

COMBINATORIA 2^{BASIC}

Titolo nota

09/09/2011

$$\{1, 2, \dots, n\} \xrightarrow{\sigma} \{1, 2, \dots, n\}$$

Biiettiva

| | | | | | | | | | |
|------------|---|---|---|---|---|---|---|---|---|
| λ | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | |
| σ : | 2 | 6 | 3 | 5 | 7 | 4 | 1 | 9 | 8 |

Tutte le permutazioni su n elementi
 $n!$

$$\sigma: \begin{pmatrix} 1 & 2 & 3 & 4 & 5 \\ 5 & 2 & 1 & 4 & 3 \end{pmatrix}$$

$$\tau: \begin{pmatrix} 1 & 2 & 3 & 4 & 5 \\ 2 & 1 & 3 & 5 & 4 \end{pmatrix}$$

$$\sigma \circ \tau: \begin{pmatrix} 1 & 2 & 3 & 4 & 5 \\ 5 & 2 & 1 & 3 & 4 \end{pmatrix}$$

$$\tau \circ \sigma: \begin{pmatrix} 1 & 2 & 3 & 4 & 5 \\ 4 & 1 & 2 & 5 & 3 \end{pmatrix}$$

→ (1 5 3)

$$\begin{matrix} 1 & 2 & 3 & 4 & 5 \\ 1 & 3 & 2 & 4 & 5 \end{matrix}$$

$$\in \{1, \dots, n\}$$

$$\sigma = (a_1 \ a_2 \ \dots \ a_k)$$

$\sigma(a_1) = a_2$
 $\sigma(a_2) = a_3$
 \dots
 $\sigma(a_{k-1}) = a_k$
 $\sigma(a_k) = a_1$

TUTTI GLI
ALTRI
FISSI

$$(1 \ 2 \ 3) \quad (4 \ 5)$$

| | | | | |
|---|---|---|---|---|
| 1 | 2 | 3 | 4 | 5 |
| 2 | 3 | 1 | 5 | 4 |

$$a_1 \rightarrow a_2 \xrightarrow{\sigma} a_3 \rightarrow \dots$$

$$a_k = a_h \quad h < k \quad \exists \text{ cassetta}$$

$\exists k$ più piccolo con questa proprietà minimo

$$a_k = a_1 \quad h=1$$

$$a_k = a_h$$

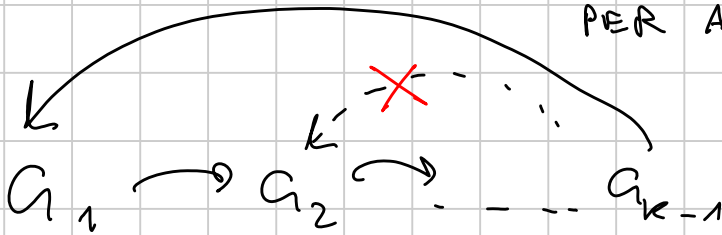
$$k > h > 1$$

PER ASSURDO

$$a_k = \sigma(a_{k-1})$$

$$a_h = \sigma(a_{h-1})$$

$$a_{k-1} = a_{h-1}$$



$$\sigma = (a_1 a_2 \dots a_{k-1}) \sigma'$$

| | | | | | | | | |
|---|---|---|---|---|---|---|---|---|
| 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 |
| 2 | 6 | 3 | 5 | 7 | 4 | 1 | 9 | 8 |

$$(1 \ 2 \ 6 \ 4 \ 5 \ 7) (8 \ 9)$$

segno

$$\text{sgn}(\sigma) = \prod_{1 \leq i < j \leq n} \frac{\sigma(i) - \sigma(j)}{i - j} \in \{ \pm 1 \}$$

$$\sigma: \{1 \dots n\} \rightarrow \{1 \dots n\}$$

$$\text{sgn}(\sigma \circ \tau) = \text{sgn}(\sigma) \cdot \text{sgn}(\tau)$$

$$\text{sgn}(\sigma \circ \tau) = \prod_{1 \leq i < j \leq n} \frac{\sigma \circ \tau(i) - \sigma \circ \tau(j)}{i - j} =$$

