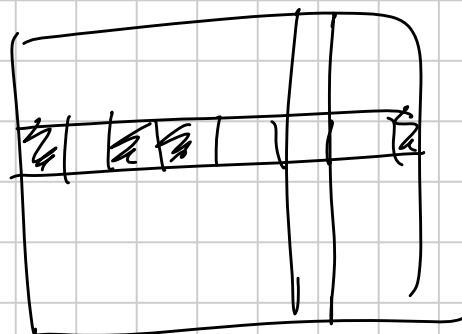
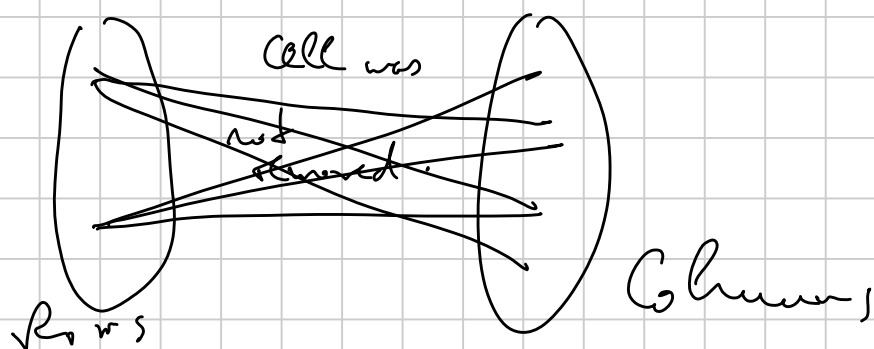


Problem 1. A $2n \times 2n$ chessboard; remove k cells, but not more than n on any row or column.

Prove that one can still place $2n$ rooks so that they do not attack each other



Solution - If no cells removed $(2n)!$ ways.



Answer = perfect matching

Marriage Theorem. Bipartite graph $|A|=|B|$,

if for any $S \subseteq A$, $|S| \leq |N(S)|$.

\implies \exists a perfect matching.

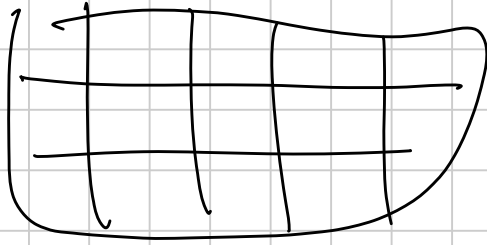
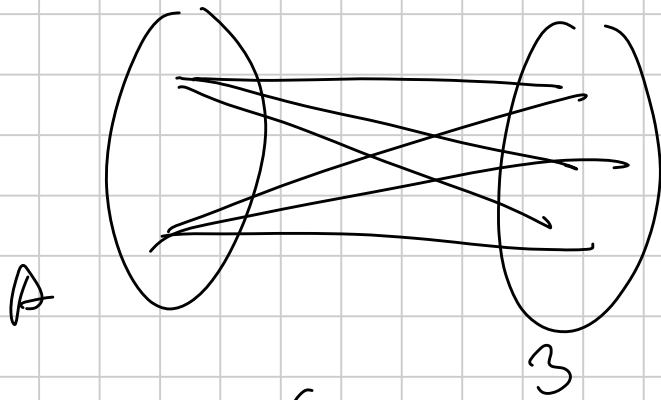
If $|S| \leq n$, ok: we are row $\in S$

If $|S| > n$, $|N(S)| = 2n$

By M.T. \implies Perfect matching can be done.

Particular case

k -regular



Set (finite) = n elements
 π partition into m classes
 of k elements each.

Find n elements which represent both partitions.

G group, H subgroup

$$G = \bigcup_{i=1}^r Hx_i = \bigcup_{i=1}^r y_i H$$

index x_i index y_i
 z_i τ_i

Problem 2. (A.R.R) Island A , n families.

