha}$     ISOMORFISMO!!!

---

$x^2 + 2 = y^3$

$(x + \sqrt{-2})(x - \sqrt{-2}) = y^3$

$\sqrt{-1}, \sqrt{-2}, \frac{1+\sqrt{-3}}{2}, \frac{1+\sqrt{-7}}{2}$   
 $\frac{1+\sqrt{-13}}{2}, \frac{1+\sqrt{-17}}{2}, \frac{1+\sqrt{-19}}{2}$

$$(x + \sqrt{-2}, x - \sqrt{-2}) = (x + \sqrt{-2}, 2\sqrt{-2}) = (x + \sqrt{-2}, \sqrt{-2}^3)$$

$$\begin{aligned} x + \sqrt{-2} &= (a + b\sqrt{-2})^3 \\ &= a^3 + 3a^2b\sqrt{-2} - 6ab^2 - 2b^3\sqrt{-2} \end{aligned}$$

$$b \neq 1 \quad b = \pm 1$$

$$1 = 3a^2 - 2 \quad \begin{cases} 3 = 3a^2 \\ a = 1 \end{cases}$$

$$1 = -3a^2 + 2 \pmod{5}$$

$$a = 1 \quad b = -1 \quad -5$$

$$(1 - \sqrt{-2})^3 =$$

$$5^2 + 2 = 27$$

---


$$a^2 + 2 = 3^2$$

$$(a + \sqrt{-2})(a - \sqrt{-2}) = (1 + \sqrt{-2})^b (1 - \sqrt{-2})^b$$

$$(1 + \sqrt{-2})^b = \pm (a \pm \sqrt{-2})$$

$$\text{Immag} (1 + \sqrt{-2})^b = \pm 1$$

$$\frac{(1 + \sqrt{-2})^b - (1 - \sqrt{-2})^b}{(1 + \sqrt{-2}) - (1 - \sqrt{-2})} = \pm 1 \quad \begin{matrix} b = 1 \\ 5^2 + 2 = 3^2 \end{matrix}$$

$$(1 - \sqrt{-2})^2 = 1 - 2 - 2\sqrt{-2} = - \boxed{1 + 2\sqrt{-2}}$$

$$\uparrow 2\sqrt{-2}$$

$$(1 + \sqrt{2})^k - (1 - \sqrt{2})^k = \pm [(1 + \sqrt{2}) - (1 - \sqrt{2})]$$

$$(1 + \sqrt{2})^k \equiv \pm (-1) \pmod{1 + 2\sqrt{2}}$$

$$(-1 + 2\sqrt{2})^k \equiv 1$$

$$(-2)^k \equiv 1 \pmod{1 + 2\sqrt{2}}$$

$$1 + 2\sqrt{2} \mid (-2)^k - 1$$

$$\exists y \mid (-2)^{k-1}$$

$$a + b\sqrt{2} \mid n$$

$$\Rightarrow \frac{1a + b\sqrt{2}}{(a,b)} \mid n$$

$\exists \mid k$

$$a^2 + 2 = 3^k$$

$$\frac{(1 + \sqrt{2})^k - (1 - \sqrt{2})^k}{(1 + \sqrt{2}) - (1 - \sqrt{2})}$$

$$a_0 = 0$$

$$a_1 = 1$$

$$a_{n+2} = Ka_{n+1} + J \cdot a_n$$

$$a_0 = 0$$

$$a_1 = 1$$

$$(a_n, a_m) = a_{\min(n,m)}$$

$$m' + \text{ilc. int. } a \mid a_m$$

$$\Rightarrow \boxed{a \mid a_m \Leftrightarrow m \mid n}$$

$$a_n = 3^k$$

$$a_n = \frac{a^n - 1}{a - 1}$$

$$a_0 = 0 \quad a_1 = 1 \quad a_{n+2} = k a_{n+1} + j a_n$$

$$(k, j) \in \mathbb{Z}$$

$$R_n = \frac{z^n - \bar{z}^n}{z - \bar{z}} \quad (x-z)(x-\bar{z}) = x^2 - kx - j$$