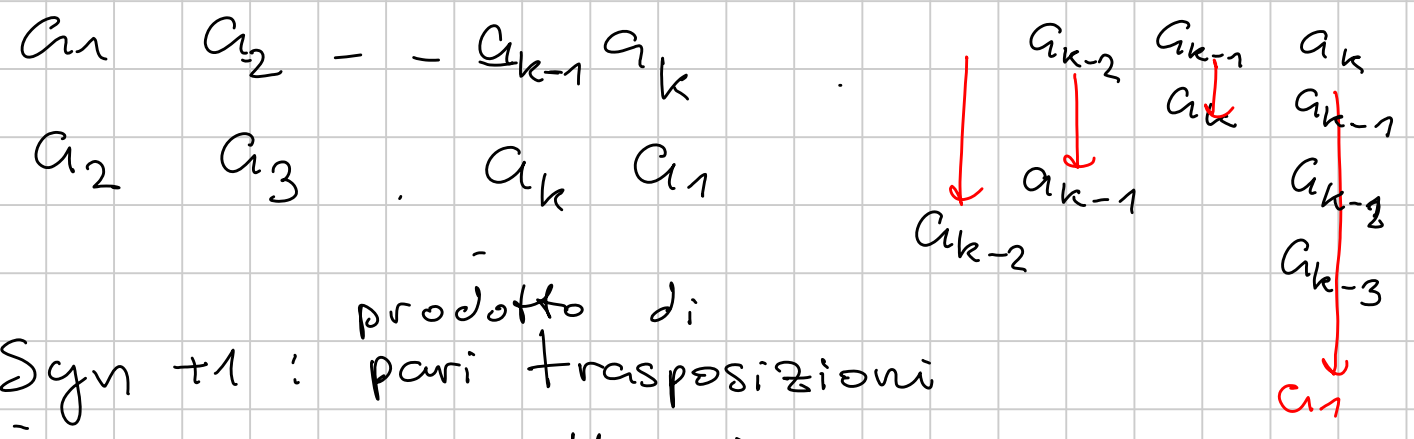
$$\prod_{1 \leq i < j \leq n} \frac{\sigma(\tau(i)) - \sigma(\tau(j))}{\tau(i) - \tau(j)} \frac{\tau(i) - \tau(j)}{i - j}$$

$\underbrace{\hspace{150px}}_{\text{sgn } \sigma}$
 $\underbrace{\hspace{150px}}_{\text{sgn } \tau}$

$$\begin{aligned} \tau(i) &= k & k > l \\ \tau(j) &= l \\ k < l & \quad \frac{\sigma(k) - \sigma(l)}{k - l} = \frac{\sigma(l) - \sigma(k)}{l - k} \end{aligned}$$

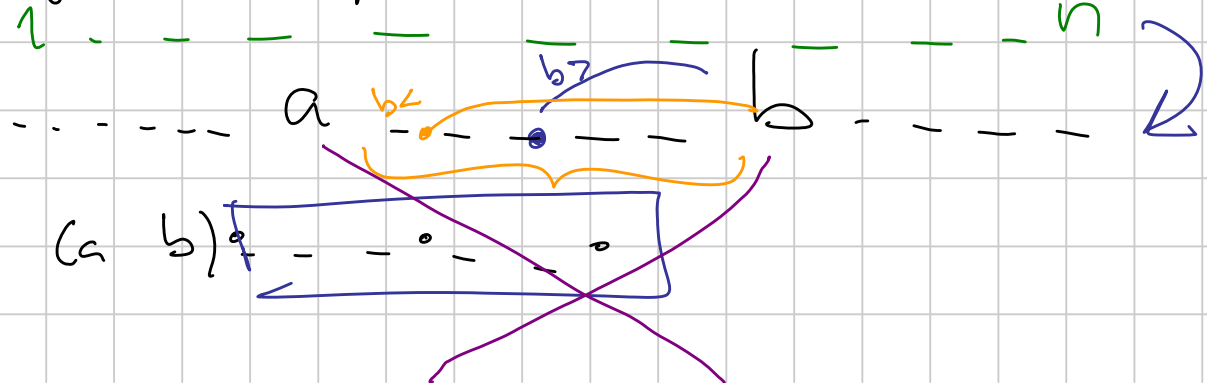
Il segno di σ conta la parità del numero di coppie (i, j) con $1 \leq i < j \leq n$ tali che $i < j$ ma $\sigma(i) > \sigma(j)$