Ministry Agriculture : partitioned A into n equal

Ministry Hunting : \dots

Ministry Family : we should allocate to each family one each so that they interested

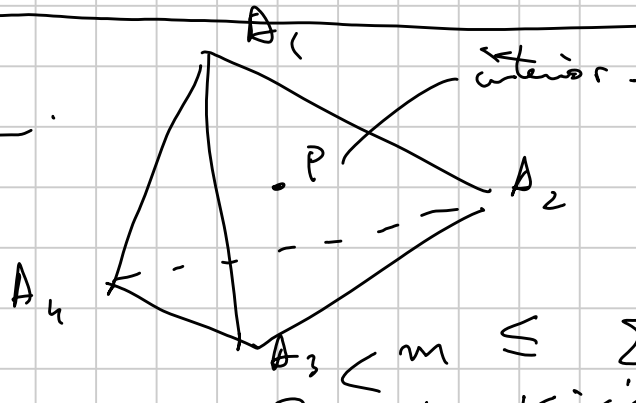
They succeeded!

Ministry Religion : declared a miracle!

a) Math. Institute : \rightarrow reported: it could always be done!

b) Find the fraction λ of area A inside G that by a dense allocation, each family receives a $|C_i \cap H_i| \geq \lambda$.

GEOMETRY



$$0 < m \leq \sum_{1 \leq i < j \leq 4} \frac{A_i \cdot PA_j}{A_j} \leq M$$

Claim
 $3\bar{u}$
 $4\bar{u}$

\therefore $\Sigma \approx 3\bar{u} + 3 \cdot 0 = 3\bar{u}$

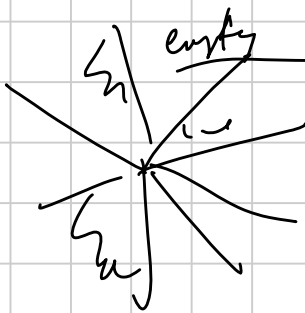
\therefore $\Sigma \approx 2 \cdot 0 + 4 \cdot \bar{u} = 4\bar{u}$



Convex polyhedron

$\Sigma \text{ angles} < 2\pi$

"cut" along edges \rightarrow put on a plane



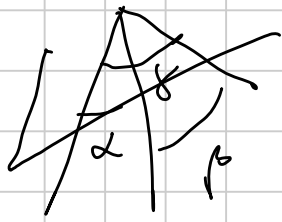
$2\pi - \Sigma \text{ angles} = \text{"defect"}$

$\Sigma \text{ defects} = \boxed{4\pi}$
Polyhedron

Euler: $\underline{V + F = E + 2}$

Proof. \sum angles $< 2\pi$
 around a vertex

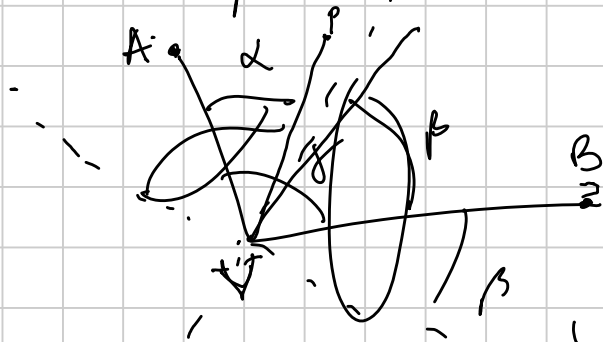
show Lemma. Vertex of a ~~hexagon~~ hedron.