$$\parallel \frac{z^{n+1} + z^{n-1} - jz - k}{z + \bar{z}}$$

LA SUA FORMA FISSA DAL CONIUGATO

$$m = kn$$

$$z^n \mid a_m$$

$$\bar{z}^n = \bar{z}^m$$

$$\frac{z^{kn} - \bar{z}^{kn}}{z^n - \bar{z}^n} \in \mathbb{Z}$$

$$F_n \mid F_{kn}$$

$$n = dx$$

$$m = dy$$

$$d \in (n, m)$$

$$l_{i,j} = \frac{a_{di}}{a_d} = \frac{z^{di} - \bar{z}^{di}}{z^d - \bar{z}^d} \in \mathbb{Z} \quad x^2 - cx - e = 0$$

$$(c, e) = 1$$

$$a_{(n,m)} \mid (a_n, a_m)$$

$$(z\bar{z}, z+\bar{z}) = 1$$

$$\left( \frac{a_n}{a_{(n,m)}}, \frac{a_m}{a_{(n,m)}} \right) = 1$$

$$(z\bar{z})^d, (z+\bar{z})^d \neq 1$$

$$\Downarrow$$

$$(z\bar{z}, (z+\bar{z})^d) \neq 1$$

$$b_n \quad b_0 = 0 \quad b_1 = 1$$

$$(z\bar{z}, z+\bar{z}) \neq 1 \quad \text{A. S. V. 20}$$

$$b_{n+2} = cb_{n+1} + db_n$$



$P \nmid R$  INDUZIONE;  $(Q_n, Q_{n+1}) \equiv 1$   $v = \beta x$   
 $w = \alpha y$   
 $(x, y) = 1$

$P \mid \begin{pmatrix} Q_n \\ Q_{n+1} \end{pmatrix}$   $P \mid Q_x$   $P \mid Q_y$   $P \mid Q_{xx}$   $P \mid Q_{yy}$

$Kx+1 = y$  ASSURDO!!  $P \mid Q_{xx}$   $P \mid Q_{yy+1}$

---

$X_0 = 2$   $X_{n+1} = 2X_n^2 - 1$

$(n, X_n) \equiv 1$

$\text{ord}(\sigma \cdot 2^n) = \frac{(e^\sigma)^{2^n} + (e^\sigma)^{-2^n}}{2}$

$e^\sigma = 2 + \sqrt{3}$

$X_n = \frac{(2 + \sqrt{3})^{2^n} + (2 - \sqrt{3})^{2^n}}{2}$

$P$  disp.  $P \mid X_n$   $\left(\frac{\beta}{P}\right) = -1$   $M X_n$

$\sqrt{3}$  esiste in  $\mathbb{Z}_P$   $\Downarrow$   $\frac{1}{a}$   $2^{n+2} \mid P_1$

$a^{2^n} \equiv -a^{-2^n}$   $a^{2^{n+1}} \equiv 1$

$\left(\frac{\beta}{P}\right) = -1$   $\text{ord}(a) = 2^{n+2}$

$\{a + b\sqrt{3} \mid a, b \in \mathbb{Z}_P\} = \mathbb{Z}_P[\sqrt{3}]$

$$\left| \mathbb{Z}_p \setminus \{0\} \right| = p-1$$

$$p \mid \frac{(2+\sqrt{3})^{2^n} + (2-\sqrt{3})^{2^n}}{2}$$

$$(2+\sqrt{3})^{2^n} = -(2-\sqrt{3})^{2^n} \quad \text{ovvero } (2+\sqrt{3}) = z^{n+2}$$

$$(2+\sqrt{3})^{2^{n+1}} = -1$$

$$z^{n+2} \mid p-1$$

$$z^{n+1} \mid p-1 \quad \vee \quad z^{n+1} \mid p+1$$

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$$z^n - 1 \mid X_{n-2} \Leftrightarrow z^n - 1 \text{ è PRIMO}$$


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## Equazioni tipo Pell

$$\mathbb{Z}[\sqrt{d}] = \{ a + b\sqrt{d} : a \in \mathbb{Z}, b \in \mathbb{Z} \}$$

$$z = a + b\sqrt{d} \quad \bar{z} = a - b\sqrt{d}$$

$$N(z) = a^2 - db^2 = z\bar{z} \quad (\text{moltiplicativa:})$$

$$N(z_1 z_2) = N(z_1) N(z_2)$$

$$(x + \sqrt{d}y)(a + \sqrt{d}b) = 1$$

$$x^2 - dy^2 = 1 \quad (d \text{ non quadrato})$$

$\exists$  una soluzione  $z_0 = x_0 + y_0\sqrt{d}$   $N(z_0) = 1$

con  $z_0$  minima tra le soluzioni  $N(z) = 1$

e  $z > 1$ , con la proprietà che ogni altra

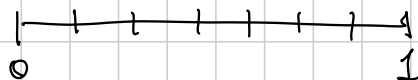
si scrive  $\pm z_0^n$ ,  $n \in \mathbb{Z}$