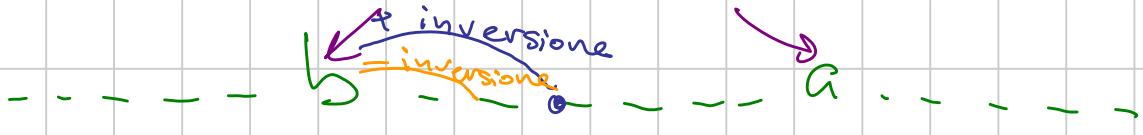
$$(a_1 \ a_2 \ \dots \ a_k) = (a_1 \ a_2) (a_2 \ a_3) \dots (a_{k-1} \ a_k)$$



pari $\text{sgn} + 1$: prodotto di pari trasposizioni
 dispari $\text{sgn} - 1$: prodotto di dispari trasposizioni

Segno parità del numero di inversioni



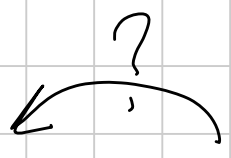


parità del numero di inversioni
 parità = del numero di trasposizioni

$$(a_1 a_2 \dots a_k) = (a_1 a_2) (a_2 a_3) \dots (a_{k-1} a_k)$$

un ciclo di lunghezza pari è
 una permutazione dispari
 dispari è
 pari.

| | | | |
|----|----|----|----|
| 1 | 2 | 3 | 4 |
| 5 | 6 | 7 | 8 |
| 9 | 10 | 11 | 12 |
| 13 | 14 | 15 | |

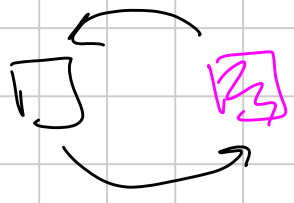


| | | | |
|----|----|----|----|
| 1 | 2 | 3 | 4 |
| 8 | 6 | 7 | 8 |
| 9 | 10 | 11 | 12 |
| 13 | 15 | 14 | |

NO

MOSSA = TRASPOSIZIONE

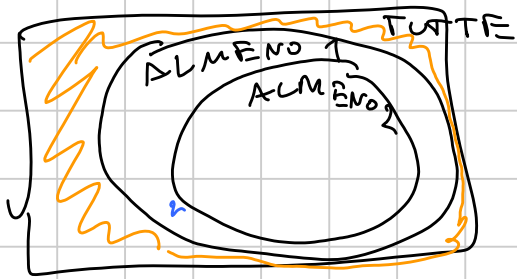
VOGLIAMO OTTENERE UNA PERMUTAZIONE
 DISPARI
 NUMERO DISPARI DI MOSSE



NO

PERMUTAZIONI SENZA PUNTI
 FISSI

$$n! - \binom{n}{1}(n-1)! + \binom{n}{2}(n-2)! - \binom{n}{3}(n-3)! + \dots + \binom{n}{1}(n-1)!$$

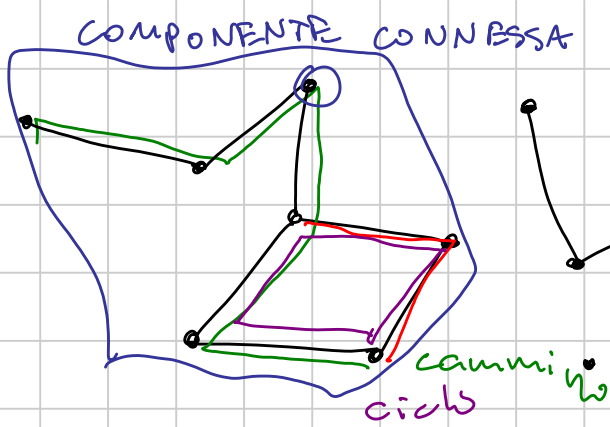


$$n! - \frac{n}{1} (n-1)! + \frac{n(n-1)}{2!} (n-2)! - \frac{n(n-1)(n-2)}{3!} (n-3)! + \dots + (-1)^n \frac{n!}{n!} 0!$$