α, β, γ have the "triangular" ineq. property

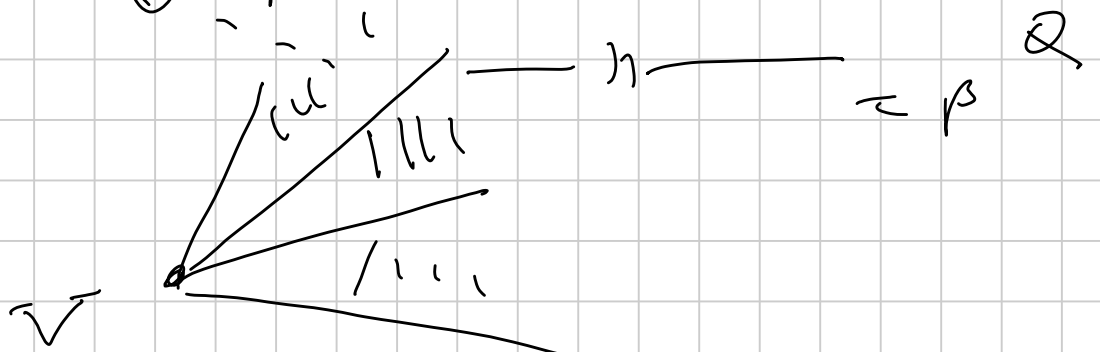
$$\begin{cases} \alpha + \beta \geq \gamma \\ \beta + \gamma \geq \alpha \\ \gamma + \alpha \geq \beta \end{cases}$$

In the plane draw the locus of γ :

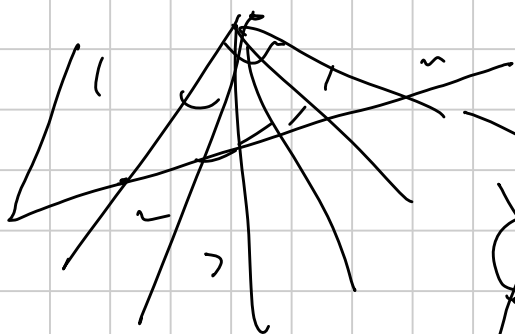


Assume $\alpha + \beta < \gamma$.

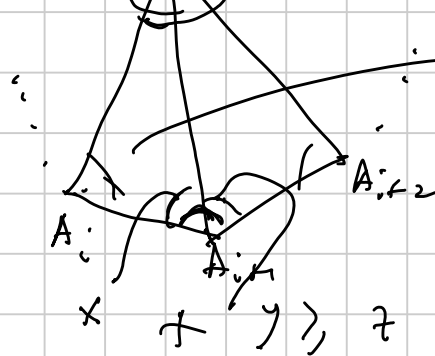
Locus in \mathbb{R}^3 of points P
 \hookrightarrow that $\angle PVA = \alpha$?



$$\sum_{i=1}^n \alpha_i$$



a pyramid.



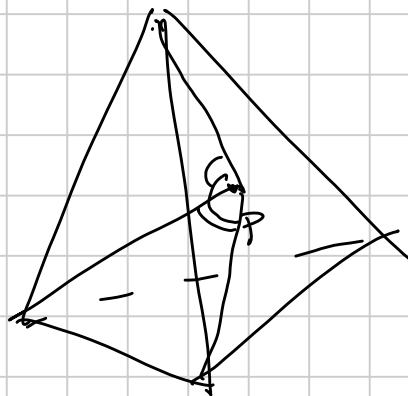
\sum along all lateral faces $\approx n\pi$