**Teo (Dirichlet)**  $\forall \alpha$  irrazionale,  $\exists$  infiniti

$$p/q \in \mathbb{Q} \quad \text{t.c.} \quad |\alpha - p/q| < 1/q^2$$

$$\text{Idea} \quad |\alpha - p/q| < \frac{1}{q(n+1)} \quad |nq - p| < \frac{1}{n+1}$$

$$(q \leq n) \quad 0, \{ \alpha \}, \{ 2\alpha \}, \dots, \{ n\alpha \}, 1$$



$$|\sqrt{d} - p/q| < \frac{1}{q^2}$$

$$\begin{aligned} p^2 - dq^2 &= (p + q\sqrt{d}) \underbrace{(p - q\sqrt{d})} = \\ &\leq q \cdot \frac{1}{q^2} (p + q\sqrt{d}) \leq \left(\frac{p}{q} + \sqrt{d}\right) \leq 2\sqrt{d} + 1 \end{aligned}$$

$\exists$  un  $M \leq 2\sqrt{d} + 1$  t.c.

$$x^2 - dy^2 = M \quad \leftarrow$$

ha  $\infty$  soluz.

$(p, q) \pmod{M}$  ha  $M^2$  possibilità

$$p_1 \equiv p_2 \pmod{M}, \quad q_1 \equiv q_2 \pmod{M}, \quad p_i^2 - dq_i^2 = M$$

$$N(p_1 + \sqrt{d}q_1) = M, \quad N(\dots) = M$$

$$\frac{p_1 + \sqrt{d}q_1}{p_2 + \sqrt{d}q_2} \cdot \frac{p_2 - \sqrt{d}q_2}{p_2 - \sqrt{d}q_2} = \frac{(p_1p_2 - dq_1q_2) + \sqrt{d}(q_1p_2 - q_2p_1)}{M}$$

$$\ast q_1p_2 \equiv q_2p_1 \pmod{M}$$

$$\ast p_1p_2 \equiv dq_1q_2 \pmod{M} \Leftrightarrow M = p_1^2 - dq_1^2 \equiv 0 \pmod{M}$$

$$z_0 = \min \left\{ z \text{ t.c. } N(z) = 1, z > 1 \right\}$$

Prendiamo una soluzione a  $N(z) = 1$ .

$$\text{Osserviamo che } N(z \cdot z_0) = N(z)N(z_0) = 1 \cdot 1 = 1$$

$$z_0 = x_0 + y_0 \sqrt{d} \rightarrow \bar{z}_0 = z_0^{-1}$$

$$(x_0 + y_0 \sqrt{d})(x_0 - y_0 \sqrt{d}) = 1$$

$$1 \leq z - \bar{z}_0^{-k} < z_0$$

Soluzione della Pell  $\rightarrow e^c 1$

$$z = \bar{z}_0^k$$

$$z = \pm \bar{z}_0^k$$

$$x^2 - dy^2 = m$$

Caso  $m = -1$

Ha soluz.  $\Leftrightarrow z_0$  è un quadrato in  $\mathbb{Z}[\sqrt{d}]$ .

Se  $z_0 = z_1^2$ , allora  $N(z_1) = -1$

$$\begin{cases} (N(z_1))^2 = N(z_1^2) = N(z_0) = 1 \\ z_1 < z_0 \rightarrow N(z_1) \neq 1 \end{cases}$$

Viceversa, se  $z_1$  è t.c.  $N(z_1) = -1$ , allora

$$N(z_1^2) = 1$$

Tutte le soluzioni:  $z = \pm z_1 \cdot z_0^k \quad k \in \mathbb{N}$   
 $= \pm z_1^{2k+1}$