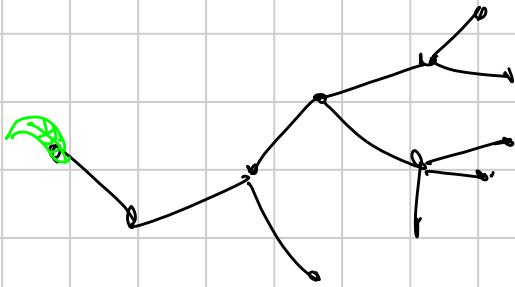
$$n! \left(1 - 1 + \frac{1}{2!} - \frac{1}{3!} + \frac{1}{4!} - \dots + (-1)^n \frac{1}{n!} \right)$$

L'intero più vicino $\frac{1}{n!} \sim \frac{1}{e}$

$$e = 2,71828 \dots$$



vertici
archi



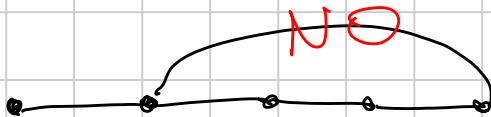
ALBERO

tagliando un arco:
non piú connesso

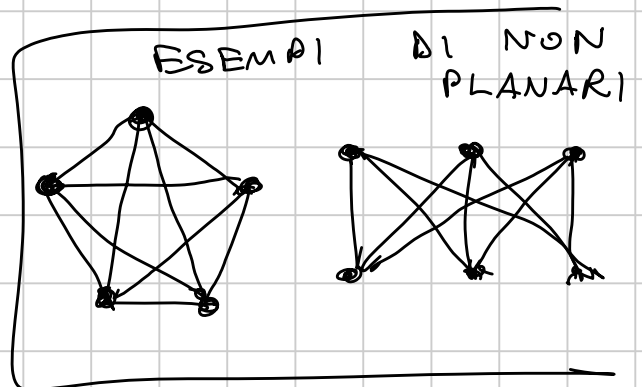
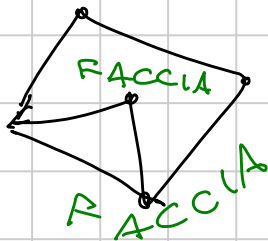
aggiungendo un arco: ottengo un ciclo