$$\sum \alpha_i \approx (n-2)\pi$$

$$\sum_{i,j}^6 \mathbb{1}_{A_i P A_j} \leq 4\pi$$

$$\sum \leq 2\pi$$

for all groups of 3 out of 4 vertices



Four times $\frac{2\sum}{4} \leq \frac{4\pi}{4}$

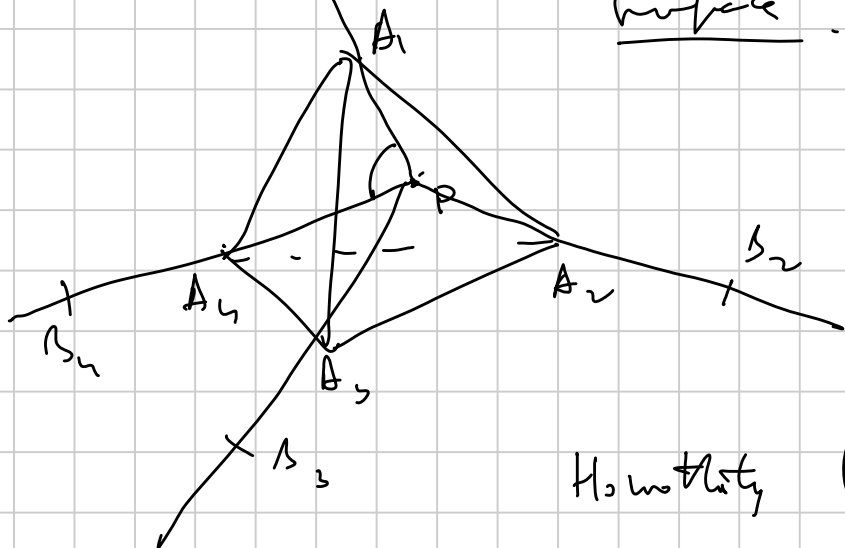
$$3 \leq \sum \leq \widehat{A_1 P A_3} + \widehat{A_3 P A_4} + \widehat{A_4 P A_2} + \widehat{A_2 P A_1}$$

(1, 3, 4, 2)

Lateral. Why is this happening (is it?)

One way → you do a nice drawing on a sphere

"geodesic" = "arc" of minimal length that connects 2 points on a surface

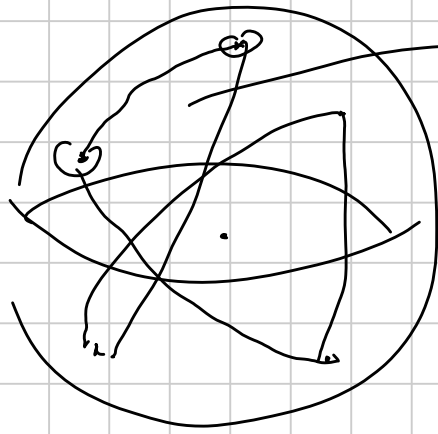
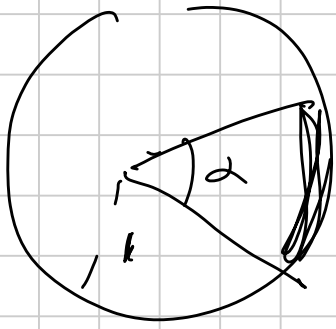


Can take B_i such that $P B_i = \text{rad.}$

∩ center of the sphere $\triangle B_1 B_2 B_3 B_4$ inside Δ

Homothety $\Rightarrow 1$

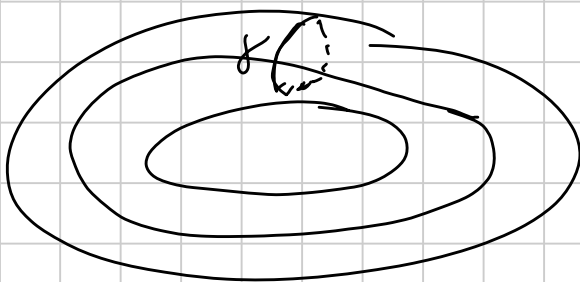
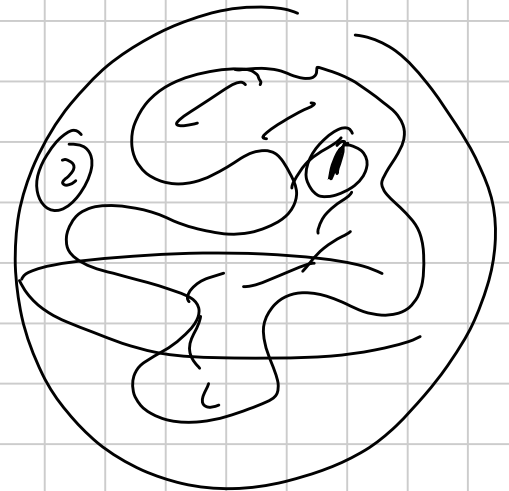
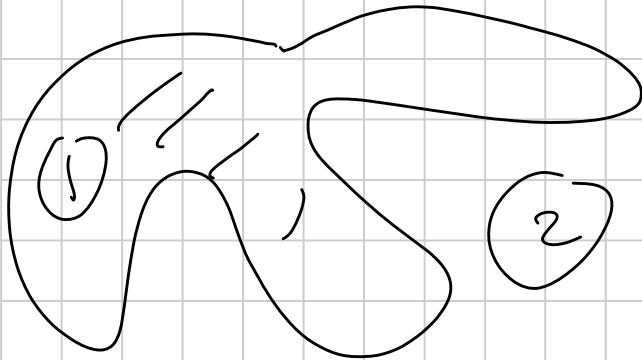
The advantage: measure of \angle = length of the arc great circle is subtended