Caso  $x^2 - dy^2 = m$  generico

Esistono al più  $\varphi(|m|)$  soluzioni,

$z_1, z_2, \dots, z_l$  con  $l \leq \varphi(|m|)$

con la proprietà che  $z$  è soluzione <sup>prim.</sup>  $\Leftrightarrow$

$$z = \pm z_k \cdot z_0^n \quad k \in \{1, \dots, l\}, \quad n \in \mathbb{Z}$$

$z$  si dice primitiva se  $(z = x + \sqrt{d}y)$  si ha

$$(x, y) = 1$$

$$\Rightarrow (z = x + \sqrt{d}y) = xy^{-1} \pmod{m}$$

$$1) \Rightarrow (z \cdot z_0^n) = \Rightarrow (z)$$

$$z \cdot z_0^n = (x + y\sqrt{d})(a + b\sqrt{d}) \quad N(a + b\sqrt{d}) = 1$$

$$= (ax + bdy) + \sqrt{d}(ay + bx)$$

$$a^2 - bd = 1$$

$$a^2 = 1 + db^2$$

$$\frac{ax + bdy}{ay + bx} \cdot \frac{ax - bdy}{ax - bdy} =$$

$$= \frac{a^2 x^2 - b^2 d^2 y^2}{\underbrace{a^2 xy + abx^2 - abdy^2 - b^2 dxy}_{xy}} = \frac{db^2(x^2 - dy^2) + x^2}{xy + Nab}$$

$$= \frac{Ndb^2 + x^2}{Nab + xy} \equiv \frac{x}{y} \pmod{N}$$

$$\text{Se } \mathfrak{s}(\omega_1) = \mathfrak{s}(\omega_2), \quad \frac{\omega_1}{\omega_2} \in \mathcal{K}[\sqrt{d}]$$

$$N\left(\frac{\omega_1}{\omega_2}\right) = \frac{N(\omega_1)}{N(\omega_2)} = 1 \quad \longrightarrow \quad \omega_2 = \omega_1 \cdot z_0^n$$

*Remark* Se  $z \in \mathcal{K}[\sqrt{d}]$ ,  $N(z) = m$ ,

$$x = \frac{z + \bar{z}}{2}, \quad y = \frac{z - \bar{z}}{2\sqrt{d}}$$

$z = z_0^n \quad \longrightarrow \quad x, y$  rispettano relaz. per ricorrenza

$$x_0 = 1, \quad y_0 = 0$$

$$x_{n+2} = (z_0 + \bar{z}_0) x_{n+1} - (z_0 \bar{z}_0) x_n$$

## Bound sulle soluzioni

Se  $x^2 - dy^2 = m$  ha una soluzione, allora ne ha una in cui  $|x| \leq \frac{1+z_0}{2\sqrt{z_0}} \sqrt{|m|}$

Moltiplicando per  $z_0^n$  opportuno, trovate una soluzione  $\in \left[ \sqrt{\frac{m}{z_0}}, \sqrt{m \cdot z_0} \right)$  che chiamo  $w$

$$2|x| = |w + \bar{w}| = \left| w + \frac{m}{w} \right| \leq$$

$$\leq \max_{w \in \left[ \sqrt{\frac{m}{z_0}}, \sqrt{m \cdot z_0} \right]} \left| w + \frac{m}{w} \right| \quad (m > 0)$$

$$= \sqrt{|m|z_0} + \sqrt{\frac{|m|}{z_0}} = \sqrt{|m|} \frac{z_0 + 1}{\sqrt{z_0}}$$



Trovare  $z_0$  + successioni di Farey  
FAREY

$$\mathcal{F}_n = \left\{ \frac{a}{b} : (a,b)=1, 0 < b \leq n, 0 \leq \frac{a}{b} \leq 1 \right\}$$

$$\begin{array}{l} | \quad \frac{0}{1} \quad \frac{1}{1} \\ | \quad \frac{0}{1} \quad \frac{1}{2} \quad \frac{1}{1} \\ | \quad \frac{0}{1} \quad \frac{1}{3} \quad \frac{1}{2} \quad \frac{2}{3} \quad \frac{1}{1} \end{array}$$

Algoritmo:  $\ast \frac{n}{1} < \sqrt{D} < \frac{n+1}{1}$

$\ast$  Ad ogni passo, se  $\frac{a}{b} < \sqrt{D} < \frac{c}{d}$ ,

calcolate  $\frac{a+c}{b+d}$ .