$$V = e + 1$$

"edges" ~ archi



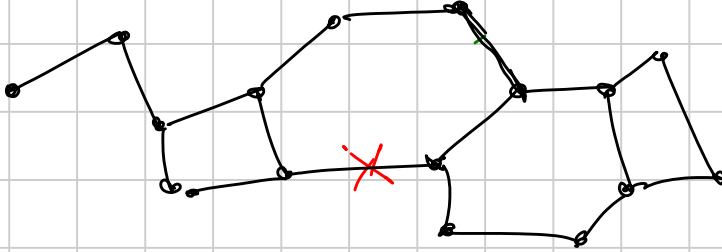
• $\uparrow \geq 0 + 1$



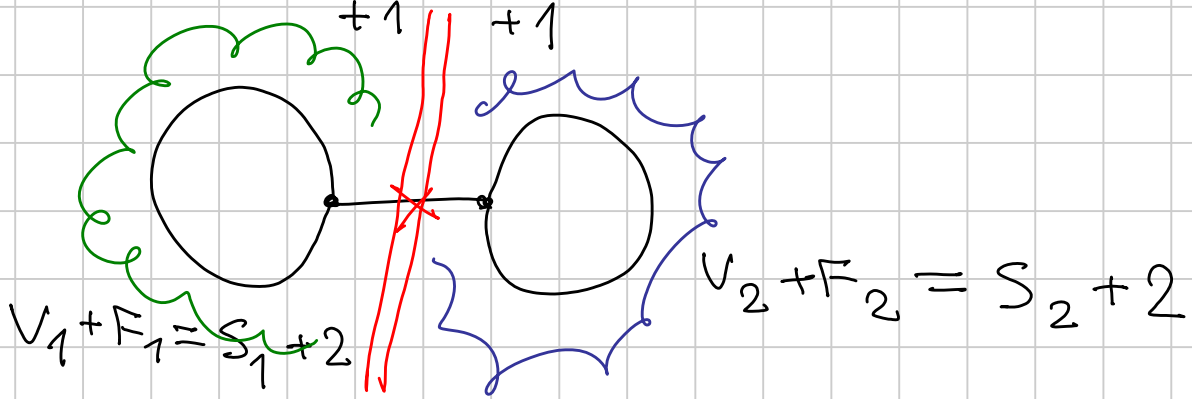
GRAFICI PLANARI CONNESSI (NON VUOTI)

$$V + F = S + 2$$

- $1 + 1 = 0 + 2$



$$V + F = S + 2$$



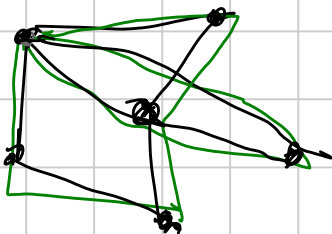
$$(V_1 + V_2) + F = (S_1 + S_2 + 1) + 2$$