length of γ $\geq 2\pi r$.

WRONG PICTURES

Combi prob 1. If on surface of sphere of radius r , curve γ simple (does not auto-intersect), closed, rectifiable (its length can be calculated). And if and a curve, divides surface of sphere into 2 regions, equal area, Then length $\lambda(\gamma) \geq 2\pi r$.

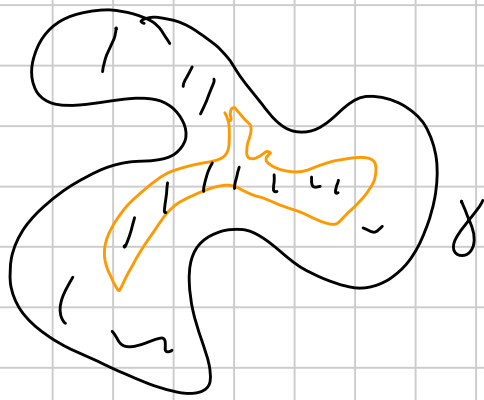


Torus (doughnut)

Proof. $\varphi: S^2 \rightarrow S^2$
 $P \rightarrow P' = \text{antipode}$

$\varphi(\gamma) = \underline{\text{curve}}$.

γ blue
 $\varphi(\gamma)$ red

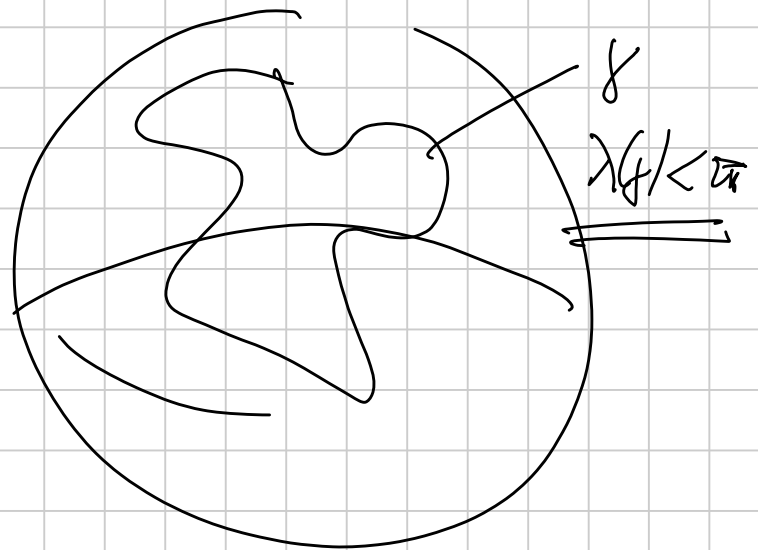
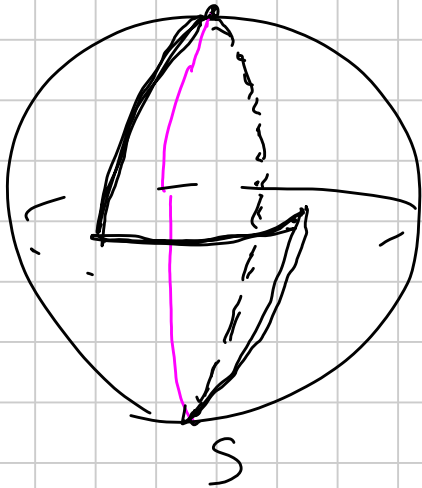