• Se  $x=(a+c), y=(b+d)$  e' soluzione Pell  $\rightarrow$  FINE

• Altrimenti,  $\frac{a}{b} < \sqrt{D} < \frac{a+c}{b+d}$  o

$$\frac{a+c}{b+d} < \sqrt{D} < \frac{c}{d}$$

Lavoriamo in  $\mathcal{F}_n$ . Diciamo che  $\frac{a}{b} < \frac{c}{d}$  siano

consecutive. Allora

i)  $b+d > n$

iii)  $bc - ad = 1$

ii)  $(a, c) = (b, d) = 1$

$$i) \frac{a}{b} < \underbrace{\frac{a+c}{b+d}}_{\leq n} < \frac{c}{d} \quad \text{ASSURDO}$$

ii) Segue da iii + Bézout

iii) Induzione su  $n$ .  $\frac{a}{b} < \frac{c}{d}$  in  $\mathcal{F}_n$

\* o  $\frac{a}{b}, \frac{c}{d}$  restano consecutive in  $\mathcal{F}_{n+1}$

$$* \text{ o } \frac{a}{b} < \frac{M}{n+1} < \frac{c}{d}$$

$$\frac{M}{n+1} - \frac{a}{b} = \frac{k_1}{b(n+1)}, \quad \frac{c}{d} - \frac{M}{n+1} = \frac{k_2}{d(n+1)}$$

$$\frac{1}{bd} \stackrel{ii}{=} \frac{c}{d} - \frac{a}{b} = \frac{k_1 d + k_2 b}{bd(n+1)}$$

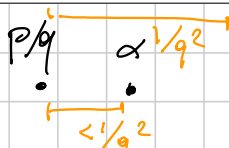
$$\underbrace{-b-d}_{\leq 0} + n + 1 = k_1 d + k_2 b - b - d = d(k_1 - 1) + b(k_2 - 1) \rightarrow k_1 = k_2 = 1$$

Teo (Hurwitz)  $\alpha$  irraz.

$$\exists \infty \frac{p}{q} \text{ t.c. } \left| \alpha - \frac{p}{q} \right| < \frac{1}{\sqrt{5} q^2}$$

$$\frac{a}{b} < \alpha < \frac{c}{d} \quad \text{una tra } \frac{a}{b}, \frac{c}{d}, \frac{a+c}{b+d} \text{ va}$$

bene nel Teo. Hurwitz



$$\alpha - p/q < 1/q^2$$

Esisterà  $a/b$  t.c.  $\alpha > a/b > p/q$ , e  $b \leq q$ ?

$$\frac{a}{b} - p/q \geq \frac{1}{bq} \geq \frac{1}{q^2}$$

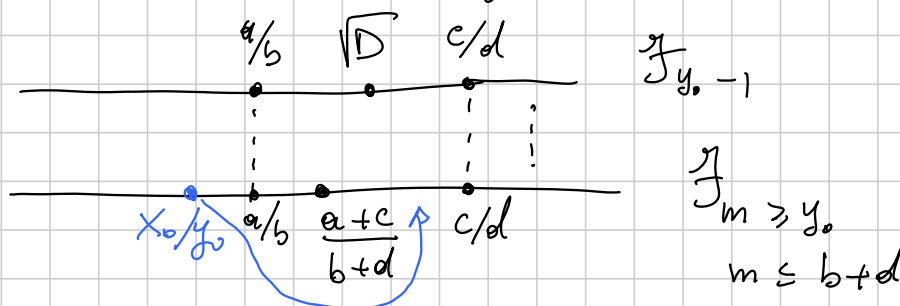
Se  $x + y\sqrt{d}$  risolve la Pell,  $x^2 - dy^2 = 1$

$$x/y > \sqrt{d}$$

$$\frac{x}{y} - \sqrt{d} = \frac{1}{y} \frac{(x - y\sqrt{d})(x + y\sqrt{d})}{(x + y\sqrt{d})} = \frac{1}{y^2(\frac{x}{y} + \sqrt{d})} < \frac{1}{2y^2\sqrt{d}} < \frac{1}{2y^2}$$

Sia  $x_0, y_0$  la più piccola soluz. Pell;  $x_0/y_0$

starà in una  $J_m = J_{y_0}$ .