SPERIAMO

$$(V_1 + V_2) + (F_1 + F_2 - 1) = (S_1 + S_2 + 1) + 2 + 1$$

IP (NO) $\begin{cases} V_1 + F_1 = S_1 + 2 \\ V_2 + F_2 = S_2 + 2 \end{cases}$

$$(V_1 + V_2) + (F_1 + F_2) = (S_1 + S_2) + 4$$

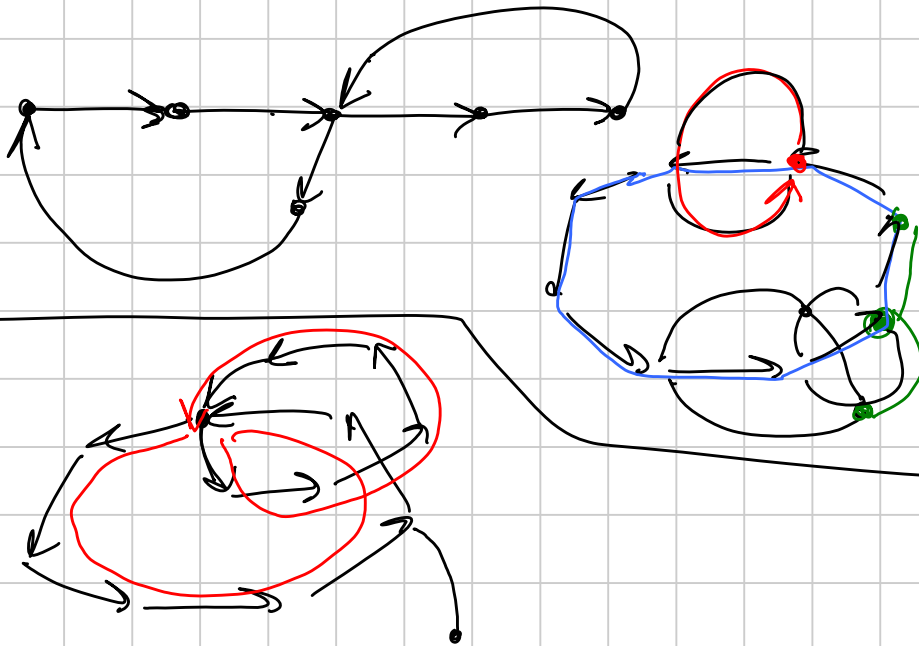


CICLO EULERIANO

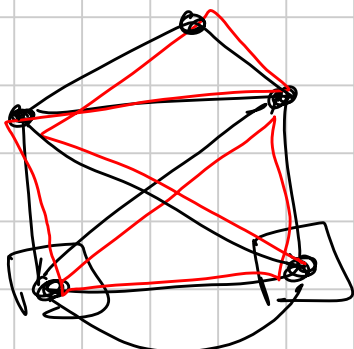
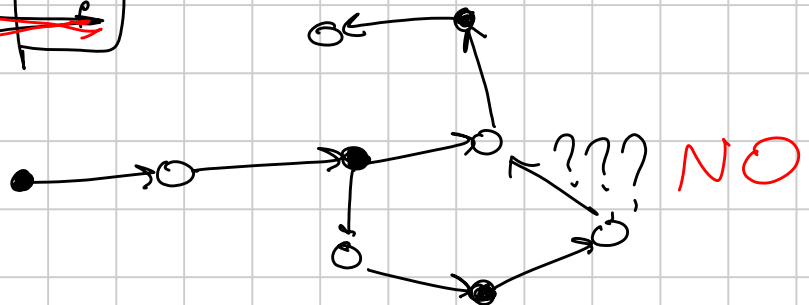
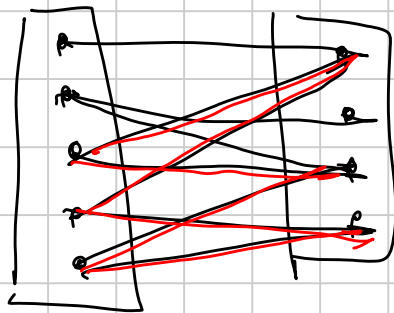
ESISTE SSE

OGNI VERTICE HA ORDINE PARI

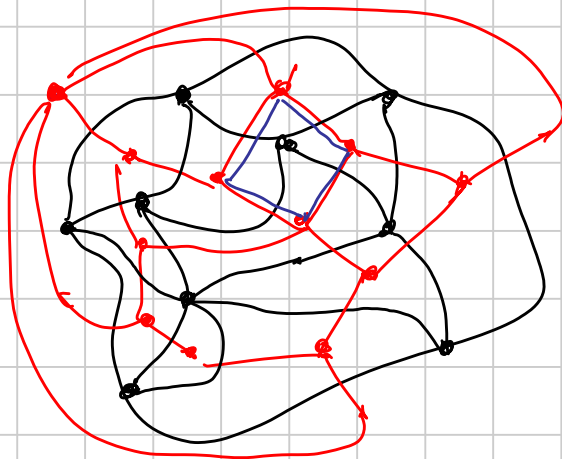
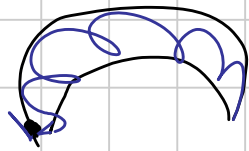
CICLO DI EUCLERO \Rightarrow CARATTERIZZAZIONE
 CICLO: ESCE ED ENTRA IN OGNI VERTICE



IL CICLO PIU' GRANDE
 COMPRENDE TUTTI GLI
 ARCHI



$$\sum_{\text{vertici}} \text{NUMERO DI ARCHI DAL VERTICE} = 2e$$

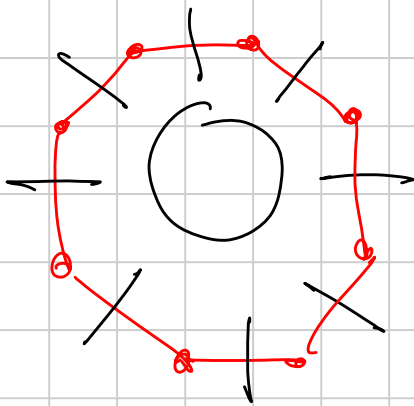
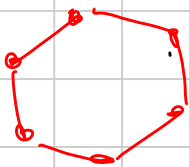


• CARTINA CON 2
COLORI

• FERROVIA LUNGO
(CONFINI)