$$|\gamma \cap \mathcal{C}(P)| \geq \frac{2(\text{pairs})}{P}$$

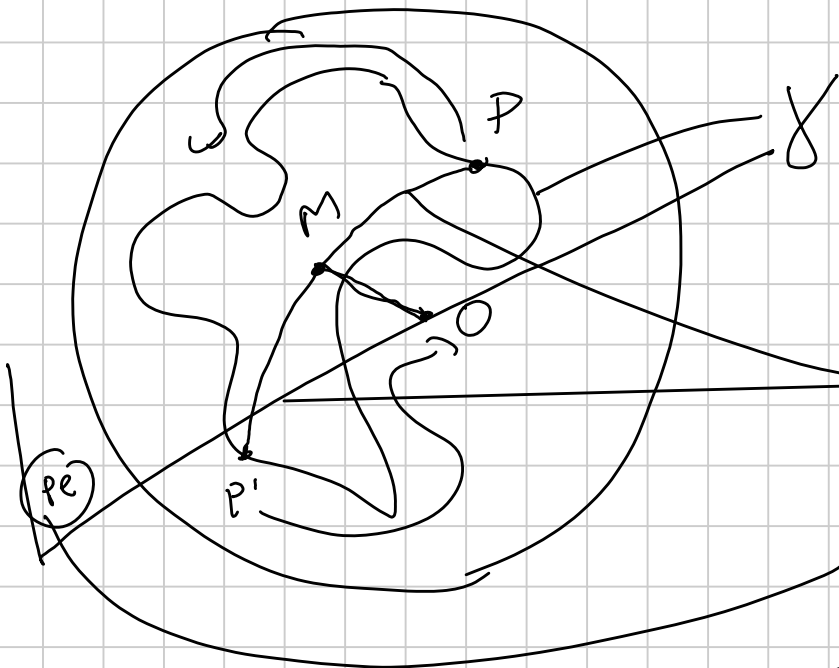
$$P, P' \in \gamma \quad (\mathcal{C}(P))$$

$$\lambda(\gamma) = \widehat{PP'} + \widehat{P'P} \geq \frac{1}{2} \text{circle} + \frac{1}{2} \text{circle} = \text{full great circle} = \pi$$