Sappiamo  $|x_0/y_0 - \sqrt{d}| < \frac{1}{2y_0^2}$ . Se  $\frac{x_0}{y_0} \notin (\frac{a}{b}, \frac{c}{d})$

succede che  $a/b$  è un'approx di  $\sqrt{d}$  migliore di  $x_0/y_0$  (IMPOSSIBILE, perché  $b \leq y_0$ )

$$\frac{c}{d} - \frac{a+c}{b+d} = \frac{cb + \cancel{cd} - ad - \cancel{cd}}{d(b+d)} = \frac{1}{d(b+d)}$$

Esempio  $a^2 = 19n+1$  e  $b^2 = 95n+1$  sono quadrati:

$$b^2 = 5a^2 - 4 \quad b^2 - 5a^2 = -4$$

$$\varphi(141) = 2 \quad b \equiv a \pmod{4}$$

$$b \equiv 3a \pmod{4}$$

$$(1 + \sqrt{5} \cdot 1)(1 - \sqrt{5} \cdot 1) = -4$$

$$(11 + 5\sqrt{5})(11 - 5\sqrt{5}) = -4$$

Mancano: \* non primitive  $x^2 - 5y^2 = -1$

$$* z_0 \quad x^2 - 5y^2 = 1 \quad \frac{2}{1} < \sqrt{5} < \frac{3}{1}$$

$$\frac{2}{1} < \sqrt{5} < \frac{5}{2} \quad \frac{7}{3} > \sqrt{5}$$

$$\frac{2}{1} < \sqrt{5} < \frac{7}{3}$$

$$\frac{9}{4}$$

Non primitive:  $(2 + \sqrt{5})^{2k+1}$

$$(1 + \sqrt{5})(9 + 4\sqrt{5})^k, \quad (11 + 5\sqrt{5})(8 + 4\sqrt{5})^k$$

$$(2 + \sqrt{5})^{2k+1}$$

$$\varphi^2 = \varphi + 1$$

$$(1 + \sqrt{5})^2 = 6 + 2\sqrt{5}$$

$$(1 + \sqrt{5})^3 = 16 + 8\sqrt{5} = 8(2 + \sqrt{5})$$

$$(2 + \sqrt{5})^2 = 9 + 4\sqrt{5} = \left(\frac{1 + \sqrt{5}}{2}\right)^6$$

$$(1 + \sqrt{5})^5 = 16 (3 + \sqrt{5})(2 + \sqrt{5}) = 16 (11 + 5\sqrt{5})$$

$$11 + 5\sqrt{5} = 2 \left( \frac{1 + \sqrt{5}}{2} \right)^5$$

$$z = 2 \left( \frac{1 + \sqrt{5}}{2} \right)^{2k+1}$$

$$y = \frac{z \left( \frac{1 + \sqrt{5}}{2} \right)^{2k+1} - z \left( \frac{1 - \sqrt{5}}{2} \right)^{2k+1}}{2\sqrt{5}} = F_{2k+1}$$

$$F_{2k+1}^2 - 1 \equiv 0 \pmod{19}$$

$$\left( \frac{5}{19} \right) = \left( \frac{19}{5} \right) = \left( \frac{4}{5} \right) = 1$$

$$\text{Es} \quad 3^m - 2 = x^2$$

$x$  dispari  $\rightarrow n$  dispari

$$3y^2 - 2 = x^2$$

$$x^2 - 3y^2 = -2$$

$$(1 + \sqrt{3})$$

$$x^2 - 3y^2 = 1$$

$$(2 + \sqrt{3})$$

$$y_n = \frac{(1 + \sqrt{3})(2 + \sqrt{3})^n - (1 - \sqrt{3})(2 - \sqrt{3})^n}{2\sqrt{3}} = \textcircled{*}$$

$$(1 + \sqrt{3})^2 = 4 + 2\sqrt{3} = 2(2 + \sqrt{3})$$

$$\textcircled{*} \frac{(1 + \sqrt{3})^{2n+1} - (1 - \sqrt{3})^{2n+1}}{2^{n+1} \sqrt{3}}$$

$$9 | y_n \Leftrightarrow n \equiv 4 \pmod{9}$$

$$y_4 = 153 = 9 \cdot 17$$

$$y_{2m+1} = 2^m y_m = \frac{(1 + \sqrt{3})^{2m+1} - (1 - \sqrt{3})^{2m+1}}{2\sqrt{3}}$$

$$(b_{2n+1}, b_g) = b_{(g, 2n+1)} = b_g \equiv 0 \pmod{17}$$

$$17 \mid 2^m y_n \quad 17 \mid y_n$$