GRAFO DEI CONFINI = CICLO DI EULERO

GRAFO DEGLI STATI BIPARTITO



$$\sum_{\text{vertici grafo confini}} \text{deg}(v) = \text{PARI} =$$

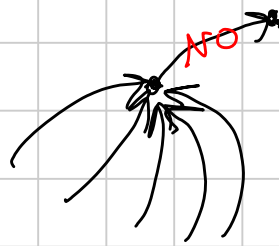
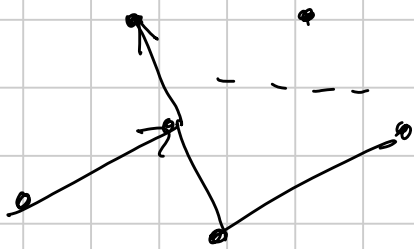
DENTRO IL CICLO

$$= 2 \text{ archi interni } +$$

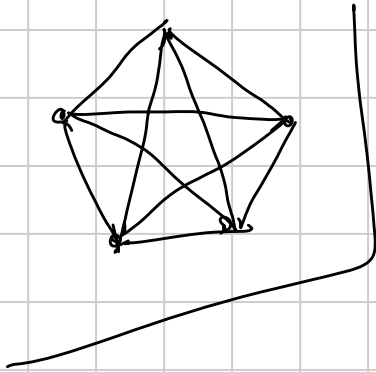
1 archi corrispondenti agli archi del ciclo

SONO PARI

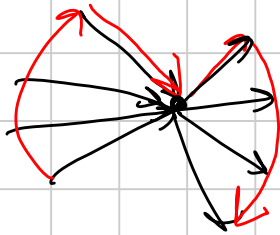
GRAFI ORIENTATI



$\forall p$: PER OGNI COPPIA, ^(a,b) POSSIAMO
 ANDARE da a a b o VICEVERSA



GRAFO COMPLETO
 OGNI ARCO È ORIENTATO
 OGNI VERTICE UNA E UNA
 SOLA VOLTA



INVARIANTI

$$a_1 \quad a_2 \quad a_3 \quad \dots \quad a_{n-1} \quad a_n \in \mathbb{R}$$

↓ ↓ ↓

$$a_1 \quad \frac{a_2 + a_3 + \dots + a_{n-1}}{3} \quad \frac{a_2 + a_3 + \dots + a_{n-1}}{3} \quad \dots \quad \frac{a_2 + a_3 + \dots + a_{n-1}}{3}, a_n$$

SOMMA È UN INVARIANTE

SOMMA DEI QUADRATI CALA

GNOMI CONFORMISTI

Gennaio, Febbraio, - - - Dicembre

Azzurro

Rosa

Alcune coppie: Amici

Si ridipinge da conformisti

Processo è finito

Felicità (gnomo) = amici con casa stesso colore → amici con casa dell'altro colore

$$F = \sum_{\text{gnomi}} f(\text{gnomo})$$

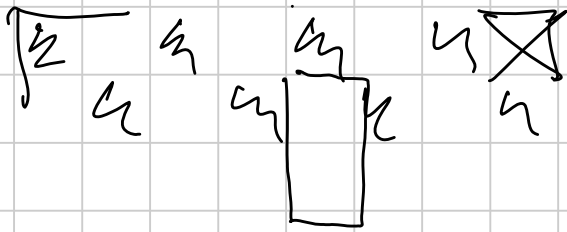
$$- \quad \rightarrow \quad +$$

$$m > n$$

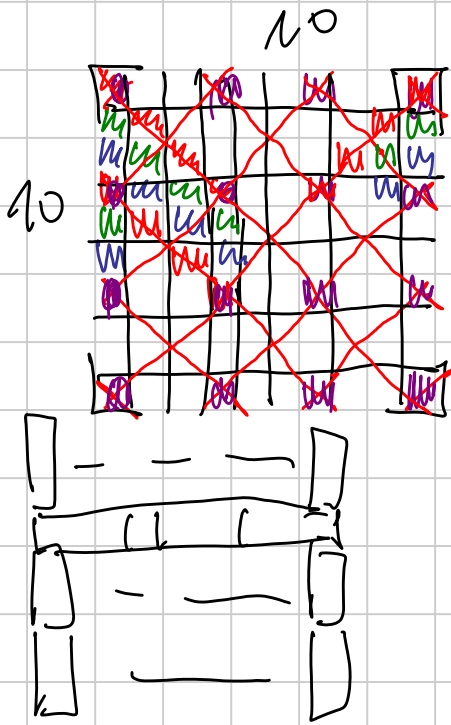
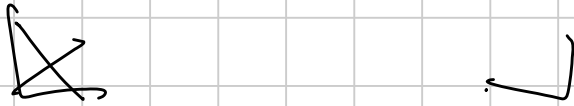
$$+2m - 2n = 2(m - n) > 0$$