□



$\gamma \subset \text{Hemisphere}$



$$\exists! P'$$

$$\lambda(PP') = \lambda(P'P)$$

arc of great circle \$PP'\$

plane \perp OP'
" " " " middle

$\gamma \subset \text{Hemipl.}$

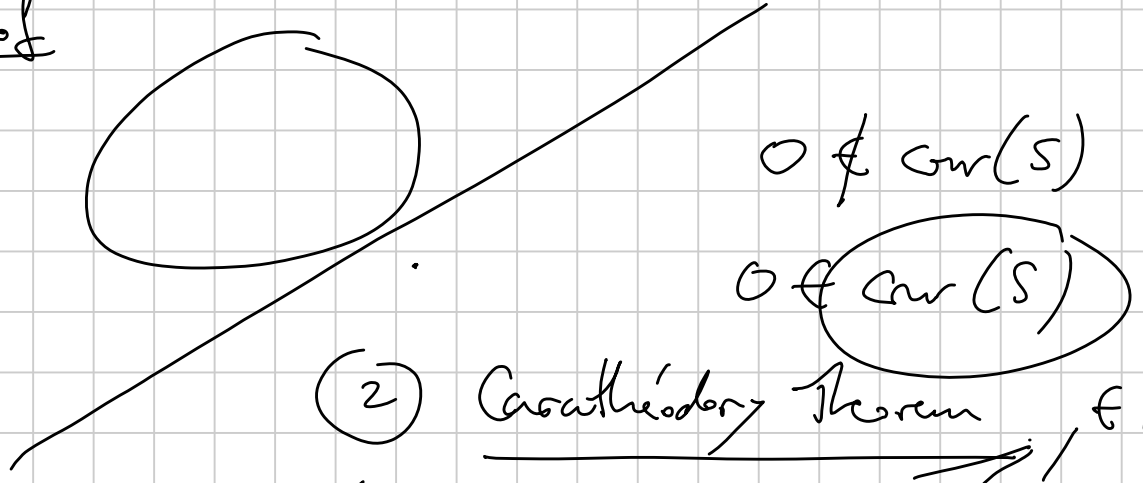
Either (PP') $PP' \cap$ "equator" Assume not

equator $\cap P' \rightarrow$ symmetric with respect to OP
 $\circ \subset S$
 same length $= \frac{1}{2} \lambda(y)$
 x' also on equator (opposite)

Instead looking at γ : $(PP' \cup \text{symmetric arc})$
 contains $x, x' = \lambda(y)$

My solution : (1) The fact that a set $S \subset S^2$ is not contained in ANY hemisphere $\Leftrightarrow \text{Conv}(S) \ni 0$.

Proof



(2) Carathéodory Theorem $\in \mathbb{R}^n$
 There exist $k \leq n$ such $A_1, \dots, A_k \ni 0$

(3)

Monty Python : And now, on a different...
 problems with hate, names ...



Establish an equivalence relation among
infinite binary words?

wwwwww... $\xrightarrow{\quad}$
 $X \sim Y$ def they differ only in finitely
many positions.

Agree on a representative $X \in \hat{X}$ (Axiom of
choice)
Each one \rightarrow write ? the color on the
position that he skipped.

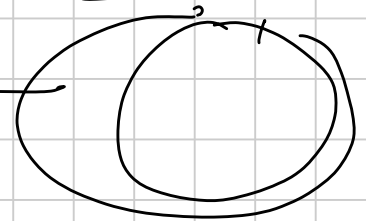
Finite variations $1 \rightarrow 1 \dots n$
 \downarrow 1, 2, ..., n

The n persons again
not allowed to speak

n — a they
write down
a color, they
also pronounce it.

But instead $\xrightarrow{\quad}$

(n) people, (n) colors of hats



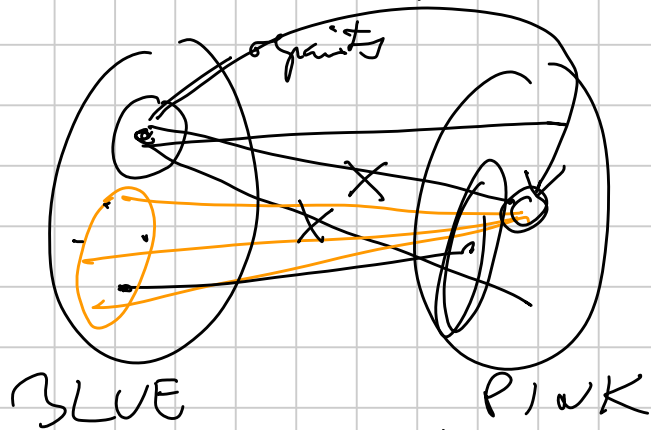
Strategy 1st give themselves some number $1 \dots n$
2nd like the color some color $1 \dots n$

Each will write: $\left[\left(\sum \text{color he sees} \right) - \text{his position} \right] \pmod{n}$
 $= \underbrace{\sum}_{\text{constant}} \text{all color} - \text{the color that person has}$

In a forest there n dwarves $\begin{cases} \text{blue house} \\ \text{pink} \end{cases}$

Each one, in turn, visits all his friends.
 If strictly $> 1/2$ his friends have a lower
 different color than his \rightarrow he repaints.
 Show, after some time \rightarrow there is no need
 to repaint.

Sikinia.



The number of edges between B/P is \downarrow strictly

The 2m dwarves — game:

In cabin there are $2m$ boxes, each
 containing the name of one dwarf.

One by one, each enters cabin, allowed to
 open $\leq m$ boxes \rightarrow finding his name.

They win if ALL find their names.
 What strategy to maximize the chance to win?

From the "idiot" $\rightarrow \left(\frac{1}{2}\right)^{2m}$ chance
 a $\frac{1}{m}$ chance by improvement

Sketch: $k \rightarrow$ opens box k — if D
 if not opens box \leftarrow
 $\sigma(1) \sigma(2) \dots \sigma(2m)$

$$\sigma = (1, \sigma(1), \sigma(\sigma(1)) \dots) (3, \sigma(3) \dots) \dots (k) =$$

$$\Rightarrow \sigma = \prod_{i=1}^n f_i, \quad \sum \lambda(f_i) = 2n.$$

They are finite if $\lambda(f_i) \leq n, \forall i$.

Estimate $\frac{\# \sigma \text{ with cycles } \leq n}{\# \text{ perm}} = \frac{1}{(2n)!}$.

Count N of σ which do have cycles of $\lambda > n$

Let's count $\# M_k$ of σ with a cycle $\lambda = n+k$
 $k=1, 2, \dots, n$

$$\binom{2n}{n+k} \cdot (n+k-1)! \cdot (n-k)! = \frac{2n!}{n+k}$$

$$\binom{2n}{n+k} \cdot (n+k-1)! \cdot (n-k)! = \frac{2n!}{n+k}$$

$$p = \frac{N}{(2n)!} \approx \frac{1}{n+1} + \frac{1}{n+2} + \dots + \frac{1}{2n} \approx \ln 2$$

$$p \approx 1 - \ln 2 \approx 0.7 \approx 30\%$$

LAST → (Martin GARDNER)

You are given N folded papers, each a

You pick the largest

real number
distinct.

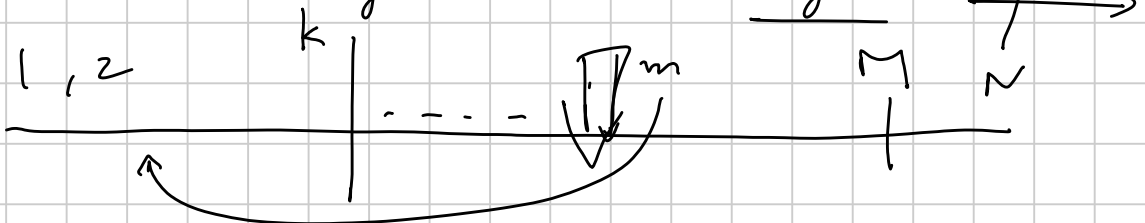
The "idiot" — pick one $\frac{1}{N}$

the only possible strategy is:

— pick $1 \leq k \leq N-1$

— open k papers — put them aside.

— start opening, until hit one
larger than those rejected



$k \rightarrow p_k =$ probability of winning.
find which k max

$$k \approx \frac{N}{e}, \quad p_k \approx \frac{1}{e}$$

1...2...3, SCAPPATE!!! XD

MA MEGLIO MALE CHE CI SONO I CANNONI

"PUÒ BACIARE LA SPOSA!"

"MA È UN LUPO!!!"

"NO! PRENDIAMO X E MO' NON LA FA LA PROLETTINIA
CHE È MANDA P ALL' INFINITO!"