$$F \leq \sum_{\text{gnomi}} \text{No amici (gnomo)}$$

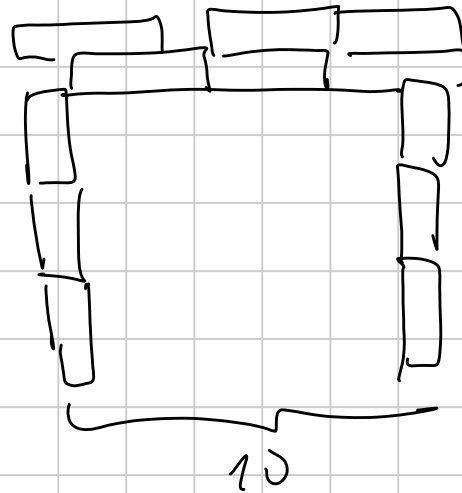
Prima o poi non si ridipinge più



NO



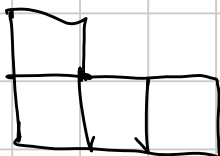
34 33 33



2011 x 2011



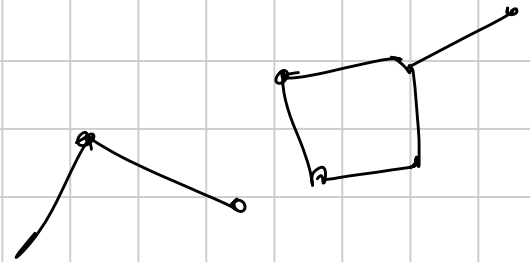
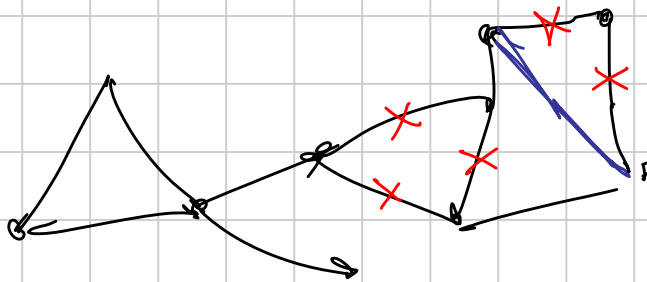
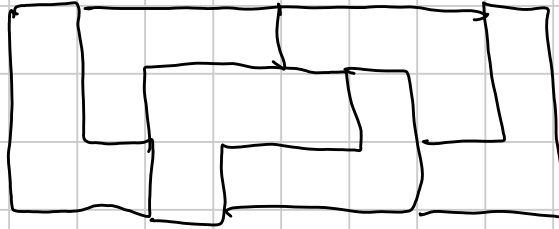
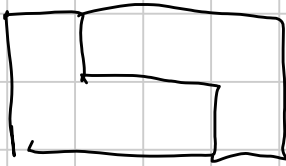
2011 NO



$4a + 9b$
 > 23

$(m, n) = 1$
 $mn - m - n$ *ultim*

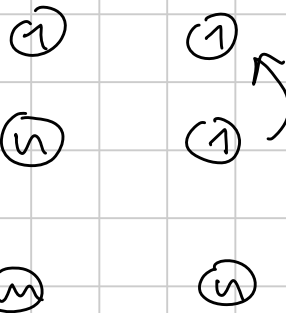
un lato pari \rightarrow strisce perpendicolari a esso
 stesso numero (pari) caselle bianche e nere



Alla fine sappiamo quanti vertici
 con arco uscente

Quanti archi

Sappiamo mosse fatte
 parità



$m \neq n$

$m > n$
 \downarrow
 $n \quad n$

perdente : qualunque mossa risulti in una config vincente per l'avversario

vincente:

7 mosse che risulta
in una config perdente
per l'avversario